

# Certification of the translation from SVM to JVM

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## Contents

<b>1</b>	<b>Java Source Language Some Auxiliary Definitions</b>	<b>2</b>
<b>2</b>	<b>unique</b>	<b>2</b>
<b>3</b>	<b>More about Maps</b>	<b>3</b>
<b>4</b>	<b>Java types</b>	<b>4</b>
<b>5</b>	<b>Class Declarations and Programs</b>	<b>5</b>
<b>6</b>	<b>Relations between Java Types</b>	<b>7</b>
<b>7</b>	<b>Java Values</b>	<b>12</b>
<b>8</b>	<b>Program State</b>	<b>13</b>
<b>9</b>	<b>Java Virtual Machine State of the JVM</b>	<b>16</b>
<b>10</b>	<b>Frame Stack</b>	<b>16</b>
<b>11</b>	<b>Exceptions</b>	<b>17</b>
<b>12</b>	<b>Primitive operators for SVM and JVM</b>	<b>18</b>
<b>13</b>	<b>Instructions of the JVM</b>	<b>18</b>
<b>14</b>	<b>JVM Instruction Semantics</b>	<b>20</b>
<b>15</b>	<b>Exception handling in the JVM</b>	<b>27</b>
<b>16</b>	<b>Program Execution in the JVM</b>	<b>29</b>
<b>17</b>	<b>Normal form values and heap</b>	<b>30</b>

<b>18 State of the SVM</b>	<b>32</b>
18.1 Sizes Table . . . . .	33
18.2 Stack . . . . .	33
18.3 Code Store and SafeImp program . . . . .	33
18.4 Runtime State . . . . .	34
<b>19 Specification Core RTS</b>	<b>35</b>
<b>20 Useful functions and theorems from the Haskell Library or Prelude</b>	<b>54</b>
<b>21 Translation from SVM to JVM</b>	<b>59</b>
21.1 Initialisation code . . . . .	59
21.2 Generic translation functions . . . . .	60
21.3 Specific translation from each SafeImp instruction to bytecode	63
<b>22 Semantics of the SVM instructions</b>	<b>74</b>
<b>23 Certification of the translation from SVM to JVM</b>	<b>79</b>
23.1 Equivalence relations between SVM and JVM states . . . . .	79
<b>24 Lemmas on the static properties of the translation SVM to JVM</b>	<b>84</b>
<b>25 Lemmas on the dynamic properties of the translation SVM to JVM</b>	<b>101</b>
<b>26 Main correctness theorem of the translation SVM to JVM</b>	<b>314</b>

# 1 Java Source Language Some Auxiliary Definitions

```
theory JBasis imports Main begin
```

```
lemmas [simp] = Let-def
```

## 2 unique

```
constdefs
```

```
  unique :: ('a × 'b) list => bool
  unique == distinct ◦ map fst
```

```
lemma fst-in-set-lemma [rule-format (no-asm)]:
  (x, y) : set xys --> x : fst ' set xys
apply (induct-tac xys)
```

**apply** *auto*  
**done**

**lemma** *unique-Nil* [*simp*]: *unique []*  
**apply** (*unfold unique-def*)  
**apply** (*simp (no-asm)*)  
**done**

**lemma** *unique-Cons* [*simp*]: *unique ((x,y)#l) = (unique l & (!y. (x,y) ~: set l))*  
**apply** (*unfold unique-def*)  
**apply** (*auto dest: fst-in-set-lemma*)  
**done**

**lemma** *unique-append* [*rule-format (no-asm)*]: *unique l' ==> unique l -->*  
*(!(x,y):set l. !(x',y'):set l'. x' ~ = x) --> unique (l @ l')*  
**apply** (*induct-tac l*)  
**apply** (*auto dest: fst-in-set-lemma*)  
**done**

**lemma** *unique-map-inj* [*rule-format (no-asm)*]:  
*unique l --> inj f --> unique (map (%(k,x). (f k, g k x)) l)*  
**apply** (*induct-tac l*)  
**apply** (*auto dest: fst-in-set-lemma simp add: inj-eq*)  
**done**

### 3 More about Maps

**lemma** *map-of-SomeI* [*rule-format (no-asm)*]:  
*unique l --> (k, x) : set l --> map-of l k = Some x*  
**apply** (*induct-tac l*)  
**apply** *auto*  
**done**

**lemma** *Ball-set-table'*:  
*( $\forall (x,y) \in \text{set } l. P x y$ ) --> ( $\forall x. \forall y. \text{map-of } l x = \text{Some } y \text{ --> } P x y$ )*  
**apply**(*induct-tac l*)  
**apply**(*simp-all (no-asm)*)  
**apply** *safe*  
**apply** *auto*  
**done**  
**lemmas** *Ball-set-table = Ball-set-table' [THEN mp]*

**lemma** *table-of-remap-SomeD* [*rule-format (no-asm)*]:  
*map-of (map ( $\lambda((k,k'),x). (k,(k',x))$ ) t) k = Some (k',x) -->*  
*map-of t (k, k') = Some x*  
**apply** (*induct-tac t*)  
**apply** *auto*  
**done**

**end**

## 4 Java types

**theory** *Type* imports *JBasis* **begin**

— typedecl cnam changed to:

**types** *cnam* = *string*

— exceptions

**datatype**

*xcpt*

= *NullPointer*

| *ClassCast*

| *OutOfMemory*

| *ArrayIndexOutOfBounds*

| *NegativeArraySize*

| *ArrayStore*

— class names

**datatype** *cname*

= *Object*

| *Xcpt xcpt*

| *Cname cnam*

— variable or field name

— typedecl vnam changed to

**types** *vnam* = *string*

— method name

— typedecl mname changed to

**types** *mname* = *string*

— names for *This* pointer and local/field variables

**datatype** *vname*

= *This*

| *VName vnam*

— primitive type, cf. 4.2

**datatype** *prim-ty*

= *Void* — 'result type' of void methods

| *Boolean*

| *Integer*

| *RetA nat* — bytecode return addresses

— reference type, cf. 4.3

**datatype** *ref-ty*

= *NullT* — null type, cf. 4.1

| *ClassT cname* — class type  
 | *ArrayT ty* — array type  
 — any type, cf. 4.1  
**and** *ty*  
 = *PrimT prim-ty* — primitive type  
 | *RefT ref-ty* — reference type

### syntax

*NT* :: *ty*  
*Class* :: *cname* => *ty*  
*RA* :: *nat* => *ty*  
*Array* :: *ty* => *ty* (-.[ ] [90] 90)

### translations

*NT* == *RefT NullT*  
*Class C* == *RefT (ClassT C)*  
*T.[ ]* == *RefT (ArrayT T)*  
*RA pc* == *PrimT (RetA pc)*

### consts

*the-Class* :: *ty* => *cname*  
*isArray* :: *ty* => *bool*

### reodef *the-Class* { }

*the-Class (Class C)* = *C*  
*the-Class (T.[ ])* = *Object*

### reodef *isArray* { }

*isArray (T.[ ])* = *True*  
*isArray T* = *False*

end

## 5 Class Declarations and Programs

**theory** *Decl* **imports** *Type* **begin**

### types

*fdecl* = *vname* × *ty* — field declaration, cf. 8.3 (, 9.3)

*sig* = *mname* × *ty list* — signature of a method, cf. 8.4.2

*'c mdecl* = *sig* × *ty* × *'c* — method declaration in a class

*'c class* = *cname* × *fdecl list* × *'c mdecl list*  
 — class = superclass, fields, methods

$'c\ cdecl = cname \times 'c\ class$  — class declaration, cf. 8.1

$'c\ prog = 'c\ cdecl\ list$  — program

### translations

$fdecl \leq (type)\ vname \times ty$   
 $sig \leq (type)\ mname \times ty\ list$   
 $mdecl\ c \leq (type)\ sig \times ty \times c$   
 $class\ c \leq (type)\ cname \times fdecl\ list \times (c\ mdecl)\ list$   
 $cdecl\ c \leq (type)\ cname \times (c\ class)$   
 $prog\ c \leq (type)\ (c\ cdecl)\ list$

### constdefs

$class :: 'c\ prog \Rightarrow (cname \rightarrow 'c\ class)$   
 $class \equiv map-of$

$is-class :: 'c\ prog \Rightarrow cname \Rightarrow bool$   
 $is-class\ G\ C \equiv class\ G\ C \neq None$

**lemma**  $finite-is-class: finite\ \{C.\ is-class\ G\ C\}$

**apply**  $(unfold\ is-class-def\ class-def)$

**apply**  $(fold\ dom-def)$

**apply**  $(rule\ finite-dom-map-of)$

**done**

### consts

$is-type :: 'c\ prog \Rightarrow ty \Rightarrow bool$   
 $isrtype :: 'c\ prog \Rightarrow ref-ty \Rightarrow bool$

### primrec

$is-type\ G\ (PrimT\ T) = True$   
 $is-type\ G\ (RefT\ T) = isrtype\ G\ T$   
 $isrtype\ G\ (NullT) = True$   
 $isrtype\ G\ (ClassT\ C) = is-class\ G\ C$   
 $isrtype\ G\ (ArrayT\ T) = is-type\ G\ T$

### consts

$is-RA :: ty \Rightarrow bool$

**recdef**  $is-RA\ \{\}$

$is-RA\ (RA\ pc) = True$

$is-RA\ t = False$

**end**

## 6 Relations between Java Types

**theory** *TypeRel* **imports** *Decl* **begin**

— direct subclass, cf. 8.1.3

**inductive**

*subcls1* :: 'c prog => [cname, cname] => bool (- ⊢ - < C1 - [71,71,71] 70)  
**for** *G* :: 'c prog

**where**

*subcls1I*:  $\llbracket \text{class } G \ C = \text{Some } (D, \text{rest}); C \neq \text{Object} \rrbracket \implies G \vdash C < C1D$

**abbreviation**

*subcls* :: 'c prog => [cname, cname] => bool (- ⊢ - ≤ C - [71,71,71] 70)  
**where**  $G \vdash C \leq C \ D \equiv (\text{subcls1 } G)^{**} \ C \ D$

**lemma** *subcls1D*:

$G \vdash C < C1D \implies C \neq \text{Object} \wedge (\exists fs \ ms. \text{class } G \ C = \text{Some } (D, fs, ms))$

**apply** (*erule subcls1.cases*)

**apply** *auto*

**done**

**lemma** *subcls1-def2*:

*subcls1* *G* =  $(\lambda C \ D. (C, D) \in$   
 $(\text{SIGMA } C: \{C. \text{is-class } G \ C\} . \{D. C \neq \text{Object} \wedge \text{fst } (\text{the } (\text{class } G \ C)) = D\}))$   
**by** (*auto simp add: is-class-def expand-fun-eq dest: subcls1D intro: subcls1I*)

**lemma** *finite-subcls1*: *finite*  $\{(C, D). \text{subcls1 } G \ C \ D\}$

**apply** (*simp add: subcls1-def2 del: mem-Sigma-iff*)

**thm** *finite-SigmaI*

**apply** (*rule finite-SigmaI [OF finite-is-class]*)

**thm** *finite-subset*

**apply** (*rule-tac B = {fst (the (class G C))} in finite-subset*)

**apply** *auto*

**done**

**lemma** *subcls-is-class*:  $(\text{subcls1 } G)^{++} \ C \ D \implies \text{is-class } G \ C$

**apply** (*unfold is-class-def*)

**apply** (*erule tranclp-trans-induct*)

**apply** (*auto dest!: subcls1D*)

**done**

**lemma** *subcls-is-class2* [*rule-format (no-asm)*]:

$G \vdash C \leq C \ D \implies \text{is-class } G \ D \longrightarrow \text{is-class } G \ C$

**apply** (*unfold is-class-def*)

**apply** (*erule rtranclp-induct*)

**apply** (*drule-tac [2] subcls1D*)

**apply** *auto*

**done**

**constdefs**

```

class-rec :: 'c prog ⇒ cname ⇒ 'a ⇒
  (cname ⇒ fdecl list ⇒ 'c mdecl list ⇒ 'a
   ⇒ 'a) ⇒ 'a
class-rec G == wfrec {(C, D). (subcls1 G) ^--1 C D}
  (λr C t f. case class G C of
    None ⇒ arbitrary
  | Some (D, fs, ms) ⇒
    f C fs ms (if C = Object then t else r D t f))

```

**lemma** *class-rec-lemma*:  $wfP ((subcls1 G) ^--1) \implies class\ G\ C = Some\ (D, fs, ms)$   
 $\implies$

```

class-rec G C t f = f C fs ms (if C=Object then t else class-rec G D t f)
by (simp add: class-rec-def wfrec [to-pred, where r=(subcls1 G) ^--1, simpli-
fied]
  cut-apply [where r={(C, D). subcls1 G D C}, simplified, OF subcls1I])

```

**definition**

```

wf-class G = wfP ((subcls1 G) ^--1)

```

**lemma** *class-rec-func* [code func]:

```

class-rec G C t f = (if wf-class G then
  (case class G C
    of None ⇒ arbitrary
  | Some (D, fs, ms) ⇒ f C fs ms (if C = Object then t else class-rec G D t f))
  else class-rec G C t f)

```

**proof** (cases wf-class G)

**case** *False* **then show** ?thesis **by auto**

**next**

**case** *True*

**from** ⟨wf-class G⟩ **have** wf:  $wfP ((subcls1 G) ^--1)$

**unfolding** wf-class-def .

**show** ?thesis

**proof** (cases class G C)

**case** *None*

**with** wf **show** ?thesis

**by** (simp add: class-rec-def wfrec [to-pred, where r=(subcls1 G) ^--1, sim-  
plified]

cut-apply [where r={(C, D).subcls1 G D C}, simplified, OF subcls1I])

**next**

**case** (Some x) **show** ?thesis

**proof** (cases x)

**case** (fields D fs ms)

**then have** is-some:  $class\ G\ C = Some\ (D, fs, ms)$  **using** Some **by simp**

**note** class-rec = class-rec-lemma [OF wf is-some]



```

    show ?thesis unfolding class-rec by (simp add: is-some)
  qed
  qed
  qed

```

**consts**

```

method  :: 'c prog × cname => ( sig  → cname × ty × 'c)
method' :: 'c prog × cname => ( sig  → cname × ty × 'c)
field   :: 'c prog × cname => ( vname → cname × ty   )
fields  :: 'c prog × cname => ((vname × cname) × ty) list

```

— methods of a class, with inheritance, overriding and hiding, cf. 8.4.6

```

defs method-def: method ≡ λ(G,C). class-rec G C empty (λC fs ms ts.
    ts ++ map-of (map (λ(s,m). (s,(C,m))) ms))

```

— methods of a class, with out inheritance, overriding and hiding

```

defs method'-def: method' ≡ λ(G,C). map-of (map (λ(s,m). (s,(C,m)))
    (snd (snd (the (class G C))))))

```

```

lemma method-rec-lemma: [[class G C = Some (D,fs,ms); wfP ((subcls1 G) ^--1)]]
==>
  method (G,C) = (if C = Object then empty else method (G,D)) ++
  map-of (map (λ(s,m). (s,(C,m))) ms)
apply (unfold method-def)
apply (simp split del: split-if)
apply (erule (1) class-rec-lemma [THEN trans])
apply auto
done

```

— list of fields of a class, including inherited and hidden ones

```

defs fields-def: fields ≡ λ(G,C). class-rec G C [] (λC fs ms ts.
  map (λ(fn,ft). ((fn,C),ft)) fs @ ts)

```

```

lemma fields-rec-lemma: [[class G C = Some (D,fs,ms); wfP ((subcls1 G) ^--1)]]
==>
  fields (G,C) =
  map (λ(fn,ft). ((fn,C),ft)) fs @ (if C = Object then [] else fields (G,D))
apply (unfold fields-def)
apply (simp split del: split-if)
apply (erule (1) class-rec-lemma [THEN trans])
apply auto
done

```

```

defs field-def: field == map-of o (map (λ((fn,fd),ft). (fn,(fd,ft)))) o fields

```

**lemma** *table-of-remap-SomeD* [*rule-format (no-asm)*]:  
*map-of (map (λ((k,k'),x). (k,(k',x))) t) k = Some (k',x) -->*  
*map-of t (k, k') = Some x*  
**apply** (*induct-tac t*)  
**apply** *auto*  
**done**

**lemma** *field-fields*:  
*field (G,C) fn = Some (fd, fT) ==> map-of (fields (G,C)) (fn, fd) = Some fT*  
**apply** (*unfold field-def*)  
**apply** (*rule table-of-remap-SomeD*)  
**apply** *simp*  
**done**

— widening, viz. method invocation conversion, cf. 5.3 i.e. sort of syntactic subtyping

**inductive**  
*widen :: 'c prog => [ty , ty ] => bool (- ⊢ - ≼ - [71,71,71] 70)*  
**for** *G :: 'c prog*  
**where**  
*refl [intro!, simp]: G ⊢ T ≼ T — identity conv., cf. 5.1.1*  
*| subcls : G ⊢ C ≼ C D ==> G ⊢ Class C ≼ Class D*  
*| null [intro!]: G ⊢ NT ≼ RefT R*

**lemmas** *refl = HOL.refl*

— casting conversion, cf. 5.5 / 5.1.5  
— left out casts on primitive types

**inductive**  
*cast :: 'c prog => [ty , ty ] => bool (- ⊢ - ≼? - [71,71,71] 70)*  
**for** *G :: 'c prog*  
**where**  
*widen: G ⊢ C ≼ D ==> G ⊢ C ≼? D*  
*| subcls: G ⊢ D ≼ C C ==> G ⊢ Class C ≼? Class D*

**lemma** *widen-PrimT-RefT [iff]*: (*G ⊢ PrimT pT ≼ RefT rT*) = *False*  
**apply** (*rule iffI*)  
**apply** (*erule widen.cases*)  
**apply** *auto*  
**done**

**lemma** *widen-RefT*: *G ⊢ RefT R ≼ T ==> ∃ t. T = RefT t*  
**apply** (*ind-cases G ⊢ RefT R ≼ T*)  
**apply** *auto*  
**done**

**lemma** *widen-RefT2*:  $G \vdash S \preceq_{\text{RefT}} R \implies \exists t. S =_{\text{RefT}} t$   
**apply** (*ind-cases*  $G \vdash S \preceq_{\text{RefT}} R$ )  
**apply** *auto*  
**done**

**lemma** *widen-Class*:  $G \vdash \text{Class } C \preceq T \implies \exists D. T =_{\text{Class}} D$   
**apply** (*ind-cases*  $G \vdash \text{Class } C \preceq T$ )  
**apply** *auto*  
**done**

**lemma** *widen-Class-NullT* [*iff*]:  $(G \vdash \text{Class } C \preceq NT) = \text{False}$   
**apply** (*rule iffI*)  
**apply** (*ind-cases*  $G \vdash \text{Class } C \preceq NT$ )  
**apply** *auto*  
**done**

**lemma** *widen-Class-Class* [*iff*]:  $(G \vdash \text{Class } C \preceq \text{Class } D) = (G \vdash C \preceq C D)$   
**apply** (*rule iffI*)  
**apply** (*ind-cases*  $G \vdash \text{Class } C \preceq \text{Class } D$ )  
**apply** (*auto elim: widen.subcls*)  
**done**

**lemma** *widen-NT-Class* [*simp*]:  $G \vdash T \preceq NT \implies G \vdash T \preceq \text{Class } D$   
**by** (*ind-cases*  $G \vdash T \preceq NT$ , *auto*)

**lemma** *cast-PrimT-RefT* [*iff*]:  $(G \vdash \text{PrimT } pT \preceq_{\text{?}} \text{RefT } rT) = \text{False}$   
**apply** (*rule iffI*)  
**apply** (*erule cast.cases*)  
**apply** *auto*  
**done**

**lemma** *cast-RefT*:  $G \vdash C \preceq_{\text{?}} \text{Class } D \implies \exists rT. C =_{\text{RefT}} rT$   
**apply** (*erule cast.cases*)  
**apply** *simp* **apply** (*erule widen.cases*)  
**apply** *auto*  
**done**

**theorem** *widen-trans*[*trans*]:  $\llbracket G \vdash S \preceq U; G \vdash U \preceq T \rrbracket \implies G \vdash S \preceq T$

**proof** –

**assume**  $G \vdash S \preceq U$  **thus**  $\wedge T. G \vdash U \preceq T \implies G \vdash S \preceq T$

**proof** *induct*

**case** (*refl*  $T T'$ ) **thus**  $G \vdash T \preceq T'$  .

**next**

**case** (*subcls*  $C D T$ )

**then obtain**  $E$  **where**  $T =_{\text{Class}} E$  **by** (*blast dest: widen-Class*)

**with** *subcls* **show**  $G \vdash \text{Class } C \preceq T$  **by** *auto*

**next**

**case** (*null*  $R RT$ )

```

    then obtain rt where  $RT = RefT\ rt$  by (blast dest: widen-RefT)
    thus  $G \vdash NT \preceq RT$  by auto
  qed
qed
end

```

## 7 Java Values

**theory** *Value* imports *Type* begin

— typed decl *loc'* locations, i.e. abstract references on objects

**types** *loc'* = *nat* — locations, i.e. abstract references on objects

**datatype** *loc*  
 = *XcptRef* *xcpt* — special locations for pre-allocated system exceptions  
 | *Loc* *loc'* — usual locations (references on objects)

**datatype** *val*  
 = *Unit* — dummy result value of void methods  
 | *Null* — null reference  
 | *Bool* *bool* — Boolean value  
 | *Intg* *int* — integer value, name *Intg* instead of *Int* because of clash with  
 HOL/Set.thy  
 | *Addr* *loc* — addresses, i.e. locations of objects  
 | *RetAddr* *nat* — return address of JSR instruction, for bytecode only

**consts**

```

the-Bool :: val => bool
the-Intg  :: val => int
the-Addr  :: val => loc
the-RetAddr :: val => nat

```

**primrec**

```
the-Bool (Bool b) = b
```

**primrec**

```
the-Intg (Intg i) = i
```

**primrec**

```
the-Addr (Addr a) = a
```

**primrec**

```
the-RetAddr (RetAddr r) = r
```

**consts**

*defpval* :: *prim-ty* => *val* — default value for primitive types  
*default-val* :: *ty* => *val* — default value for all types

**primrec**

*defpval* *Void* = *Unit*  
*defpval* *Boolean* = *Bool False*  
*defpval* *Integer* = *Intg 0*  
*defpval* (*RetA pc*) = *RetAddr pc*

**primrec**

*default-val* (*PrimT pt*) = *defpval pt*  
*default-val* (*RefT r*) = *Null*

**end**

## 8 Program State

**theory** *State* **imports** *TypeRel Value* **begin**

**types**

*fields-* = (*vname* × *cname* → *val*) — field name, defining class to value  
*entries-* = *nat* → *val* — array index to value

**datatype**

*heap-entry* = *Obj cname fields-* — class instance with class name and fields  
| *Arr ty nat entries-* — array with component type, length, and entries

**constdefs**

*obj-ty* :: *heap-entry* ⇒ *ty*  
*obj-ty entry* ≡ *case entry of Obj C fs* ⇒ *Class C* | *Arr T len entries* ⇒ *T*.[]

*init-vars* :: (*'a* × *ty*) *list* => (*'a* → *val*)  
*init-vars* ≡ *map-of* ∘ *map* ( $\lambda(n, T). (n, \text{default-val } T)$ )

**consts**

*the-obj* :: *heap-entry* ⇒ *cname* × *fields-*  
*the-arr* :: *heap-entry* ⇒ *ty* × *nat* × *entries-*

**recldef** *the-obj* {}

*the-obj* (*Obj C fs*) = (*C, fs*)

**recldef** *the-arr* {}

*the-arr* (*Arr T len entries*) = (*T, len, entries*)

**types** *ahheap* = *loc* → *heap-entry* — "heap" used in a translation below

*locals* = *vname* → *val* — simple state, i.e. variable contents

*state* = *ahheap* × *locals* — heap, local parameter including This

*xstate* = *xcpt option* × *state* — state including exception information

**syntax**

$heap \quad :: \text{state} \Rightarrow \text{aheap}$   
 $locals \quad :: \text{state} \Rightarrow \text{locals}$   
 $Norm \quad \quad :: \text{state} \Rightarrow \text{xstate}$

**translations**

$heap \quad \Rightarrow \text{fst}$   
 $locals \Rightarrow \text{snd}$   
 $Norm \ s \ == \ (None, s)$

**constdefs**

$\text{new-Addr} \quad :: \text{aheap} \Rightarrow \text{loc} \times \text{xcpt option}$   
 $\text{new-Addr} \ h \ == \ \text{SOME} \ (a, x). \ (h \ a = \text{None} \wedge \ x = \text{None}) \ | \ x = \text{Some} \ \text{OutOfMemory}$

$\text{raise-if} \quad :: \text{bool} \Rightarrow \text{xcpt} \Rightarrow \text{xcpt option} \Rightarrow \text{xcpt option}$   
 $\text{raise-if} \ c \ x \ xo \ == \ \text{if} \ c \wedge \ (xo = \text{None}) \ \text{then} \ \text{Some} \ x \ \text{else} \ xo$

$np \quad \quad :: \text{val} \Rightarrow \text{xcpt option} \Rightarrow \text{xcpt option}$   
 $np \ v \ == \ \text{raise-if} \ (v = \text{Null}) \ \text{NullPointer}$

$c\text{-hupd} \quad :: \text{aheap} \Rightarrow \text{xstate} \Rightarrow \text{xstate}$   
 $c\text{-hupd} \ h' \ == \ \lambda(xo, (h, l)). \ \text{if} \ xo = \text{None} \ \text{then} \ (None, (h', l)) \ \text{else} \ (xo, (h, l))$

$\text{cast-ok} \quad :: \ 'c \ \text{prog} \Rightarrow \ \text{cname} \Rightarrow \ \text{aheap} \Rightarrow \ \text{val} \Rightarrow \ \text{bool}$   
 $\text{cast-ok} \ G \ C \ h \ v \ == \ v = \text{Null} \vee \ G \vdash \text{obj-ty} \ (\text{the} \ (h \ (\text{the-Addr} \ v))) \preceq \ \text{Class} \ C$

**lemma**  $\text{obj-ty-def2}$   $[\text{simp}]$ :  $\text{obj-ty} \ (\text{Obj} \ C \ fs) = \ \text{Class} \ C$   
**apply**  $(\text{unfold} \ \text{obj-ty-def})$   
**apply**  $(\text{simp} \ (\text{no-asm}))$   
**done**

**lemma**  $\text{obj-ty-def3}$   $[\text{simp}]$ :  $\text{obj-ty} \ (\text{Arr} \ T \ \text{len} \ \text{entries}) = \ T.[]$   
**by**  $(\text{unfold} \ \text{obj-ty-def}) \ \text{simp}$

**lemma**  $\text{raise-if-True}$   $[\text{simp}]$ :  $\text{raise-if} \ \text{True} \ x \ y \neq \ \text{None}$   
**apply**  $(\text{unfold} \ \text{raise-if-def})$   
**apply**  $\text{auto}$   
**done**

**lemma** *raise-if-False* [*simp*]: *raise-if False x y = y*  
**apply** (*unfold raise-if-def*)  
**apply** *auto*  
**done**

**lemma** *raise-if-Some* [*simp*]: *raise-if c x (Some y) ≠ None*  
**apply** (*unfold raise-if-def*)  
**apply** *auto*  
**done**

**lemma** *raise-if-Some2* [*simp*]:  
*raise-if c z (if x = None then Some y else x) ≠ None*  
**apply** (*unfold raise-if-def*)  
**apply** (*induct-tac x*)  
**apply** *auto*  
**done**

**lemma** *raise-if-SomeD* [*rule-format (no-asm)*]:  
*raise-if c x y = Some z → c ∧ Some z = Some x | y = Some z*  
**apply** (*unfold raise-if-def*)  
**apply** *auto*  
**done**

**lemma** *raise-if-NoneD* [*rule-format (no-asm)*]:  
*raise-if c x y = None → ¬ c ∧ y = None*  
**apply** (*unfold raise-if-def*)  
**apply** *auto*  
**done**

**lemma** *np-NoneD* [*rule-format (no-asm)*]:  
*np a' x' = None → x' = None ∧ a' ≠ Null*  
**apply** (*unfold np-def raise-if-def*)  
**apply** *auto*  
**done**

**lemma** *np-None* [*rule-format (no-asm), simp*]: *a' ≠ Null → np a' x' = x'*  
**apply** (*unfold np-def raise-if-def*)  
**apply** *auto*  
**done**

**lemma** *np-Some* [*simp*]: *np a' (Some xc) = Some xc*  
**apply** (*unfold np-def raise-if-def*)  
**apply** *auto*  
**done**

**lemma** *np-Null* [*simp*]: *np Null None = Some NullPointer*  
**apply** (*unfold np-def raise-if-def*)  
**apply** *auto*  
**done**

```

lemma np-Addr [simp]: np (Addr a) None = None
apply (unfold np-def raise-if-def)
apply auto
done

```

```

lemma np-raise-if [simp]: (np Null (raise-if c xc None)) =
  Some (if c then xc else NullPointer)
apply (unfold raise-if-def)
apply (simp (no-asm))
done

```

```

syntax (xsymbols)
  conf    :: 'c prog => aheap => val => ty => bool
            (-, - ⊢ - :: ≤ - [51,51,51,51] 50)

```

```

end

```

## 9 Java Virtual Machine State of the JVM

```

theory JVMState
imports State
begin

```

For object initialization, we tag each location with the current init status. The tags use an extended type system for object initialization (that gets reused in the BV).

We have either

- usual initialized types, or
- a class that is not yet initialized and has been created by a *New* instruction at a certain line number, or
- a partly initialized class (on which the next super class constructor has to be called). We store the name of the class the super class constructor has to be called of.

```

datatype init-ty = Init ty | UnInit cname nat | PartInit cname

```

## 10 Frame Stack

```

types

```



*opstack* = *val list*  
*locvars* = *val list*  
*p-count* = *nat*  
*ref-upd* = (*val* × *val*)  
*sheap* = *cname* × *vname* → *val*

*frame* = *opstack* ×  
*locvars* ×  
*cname* ×  
*sig* ×  
*p-count* ×  
*ref-upd*

- operand stack
- local variables (including this pointer and method parameters)
- name of class where current method is defined
- method name + parameter types
- program counter within frame
- ref update for obj init proof

## 11 Exceptions

### constdefs

*raise-system-xcpt* :: *bool* ⇒ *xcpt* ⇒ *val option*  
*raise-system-xcpt* *b x* ≡ *if b then Some (Addr (XcptRef x)) else None*

- redefines `State.new_Addr`:

*new-Addr* :: *aheap* ⇒ *loc* × *val option*  
*new-Addr* *h* ≡ *let (a, x) = State.new-Addr h*  
*in (a, raise-system-xcpt (x ≈ None) OutOfMemory)*

### types

*init-heap* = *loc* ⇒ *init-ty*  
— type tags to track init status of objects

*jvm-state* = *val option* × *sheap* × *aheap* × *init-heap* × *frame list*  
— exception flag, static heap, heap, tag heap, frames

a new, blank object with default values in all fields:

### constdefs

*blank* :: '*c prog* ⇒ *cname* ⇒ *heap-entry*  
*blank* *G C* ≡ *Obj C (init-vars (fields(G, C)))*

*blank-arr* :: *ty* ⇒ *nat* ⇒ *heap-entry*  
*blank-arr* *T len* ≡ *Arr T len (λx. if x < len then Some (default-val T) else None)*

*start-heap* :: '*c prog* ⇒ *aheap*  
*start-heap* *G* ≡ *empty (XcptRef NullPointer ↦ blank G (Xcpt NullPointer))*

```

      (XcptRef ClassCast ↦ blank G (Xcpt ClassCast))
      (XcptRef OutOfMemory ↦ blank G (Xcpt OutOfMemory))
      (XcptRef ArrayIndexOutOfBounds ↦ blank G (Xcpt ArrayIndex-
OutOfBounds))
      (XcptRef NegativeArraySize ↦ blank G (Xcpt NegativeArraySize))
      (XcptRef ArrayStore ↦ blank G (Xcpt ArrayStore))

```

```

start-sheap :: 'c prog ⇒ sheap
start-sheap G ≡ empty

```

```

start-iheap :: init-heap
start-iheap ≡ (((((λx. arbitrary)
  (XcptRef NullPointer := Init (Class (Xcpt NullPointer))))
  (XcptRef ClassCast := Init (Class (Xcpt ClassCast))))
  (XcptRef OutOfMemory := Init (Class ((Xcpt OutOfMemory))))
  (XcptRef ArrayIndexOutOfBounds := Init (Class (Xcpt ArrayIndex-
OutOfBounds))))
  (XcptRef NegativeArraySize := Init (Class (Xcpt NegativeArraySize))))
  (XcptRef ArrayStore := Init (Class (Xcpt ArrayStore))))

```

end

## 12 Primitive operators for SVM anf JVM

```

theory BinOP
imports Main
begin

```

Primitive operators

```

datatype PrimOp = Add | Substract | Times | Divide | LessThan | LessEqual
                | Equal | GreaterThan | GreaterEqual | NotEqual

```

end

## 13 Instructions of the JVM

```

theory JVMInstructions
imports JVMState BinOP begin

```

Apply Binary Operation

```

consts
  applyBinOp :: [PrimOp, val, val] ⇒ val

```

**primrec**

$applyBinOp\ Equal\ b1\ b2 = (if\ (the-Intg(b1) = the-Intg(b2))\ then\ Intg\ 1\ else\ Intg\ 0)$   
 $applyBinOp\ NotEqual\ b1\ b2 = (if\ (the-Intg(b1) \neq the-Intg(b2))\ then\ Intg\ 1\ else\ Intg\ 0)$   
 $applyBinOp\ GreaterEqual\ b1\ b2 = (if\ (the-Intg(b1) \geq the-Intg(b2))\ then\ Intg\ 1\ else\ Intg\ 0)$   
 $applyBinOp\ GreaterThan\ b1\ b2 = (if\ (the-Intg(b1) > the-Intg(b2))\ then\ Intg\ 1\ else\ Intg\ 0)$   
 $applyBinOp\ LessThan\ b1\ b2 = (if\ (the-Intg(b1) < the-Intg(b2))\ then\ Intg\ 1\ else\ Intg\ 0)$   
 $applyBinOp\ LessEqual\ b1\ b2 = (if\ (the-Intg(b1) \leq the-Intg(b2))\ then\ Intg\ 1\ else\ Intg\ 0)$   
 $applyBinOp\ Add\ b1\ b2 = Intg\ (the-Intg(b1) + the-Intg(b2))$   
 $applyBinOp\ Subtract\ b1\ b2 = Intg\ (the-Intg(b1) - the-Intg(b2))$   
 $applyBinOp\ Times\ b1\ b2 = Intg\ (the-Intg(b1) * the-Intg(b2))$   
 $applyBinOp\ Divide\ b1\ b2 = Intg\ (the-Intg(b1) div the-Intg(b2))$

Apply ifcmp Generic for integers

**consts**

$applyIf :: [PrimOp, val, val] \Rightarrow bool$

**primrec**

$applyIf\ Equal\ b1\ b2 = (the-Intg(b1) = the-Intg(b2))$   
 $applyIf\ NotEqual\ b1\ b2 = (the-Intg(b1) \neq the-Intg(b2))$   
 $applyIf\ GreaterEqual\ b1\ b2 = (the-Intg(b1) \geq the-Intg(b2))$   
 $applyIf\ GreaterThan\ b1\ b2 = (the-Intg(b1) > the-Intg(b2))$   
 $applyIf\ LessThan\ b1\ b2 = (the-Intg(b1) < the-Intg(b2))$   
 $applyIf\ LessEqual\ b1\ b2 = (the-Intg(b1) \leq the-Intg(b2))$

**datatype**

$instr = Load\ nat$	— load from local variable
$Store\ nat$	— store into local variable
$LitPush\ val$	— push a literal (constant)
$New\ cname$	— create object
$Getfield\ vname\ cname$	— Fetch field from object
$Putfield\ vname\ cname$	— Set field in object
$Checkcast\ cname$	— Check whether object is of given type
$Invoke\ cname\ mname\ (ty\ list)$	— inv. instance meth of an object
$Invoke-special\ cname\ mname\ ty\ list$	— no dynamic type lookup, for constructors
$Return$	— return from method
$Return-Void$	— return from void method
$Pop$	— pop top element from opstack
$Dup$	— duplicate top element of opstack

<i>Dup2</i>	— duplicate two element of opstack
<i>Dup-x1</i>	— duplicate next to top element
<i>Dup-x2</i>	— duplicate 3rd element
<i>Swap</i>	— swap top and next to top element
<i>IAdd</i>	— integer addition
<i>Goto int</i>	— goto relative address
<i>Ifcmpeq int</i>	— branch if int/ref comparison succeeds
<i>Throw</i>	— throw top of stack as exception
<i>Jsr int</i>	— jump to subroutine
<i>Ret nat</i>	— return from subroutine
<i>ArrLoad</i>	— load indexed entry from array
<i>ArrStore</i>	— store value into indexed array entry
<i>ArrLength</i>	— get array length
<i>ArrNew ty</i>	— create new 1-dimensional array
<i>Invoke-static cname mname (ty list)</i>	— invoke a class (static) method
<i>Getstatic vname cname</i>	— get value of static field
<i>Putstatic vname cname</i>	— set value of static field
<i>Tableswitch int int (int list)</i>	— jump according to a table
<i>BinOp PrimOp</i>	— Binary Operation
<i>Ifcmp PrimOp int</i>	— Generic Branch

## types

*bytecode* = *instr list*  
*exception-entry* = *p-count* × *p-count* × *p-count* × *cname*  
— start-pc, end-pc, handler-pc, exception type  
*exception-table* = *exception-entry list*  
*jvm-method* = *nat* × *nat* × *bytecode* × *exception-table*  
*jvm-prog* = *jvm-method prog*

end

## 14 JVM Instruction Semantics

**theory** *JVMExecInstr*  
**imports** *JVMInstructions*  
**begin**

the method name of constructors:

**consts**  
*init* :: *mname*

replace a by b in l:

**constdefs**  
*replace* :: 'a ⇒ 'a ⇒ 'a list ⇒ 'a list  
*replace a b l* == *map* ( $\lambda x. \text{if } x = a \text{ then } b \text{ else } x$ ) *l*

some lemmas about replace

**lemma** *replace-removes-elim*:  
 $a \neq b \implies a \notin \text{set } (\text{replace } a \ b \ l)$   
**by** (*unfold replace-def*) *auto*

**lemma** *replace-length* [*simp*]:  
 $\text{length } (\text{replace } a \ b \ l) = \text{length } l$  **by** (*simp add: replace-def*)

**lemma** *replace-Nil* [*iff*]:  
 $\text{replace } x \ y \ [] = []$  **by** (*simp add: replace-def*)

**lemma** *replace-Cons*:  
 $\text{replace } x \ y \ (l \# ls) = (\text{if } l = x \ \text{then } y \ \text{else } l) \# (\text{replace } x \ y \ ls)$   
**by** (*simp add: replace-def*)

**lemma** *replace-map*:  
 $\text{inj } f \implies \text{replace } (f \ x) \ (f \ y) \ (\text{map } f \ l) = \text{map } f \ (\text{replace } x \ y \ l)$   
**apply** (*induct l*)  
**apply** (*simp add: replace-def*)  
**apply** (*simp add: replace-def*)  
**apply** *clarify*  
**apply** (*drule injD, assumption*)  
**apply** *simp*  
**done**

**lemma** *replace-id*:  
 $x \notin \text{set } l \vee x = y \implies \text{replace } x \ y \ l = l$   
**apply** (*induct l*)  
**apply** (*auto simp add: replace-def*)  
**done**

single execution step for each instruction:

**consts**  
 $\text{exec-instr} :: [\text{instr}, \text{jvm-prog}, \text{sheap}, \text{aheap}, \text{init-heap}, \text{opstack}, \text{locvars},$   
 $\text{cname}, \text{sig}, \text{p-count}, \text{ref-upd}, \text{frame list}] \implies \text{jvm-state}$

**primrec**  
 $\text{exec-instr } (\text{Load } \text{idx}) \ G \ \text{shp } \text{hp } \text{ihp } \text{stk } \text{vars } \text{Cl } \text{sig } \text{pc } z \ \text{frs} =$   
 $(\text{None}, \text{shp}, \text{hp}, \text{ihp}, ((\text{vars } ! \ \text{idx}) \# \ \text{stk}, \ \text{vars}, \ \text{Cl}, \ \text{sig}, \ \text{pc}+1, \ z) \# \ \text{frs})$

$\text{exec-instr } (\text{Store } \text{idx}) \ G \ \text{shp } \text{hp } \text{ihp } \text{stk } \text{vars } \text{Cl } \text{sig } \text{pc } z \ \text{frs} =$   
 $(\text{None}, \text{shp}, \text{hp}, \text{ihp}, (\text{tl } \text{stk}, \ \text{vars}[\text{idx} := \text{hd } \text{stk}], \ \text{Cl}, \ \text{sig}, \ \text{pc}+1, \ z) \# \ \text{frs})$

$\text{exec-instr } (\text{LitPush } v) \ G \ \text{shp } \text{hp } \text{ihp } \text{stk } \text{vars } \text{Cl } \text{sig } \text{pc } z \ \text{frs} =$   
 $(\text{None}, \text{shp}, \text{hp}, \text{ihp}, (v \# \ \text{stk}, \ \text{vars}, \ \text{Cl}, \ \text{sig}, \ \text{pc}+1, \ z) \# \ \text{frs})$

$\text{exec-instr } (\text{New } C) \ G \ \text{shp } \text{hp } \text{ihp } \text{stk } \text{vars } \text{Cl } \text{sig } \text{pc } z \ \text{frs} =$   
 $(\text{let } (\text{oref}, \text{xp}') = \text{new-Addr } \text{hp};$   
 $\text{hp}' = \text{if } \text{xp}' = \text{None} \ \text{then } \text{hp}(\text{oref} \mapsto \text{blank } G \ C) \ \text{else } \text{hp};$   
 $\text{ihp}' = \text{if } \text{xp}' = \text{None} \ \text{then } \text{ihp}(\text{oref} := \text{UnInit } C \ \text{pc}) \ \text{else } \text{ihp};$

```

    stk'      = if xp'=None then (Addr oref)#stk else stk;
    pc'      = if xp'=None then pc+1 else pc
  in (xp',shp, hp', ihp', (stk', vars, Cl, sig, pc', z)#frs))

```

```

exec-instr (Getfield F C) G shp hp ihp stk vars Cl sig pc z frs =
  (let oref = hd stk;
    xp'    = raise-system-xcpt (oref=None) NullPointer;
    (oc,fs) = the-obj (the(hp(the-Addr oref)));
    stk'   = if xp'=None then the(fs(F,C))#(tl stk) else tl stk;
    pc'    = if xp'=None then pc+1 else pc
  in (xp',shp, hp, ihp, (stk', vars, Cl, sig, pc', z)#frs))

```

```

exec-instr (Putfield F C) G shp hp ihp stk vars Cl sig pc z frs =
  (let (fval,oref) = (hd stk, hd(tl stk));
    xp'          = raise-system-xcpt (oref=None) NullPointer;
    a            = the-Addr oref;
    (oc,fs)     = the-obj (the(hp a));
    hp'         = if xp'=None then hp(a ↦ Obj oc (fs((F,C) ↦ fval)))
                  else hp;
    pc'         = if xp'=None then pc+1 else pc
  in (xp',shp, hp', ihp, (tl (tl stk), vars, Cl, sig, pc', z)#frs))

```

```

exec-instr (Checkcast C) G shp hp ihp stk vars Cl sig pc z frs =
  (let oref = hd stk;
    xp'    = raise-system-xcpt (¬cast-ok G C hp oref) ClassCast;
    stk'   = if xp'=None then stk else tl stk;
    pc'    = if xp'=None then pc+1 else pc
  in (xp',shp, hp, ihp, (stk', vars, Cl, sig, pc', z)#frs))

```

```

exec-instr (Invoke C mn ps) G shp hp ihp stk vars Cl sig pc z frs =
  (let n = length ps;
    args = take n stk;
    oref = stk!n;
    xp'  = raise-system-xcpt (oref=None) NullPointer;
    dynT = the-Class (obj-ty (the(hp (the-Addr oref))));
    (dc,mh,mxs,mxl,c) = the (method (G,dynT) (mn,ps));
    frs' = (if xp'=None then
      [([],oref#(rev args)@(replicate mxl arbitrary),dc,(mn,ps),0,arbitrary)]
      else [])
  in (xp',shp, hp, ihp, frs'@(stk, vars, Cl, sig, pc, z)#frs))

```

- Because exception handling needs the pc of the Invoke instruction,
- Invoke doesn't change *stk* and *pc* yet (*Return* does that).

```

exec-instr (Invoke-special C mn ps) G shp hp ihp stk vars Cl sig pc z frs =
  (let n = length ps;
    args = take n stk;
    oref = stk!n;

```

```

addr = the-Addr oref;
x' = raise-system-xcpt (oref=None) NullPointer;
dynT = the-Class (obj-ty (the(hp addr)));
(dc,mh,mxs,mxl,c) = the (method (G,C) (mn,ps));
(addr',x'') = new-Addr hp;
xp' = if x' = None then x'' else x';
hp' = if xp' = None then hp(addr' ↦ blank G dynT) else hp;
ihp' = if C = Object then ihp(addr':= Init (Class dynT))
      else ihp(addr':= PartInit C);
ihp'' = if xp' = None then ihp' else ihp;
z' = if C = Object then (Addr addr', Addr addr') else (Addr addr', Null);
frs' = (if xp'=None then
        [([],(Addr addr')#(rev args)@(replicate mxl arbitrary),dc,(mn,ps),0,z')]
        else [])
in (xp',shp, hp', ihp'', frs'@(stk, vars, Cl, sig, pc, z)#frs)

```

```

exec-instr Return G shp hp ihp stk0 vars Cl sig0 pc z0 frs =
(if frs=[] then (None, shp, hp, ihp, []))
else let
  val = hd stk0;
  (mn,pt) = sig0;
  (stk,loc,C,sig,pc,z) = hd frs;
  (b,c) = z0;
  (a,c') = z;
  n = length pt;
  addr = stk!n;
  stk' = if mn=init then val#replace addr c stk else val#stk;
  loc' = if mn=init then replace addr c loc else loc;
  z' = if mn=init ∧ z = (addr,Null) then (a,c) else z
in (None, shp, hp, ihp, (stk',loc',C,sig,pc+1,z')#tl frs)

```

— Return drops arguments from the caller’s stack and increases  
— the program counter in the caller

—  $z$  is only updated if we are in a constructor and have initialized the  
— same reference as the constructor in the frame above (otherwise we are  
— in the last constructor of the init chain)

```

exec-instr Return-Void G shp hp ihp stk0 vars Cl sig0 pc z0 frs =
(if frs=[] then (None, shp, hp, ihp, []))
else let
  (mn,pt) = sig0;
  (stk,loc,C,sig,pc,z) = hd frs;
  (b,c) = z0;
  (a,c') = z;
  n = length pt;
  addr = stk!n;

```

$stk'$  = if  $mn=init$  then replace  $addr\ c\ stk$  else  $stk$ ;  
 $loc'$  = if  $mn=init$  then replace  $addr\ c\ loc$  else  $loc$ ;  
 $z'$  = if  $mn=init \wedge z = (addr, Null)$  then  $(a, c)$  else  $z$   
in  $(None, shp, hp, ihp, (stk', loc', C, sig, pc+1, z') \# tl\ frs)$

— All items on the current method's operand stack are discarded

*exec-instr Pop*  $G\ shp\ hp\ ihp\ stk\ vars\ Cl\ sig\ pc\ z\ frs =$   
 $(None, shp, hp, ihp, (tl\ stk, vars, Cl, sig, pc+1, z) \# frs)$

*exec-instr Dup*  $G\ shp\ hp\ ihp\ stk\ vars\ Cl\ sig\ pc\ z\ frs =$   
 $(None, shp, hp, ihp, (hd\ stk \# stk, vars, Cl, sig, pc+1, z) \# frs)$

*exec-instr Dup2*  $G\ shp\ hp\ ihp\ stk\ vars\ Cl\ sig\ pc\ z\ frs =$   
 $(None, shp, hp, ihp, (hd\ stk \# hd\ (tl\ stk) \# stk, vars, Cl, sig, pc+1, z) \# frs)$

*exec-instr Dup-x1*  $G\ shp\ hp\ ihp\ stk\ vars\ Cl\ sig\ pc\ z\ frs =$   
 $(None, shp, hp, ihp, (hd\ stk \# hd\ (tl\ stk) \# hd\ stk \# (tl\ (tl\ stk)),$   
 $vars, Cl, sig, pc+1, z) \# frs)$

*exec-instr Dup-x2*  $G\ shp\ hp\ ihp\ stk\ vars\ Cl\ sig\ pc\ z\ frs =$   
 $(None, shp, hp, ihp,$   
 $(hd\ stk \# hd\ (tl\ stk) \# (hd\ (tl\ (tl\ stk))) \# hd\ stk \# (tl\ (tl\ (tl\ stk))),$   
 $vars, Cl, sig, pc+1, z) \# frs)$

*exec-instr Swap*  $G\ shp\ hp\ ihp\ stk\ vars\ Cl\ sig\ pc\ z\ frs =$   
 $(let\ (val1, val2) = (hd\ stk, hd\ (tl\ stk))$   
in  $(None, shp, hp, ihp, (val2 \# val1 \# (tl\ (tl\ stk)), vars, Cl, sig, pc+1, z) \# frs)$

*exec-instr IAdd*  $G\ shp\ hp\ ihp\ stk\ vars\ Cl\ sig\ pc\ z\ frs =$   
 $(let\ (val1, val2) = (hd\ stk, hd\ (tl\ stk))$   
in  $(None, shp, hp, ihp, (Intg\ ((the-Intg\ val1) + (the-Intg\ val2)) \# (tl\ (tl\ stk)),$   
 $vars, Cl, sig, pc+1, z) \# frs)$

*exec-instr Ifcmpeq*  $i\ G\ shp\ hp\ ihp\ stk\ vars\ Cl\ sig\ pc\ z\ frs =$   
 $(let\ (val1, val2) = (hd\ stk, hd\ (tl\ stk));$   
 $pc' = if\ val1 = val2\ then\ nat(int\ pc+i)\ else\ pc+1$   
in  $(None, shp, hp, ihp, (tl\ (tl\ stk), vars, Cl, sig, pc', z) \# frs)$

*exec-instr Goto*  $i\ G\ shp\ hp\ ihp\ stk\ vars\ Cl\ sig\ pc\ z\ frs =$   
 $(None, shp, hp, ihp, (stk, vars, Cl, sig, nat(int\ pc+i), z) \# frs)$

*exec-instr Throw*  $G\ shp\ hp\ ihp\ stk\ vars\ Cl\ sig\ pc\ z\ frs =$   
 $(let\ xcpt = raise-system-xcpt\ (hd\ stk = Null)\ NullPointer;$   
 $xcpt' = if\ xcpt = None\ then\ Some\ (hd\ stk)\ else\ xcpt$   
in  $(xcpt', shp, hp, ihp, (stk, vars, Cl, sig, pc, z) \# frs)$

*exec-instr Jsr*  $i\ G\ shp\ hp\ ihp\ stk\ vars\ Cl\ sig\ pc\ z\ frs =$



$(None, shp, hp, ihp, (RetAddr (pc+1)\#stk, vars, Cl, sig, nat(int pc+i), z)\#frs)$

$exec-instr (Ret\ idx)\ G\ shp\ hp\ ihp\ stk\ vars\ Cl\ sig\ pc\ z\ frs =$   
 $(None, shp, hp, ihp, (stk, vars, Cl, sig, the-RetAddr (vars!idx), z)\#frs)$

$exec-instr (ArrLoad)\ G\ shp\ hp\ ihp\ stk\ vars\ Cl\ sig\ pc\ z\ frs =$   
 $(let\ idx = the-Intg (hd\ stk);$   
 $\quad aref = hd (tl\ stk);$   
 $\quad xp'' = raise-system-xcpt (aref=None)\ NullPointer;$   
 $\quad (T, l, en) = the-arr (the(hp(the-Addr\ aref)));$   
 $\quad xp' = if\ xp''=None\ then\ raise-system-xcpt (int\ l \leq idx \vee idx < 0)$   
 $ArrayIndexOutOfBounds$   
 $\quad else\ xp'';$   
 $\quad stk' = the (en (nat\ idx)) \# (tl (tl\ stk));$   
 $\quad pc' = if\ xp'=None\ then\ pc+1\ else\ pc$   
 $in\ (xp', shp, hp, ihp, (stk', vars, Cl, sig, pc', z)\#frs))$

$exec-instr (ArrStore)\ G\ shp\ hp\ ihp\ stk\ vars\ Cl\ sig\ pc\ z\ frs =$   
 $(let\ val = hd\ stk;$   
 $\quad idx = the-Intg (hd (tl\ stk));$   
 $\quad aref = hd (tl (tl\ stk));$   
 $\quad xp'' = raise-system-xcpt (aref=None)\ NullPointer;$   
 $\quad a = the-Addr\ aref;$   
 $\quad (T, l, en) = the-arr (the(hp\ a));$   
 $\quad xp''' = if\ xp''=None\ then\ raise-system-xcpt (int\ l \leq idx \vee idx < 0)\ ArrayIn-$   
 $dexOutOfBounds$   
 $\quad else\ xp'';$   
 $\quad xp' = if\ xp'''=None\ then\ raise-system-xcpt (False)\ ArrayStore\ else\ xp''';$   
 $\quad hp' = if\ xp'=None\ then\ hp(a \mapsto Arr\ T\ l\ (en(nat\ idx \mapsto val)))\ else\ hp;$   
 $\quad pc' = if\ xp'=None\ then\ pc+1\ else\ pc$   
 $in\ (xp', shp, hp', ihp, (tl (tl (tl\ stk)), vars, Cl, sig, pc', z)\#frs))$

$exec-instr (ArrLength)\ G\ shp\ hp\ ihp\ stk\ vars\ Cl\ sig\ pc\ z\ frs =$   
 $(let\ aref = hd\ stk;$   
 $\quad xp' = raise-system-xcpt (aref=None)\ NullPointer;$   
 $\quad a = the-Addr\ aref;$   
 $\quad (T, l, en) = the-arr (the(hp\ a));$   
 $\quad pc' = if\ xp'=None\ then\ pc+1\ else\ pc$   
 $in\ (xp', shp, hp, ihp, (Intg (int\ l)\#(tl\ stk), vars, Cl, sig, pc', z)\#frs))$

$exec-instr (ArrNew\ T)\ G\ shp\ hp\ ihp\ stk\ vars\ Cl\ sig\ pc\ z\ frs =$   
 $(let\ len = the-Intg (hd\ stk);$   
 $\quad xp''' = raise-system-xcpt (len < 0)\ NegativeArraySize;$   
 $\quad (aref, xp'') = new-Addr\ hp;$   
 $\quad xp' = if\ xp'''=None\ then\ xp''\ else\ xp''';$   
 $\quad hp' = if\ xp'=None\ then\ hp(aref \mapsto blank-arr\ T (nat\ len))\ else\ hp;$   
 $\quad ihp' = if\ xp'=None\ then\ ihp(aref := Init (T.\[]))\ else\ ihp;$

```

stk'      = (Addr aref)#tl stk;
pc'       = if xp'=None then pc+1 else pc
in (xp',shp, hp', ihp', (stk', vars, Cl, sig, pc', z)#frs)

```

```

exec-instr (BinOp bop) G shp hp ihp stk vars Cl sig pc z frs =
  (let (val2,val1) = (hd stk, hd (tl stk))
   in (None,shp, hp, ihp, ((applyBinOp bop val1 val2 )#(tl (tl stk))), vars, Cl, sig,
pc + 1, z)#frs))

```

```

exec-instr(Invoke-static C mn ps) G shp hp ihp stk vars Cl sig pc z frs =
  (let n      = length ps;
   args      = take n stk;
   (dc,mh,mxs,mxl,c) = the (method (G,C) (mn,ps));
   frs'      = [[(), (rev args)@(replicate mxl arbitrary),dc,(mn,ps),0,arbitrary]]
   in (None,shp,hp,ihp,frs'@(stk,vars,Cl,sig,pc,z)#frs))

```

```

exec-instr(Tablesitch n m t) G shp hp ihp stk vars Cl sig pc z frs =
  (let val = hd stk;
   v      = the-Intg val;
   pc'    = if (v < n) ∨ (v > m) then pc + nat(t!(nat(m - n + 1))) else pc +
nat(t!(nat(v - n)))
   in (None,shp,hp,ihp,((tl stk),vars,Cl,sig,pc',z)#frs))

```

```

exec-instr (Getstatic F C) G shp hp ihp stk vars Cl sig pc z frs =
  (let v = the (shp (C,F));
   pc' = pc + 1
   in (None,shp, hp, ihp, ((v#stk), vars, Cl, sig, pc', z)#frs))

```

```

exec-instr (Putstatic F C) G shp hp ihp stk vars Cl sig pc z frs =
  (let shp' = shp((C,F) ↦ hd stk );
   pc'     = pc + 1
   in (None, shp', hp, ihp, ((tl stk), vars, Cl, sig, pc', z)#frs))

```

```

exec-instr (Ifcmp bop i) G shp hp ihp stk vars Cl sig pc z frs =
  (let (val1,val2) = (hd stk, hd (tl stk));
   pc'           = if ((applyIf bop val1 val2) = True) then nat(int pc+i) else pc+1
   in (None,shp, hp, ihp, (tl (tl stk), vars, Cl, sig, pc', z)#frs))

```

end

## 15 Exception handling in the JVM

```

theory JVMExceptions
imports JVMInstructions
begin
constdefs
  match-exception-entry :: jvm-prog  $\Rightarrow$  cname  $\Rightarrow$  p-count  $\Rightarrow$  exception-entry  $\Rightarrow$ 
  bool
  match-exception-entry G cn pc ee ==
    let (start-pc, end-pc, handler-pc, catch-type) = ee in
      start-pc  $\leq$  pc  $\wedge$  pc  $<$  end-pc  $\wedge$   $G \vdash cn \preceq C$  catch-type

consts
  match-exception-table :: jvm-prog  $\Rightarrow$  cname  $\Rightarrow$  p-count  $\Rightarrow$  exception-table
     $\Rightarrow$  p-count option

primrec
  match-exception-table G cn pc [] = None
  match-exception-table G cn pc (e#es) = (if match-exception-entry G cn pc e
    then Some (fst (snd (snd e)))
```

*else match-exception-table* *G cn pc es*)

```

consts
  cname-of :: aheap  $\Rightarrow$  val  $\Rightarrow$  cname
  ex-table-of :: jvm-method  $\Rightarrow$  exception-table

translations
  cname-of hp v == the-Class (obj-ty (the (hp (the-Addr v))))
  ex-table-of m == snd (snd (snd m))

consts
  find-handler :: jvm-prog  $\Rightarrow$  val option  $\Rightarrow$  sheap  $\Rightarrow$  aheap  $\Rightarrow$  init-heap  $\Rightarrow$  frame
  list  $\Rightarrow$  jvm-state

primrec
  find-handler G xcpt shp hp ihp [] = (xcpt,shp, hp, ihp, [])
  find-handler G xcpt shp hp ihp (fr#frs) =
    (case xcpt of
      None  $\Rightarrow$  (None, shp, hp, ihp, fr#frs)
    | Some xc  $\Rightarrow$ 
      let (stk, loc, C, sig, pc, r) = fr in
        (case match-exception-table G (cname-of hp xc) pc
          (ex-table-of (snd (snd (the (method (G, C) sig)))))) of
          None  $\Rightarrow$  find-handler G (Some xc) shp hp ihp frs
        | Some handler-pc  $\Rightarrow$  (None, shp, hp, ihp, ([xc], loc, C, sig, handler-pc,
          r)#frs)))

```

Expresses that a value is tagged with an initialized type (only applies to addresses and then only if the heap contains a value for the address)

**constdefs**

*is-init* :: *aheap*  $\Rightarrow$  *init-heap*  $\Rightarrow$  *val*  $\Rightarrow$  *bool*  
*is-init hp ih v*  $\equiv$   
 $\forall loc. v = Addr\ loc \longrightarrow hp\ loc \neq None \longrightarrow (\exists t. ih\ loc = Init\ t)$

System exceptions are allocated in all heaps.

**constdefs**

*preallocated* :: *aheap*  $\Rightarrow$  *init-heap*  $\Rightarrow$  *bool*  
*preallocated hp ihp*  $\equiv \forall x. \exists fs. hp\ (XcptRef\ x) = Some\ (Obj\ (Xcpt\ x)\ fs) \wedge is-init\ hp\ ihp\ (Addr\ (XcptRef\ x))$

**lemma** *preallocatedD* [*simp, dest*]:

*preallocated hp ihp*  $\Longrightarrow \exists fs. hp\ (XcptRef\ x) = Some\ (Obj\ (Xcpt\ x)\ fs) \wedge is-init\ hp\ ihp\ (Addr\ (XcptRef\ x))$   
**by** (*unfold preallocated-def fast*)

**lemma** *preallocatedE* [*elim?*]:

*preallocated hp ihp*  $\Longrightarrow$   
 $(\bigwedge fs. hp\ (XcptRef\ x) = Some\ (Obj\ (Xcpt\ x)\ fs) \Longrightarrow is-init\ hp\ ihp\ (Addr\ (XcptRef\ x))) \Longrightarrow P\ hp\ ihp$   
 $\Longrightarrow P\ hp\ ihp$   
**by** *fast*

**lemma** *cname-of-xcp*:

*raise-system-xcpt b x = Some xcp*  $\Longrightarrow preallocated\ hp\ ihp$   
 $\Longrightarrow cname-of\ hp\ xcp = Xcpt\ x$

**proof** –

**assume** *raise-system-xcpt b x = Some xcp*  
**hence** *xcp = Addr (XcptRef x)*  
**by** (*simp add: raise-system-xcpt-def split: split-if-asm*)

**moreover**

**assume** *preallocated hp ihp*  
**then obtain** *fs* **where** *hp (XcptRef x) = Some (Obj (Xcpt x) fs) ..*  
**ultimately show** *?thesis* **by** *simp*

**qed**

**lemma** *preallocated-start*:

*preallocated (start-heap G) start-iheap*  
**apply** (*unfold preallocated-def*)  
**apply** (*unfold start-heap-def start-iheap-def*)  
**apply** (*rule allI*)  
**apply** (*case-tac x*)  
**apply** (*auto simp add: blank-def is-init-def*)  
**done**

Only program counters that are mentioned in the exception table can be returned by *match-exception-table*:

**lemma** *match-exception-table-in-et*:

*match-exception-table*  $G C pc et = Some pc' \implies \exists e \in set et. pc' = fst (snd (snd e))$

**by** (*induct et*) (*auto split: split-if-asm*)

**end**

## 16 Program Execution in the JVM

**theory** *JVMExec*

**imports** *JVMExecInstr JVMExceptions*

**begin**

**consts**

*exec* :: *jvm-prog*  $\times$  *jvm-state*  $\implies$  *jvm-state option*

— redef only used for pattern matching

**redef** *exec* {}

*exec* ( $G, xp, shp, hp, ihp, []$ ) = *None*

*exec* ( $G, None, shp, hp, ihp, (stk, loc, C, sig, pc, z) \# frs$ ) =  
(*let*

$i = fst (snd (snd (snd (snd (the (method' (G, C) sig)))))) ! pc;$

$(xcpt', shp', hp', ihp', frs') = exec-instr i G shp hp ihp stk loc C sig pc z frs$

*in Some (find-handler G xcpt' shp' hp' ihp' frs')*)

*exec* ( $G, Some xp, hp, ihp, frs$ ) = *None*

**constdefs**

*exec-all* :: [*jvm-prog, jvm-state, jvm-state*]  $\implies$  *bool*  
(- |- - *jvm*  $\rightarrow$  - [61, 61, 61] 60)

$G \mid - s \text{ -jvm} \rightarrow t == (s, t) \in \{(s, t). exec(G, s) = Some t\}^*$

**syntax** (*xsymbols*)

*exec-all* :: [*jvm-prog, jvm-state, jvm-state*]  $\implies$  *bool*  
(-  $\vdash$  - *jvm*  $\rightarrow$  - [61, 61, 61] 60)

The start configuration of the JVM: in the start heap, we call a method  $m$  of class  $C$  in program  $G$ . The *this* pointer of the frame is set to *Null* to simulate a static method invocation.

**constdefs**

*start-state* :: *jvm-prog*  $\Rightarrow$  *cname*  $\Rightarrow$  *mname*  $\Rightarrow$  *jvm-state*

*start-state*  $G C m \equiv$

*let* ( $C', rT, mxs, mxl, ins, et$ ) = *the (method (G, C) (m, [])) in*

```
(None, start-sheap G, start-heap G, start-iheap,
  [([], Null # replicate max arbitrary, C, (m, []), 0, (Null, Null))])
```

**end**

## 17 Normal form values and heap

```
theory SafeHeap
imports Main
begin
```

**types**

```
Location = nat
Constructor = string
FunName = string
```

— Normal form values

```
datatype Val = Loc Location | IntT int | BoolT bool
```

— Destructions of datatype val

```
consts the-IntT :: Val ⇒ int
```

**primrec**

```
the-IntT (IntT i) = i
```

```
consts the-BoolT :: Val ⇒ bool
```

**primrec**

```
the-BoolT (BoolT b) = b
```

— check if is constant bool

```
constdefs isBool :: Val ⇒ bool
```

```
isBool v ≡ (case v of (BoolT -) ⇒ True
  | - ⇒ False )
```

A heap is a partial mapping from locations to cells. But, as it is split into regions, the mapping tells also the region where the cell lives. The second component is the highest live region  $k$ . A consistent heap  $(h, k)$  has cells only in regions  $0 \dots k$ .

**types**

```
Cell = Constructor × Val list
Region = nat
HeapMap = Location → (Region × Cell)
Heap = HeapMap × nat
```

**consts**

```
restrictToRegion :: Heap => Region => Heap (infix ↓ 110)
```

**primrec**

$$(h,k) \downarrow k0 = (\text{let } A = \{ p . p \in \text{dom } h \ \& \ \text{fst } (\text{the } (h \ p)) \leq k0 \} \\ \text{in } (h \mid' A, k0))$$
**constdefs**

$$\text{fresh} :: \text{Location} \Rightarrow \text{HeapMap} \Rightarrow \text{bool}$$

$$\text{fresh } p \ h \equiv p \notin \text{dom } h$$
**constdefs**

$$\text{getFresh} :: \text{HeapMap} \Rightarrow \text{Location}$$

$$\text{getFresh } h \equiv \text{SOME } b. \text{fresh } b \ h$$

'copy' is a runtime support function copying the recursive part of a data structure.

**consts**

$$\text{copy} :: [\text{Heap}, \text{Region}, \text{Location}] \Rightarrow \text{Heap} \times \text{Location}$$
**constdefs**

$$\text{self} :: \text{string} \text{ --- this identifies the topmost region referenced in a function body}$$

$$\text{self} \equiv \text{"self"}$$

The constructor table tells, for each constructor, the number of arguments and a description of each one. The second nat gives the alternative  $0..n - 1$  corresponding to this constructor in every **case** of its type

$$\text{datatype } \text{ArgType} = \text{IntArg} \mid \text{BoolArg} \mid \text{NonRecursive} \mid \text{Recursive}$$
**types**

$$\text{ConstructorTableType} = (\text{Constructor} \times (\text{nat} \times \text{nat} \times \text{ArgType list})) \text{ list}$$

$$\text{ConstructorTableFun} = \text{Constructor} \rightarrow (\text{nat} \times \text{nat} \times \text{ArgType list})$$

This is the constructor table of the Safe expressions semantics. It is assumed to be a constant which somebody else will provide. It is used in the semantic function 'copy'

**consts**

$$\text{ConstructorTable} :: \text{ConstructorTableFun}$$
**constdefs**  $\text{getConstructorCell} :: \text{Cell} \Rightarrow \text{Constructor}$ 

$$\text{getConstructorCell } c \equiv \text{fst } c$$
**constdefs**  $\text{getValuesCell} :: \text{Cell} \Rightarrow \text{Val list}$ 

$$\text{getValuesCell } c \equiv \text{snd } c$$
**constdefs**  $\text{getCell} :: \text{Heap} \Rightarrow \text{Location} \Rightarrow \text{Cell}$ 

$$\text{getCell } h \ l \equiv \text{snd } (\text{the } ((\text{fst } h) \ l))$$
**constdefs**  $\text{getRegion} :: \text{Heap} \Rightarrow \text{Location} \Rightarrow \text{Region}$





```
imports SafeHeap ../JVMSAFE/BinOP
begin
```

## 18.1 Sizes Table

This gives statically inferred information about the maximum number of heap cells, of heap regions, and of stack words needed by the compiled program.

```
types ncell = nat
       sizeRegions = nat
       sizeStackS = nat
```

```
types SizesTable = ncell  $\times$  sizeRegions  $\times$  sizeStackS
```

## 18.2 Stack

```
types
       CodeLabel = nat
       Continuation = Region  $\times$  CodeLabel
```

```
datatype StackObject = Val Val | Reg Region | Cont Continuation
```

The SVM stack may contain normal form values, region arguments for functions or constructors, and continuations. A continuation  $(k_0, p)$  contains a jump  $p$  to a code sequence and an adjustment  $k_0$  for the heap watermark  $k_0$  of the SVM state.

```
types
       Stack = StackObject list
       StackOffset = nat
```

## 18.3 Code Store and SafeImp program

— Items are the components of environments and closures

```
datatype Item = ItemConst Val
             | ItemVar StackOffset
             | ItemRegSelf
```

— The SVM instruction repertory

```
datatype SafeInstr = DECREGION
                 | POPCONT
                 | PUSHCONT CodeLabel
                 | COPY
                 | REUSE
                 | CALL CodeLabel
                 | PRIMOP PrimOp
                 | MATCH StackOffset (CodeLabel list)
```

```

| MATCHD StackOffset (CodeLabel list)
| MATCHN StackOffset nat nat (CodeLabel list)
| BUILDENV (Item list)
| BUILDCLS Constructor (Item list) Item
| SLIDE nat nat

```

```

fun pushcont :: SafeInstr => bool
where
  pushcont (PUSHCONT p) = True
| pushcont -           = False

```

```

fun popcont :: SafeInstr => bool
where
  popcont POPCONT = True
| popcont -       = False

```

A Safe program, when translated into SafeImp, produces four components (1) a map from labels to pairs consisting of a code sequence and a function name. It is given as a list in order to be able to ‘traverse’ the map; (2) a map from function names to pairs consisting of a label and a list of labels. The first points to the starting sequence of the function and the second collects, for each function body, the code labels corresponding to continuations. The map is also given as a list; (3) the code label of the main expression; and (4) a constructor table collecting the properties of all the constructors.

```

types
  CodeSequence = SafeInstr list
  SVMCode      = (CodeLabel × CodeSequence × FunName) list
  ContinuationMap = (FunName × CodeLabel × CodeLabel list) list
  CodeStore      = SVMCode × ContinuationMap
  SafeImpProg    = CodeStore × CodeLabel × ConstructorTableType × SizesTable

```

## 18.4 Runtime State

```

types
  PC = CodeLabel × nat
  SVMState = Heap × Region × PC × Stack

```

```

consts
  incrPC :: PC => PC

```

```

primrec
  incrPC (l,i) = (l,i+1)

```

This is the correspondence between primitive operators in CoreSafe and SafeImp.

```

constdefs

```

```

primops :: string → PrimOp

primops ≡ map-of [( "+", Add),
                  ( "-", Subtract),
                  ( "*", Times),
                  ( "%", Divide),
                  ( "<", LessThan),
                  ( "<=", LessEqual),
                  ( "=", Equal),
                  ( ">", GreaterThan),
                  ( ">=", GreaterEqual)
                ]

— Define primitive operations
consts
  execOp :: [PrimOp, Val, Val] => Val

primrec
  execOp Equal b1 b2 = BoolT (the-IntT(b1) = the-IntT(b2))
  execOp NotEqual b1 b2 = BoolT (the-IntT(b1) ≠ the-IntT(b2))
  execOp GreaterEqual b1 b2 = BoolT (the-IntT(b1) ≥ the-IntT(b2))
  execOp GreaterThan b1 b2 = BoolT (the-IntT(b1) > the-IntT(b2))
  execOp LessThan b1 b2 = BoolT (the-IntT(b1) < the-IntT(b2))
  execOp LessEqual b1 b2 = BoolT (the-IntT(b1) ≤ the-IntT(b2))

  execOp Add b1 b2 = IntT (the-IntT(b1) + the-IntT(b2))
  execOp Subtract b1 b2 = IntT (the-IntT(b1) - the-IntT(b2))
  execOp Times b1 b2 = IntT (the-IntT(b1) * the-IntT(b2))
  execOp Divide b1 b2 = IntT (the-IntT(b1) div the-IntT(b2))

end

```

## 19 Specification Core RTS

```

theory RTSCore
imports ../JVMSAFE/JVMInstructions
begin

```

— Run-time system class name

```

constdefs
  System :: cname
  System ≡ Cname "java/lang/System"
  objectC :: cname
  objectC ≡ Cname "Object"
  heapC :: cname
  heapC ≡ Cname "rtsCore/Heap"

```

```

stackC :: cname
stackC ≡ Cname "rtsCore/Stack-S"
dirCellC :: cname
dirCellC ≡ Cname "rtsCore/DirectoryCell"
consTableC :: cname
consTableC ≡ Cname "rtsCore/ConsTable"
consDataC :: cname
consDataC ≡ Cname "rtsCore/ConsData"
cellC :: cname
cellC ≡ Cname "rtsCore/Cell"

```

— Run-time system method name

```

Init :: mname
Init ≡ "<init>"
exitM :: mname
exitM ≡ "exit"
decregion :: mname
decregion ≡ "decregion"
slide :: mname
slide ≡ "slide"
pushRegion :: mname
pushRegion ≡ "pushRegion"
popRegion :: mname
popRegion ≡ "popRegion"
copyCell :: mname
copyCell ≡ "copyCell"
copyCellAux :: mname
copyCellAux ≡ "copyCellAux"
clone :: mname
clone ≡ "clone"
copyRTS :: mname
copyRTS ≡ "copy"
reserveCell :: mname
reserveCell ≡ "reserveCell"
releaseCell :: mname
releaseCell ≡ "releaseCell"
insertCell :: mname
insertCell ≡ "insertCell"
makeHeap :: mname
makeHeap ≡ "makeHeap"
makeDirectory :: mname
makeDirectory ≡ "makeDirectory"
setMaxSize :: mname
setMaxSize ≡ "setMaxSize"
makeConsTable :: mname
makeConsTable ≡ "makeConsTable"

```

— Run-time system field name

**constdefs**

```
izqf :: vname
izqf ≡ VName "izq"
derf :: vname
derf ≡ VName "der"
idf :: vname
idf ≡ VName "id"
freef :: vname
freef ≡ VName "free"
tipoArgsf :: vname
tipoArgsf ≡ VName "tipoArgs"
Sf :: vname
Sf ≡ VName "S"
topf :: vname
topf ≡ VName "top"
k0f :: vname
k0f ≡ VName "k0"
kf :: vname
kf ≡ VName "k"
regionsf :: vname
regionsf ≡ VName "regions"
safeDirf :: vname
safeDirf ≡ VName "safeDir"
tagLf :: vname
tagLf ≡ VName "tagL"
tagGf :: vname
tagGf ≡ VName "tagG"
nargsf :: vname
nargsf ≡ VName "nargs"
tablef :: vname
tablef ≡ VName "table"
arg1 :: vname
arg1 ≡ VName "arg1"
arg2 :: vname
arg2 ≡ VName "arg2"
arg3 :: vname
arg3 ≡ VName "arg3"
arg4 :: vname
arg4 ≡ VName "arg4"
arg5 :: vname
arg5 ≡ VName "arg5"
arg6 :: vname
arg6 ≡ VName "arg6"
arg7 :: vname
arg7 ≡ VName "arg7"
arg8 :: vname
arg8 ≡ VName "arg8"
```

## Cell Class

— Cell Field Declarations

```
constdefs cellFdecl :: fdecl list
cellFdecl ≡ [(izgf, RefT (ClassT cellC)),
              (derf, RefT (ClassT cellC)),
              (idf, PrimT Integer),
              (tagGf, PrimT Integer),
              (arg1, PrimT Integer),
              (arg2, PrimT Integer),
              (arg3, PrimT Integer),
              (arg4, PrimT Integer),
              (arg5, PrimT Integer),
              (arg6, PrimT Integer),
              (arg7, PrimT Integer),
              (arg8, PrimT Integer)]
```

## ConsTable Class

— ConsTable Field Declarations

```
constdefs consTableFdecl :: fdecl list
consTableFdecl ≡ [(tablef, RefT (ArrayT (RefT (ClassT consDataC))))]
```

## **constdefs**

```
EndFor1 :: int
EndFor1 ≡ 12
bucle1 :: int
bucle1 ≡ - 14
constdefs instMakeConsTable :: instr list
instMakeConsTable ≡
  [Load 0,
   ArrNew (RefT (ClassT consDataC)),
   Putstatic tablef consTableC, (* table = new ConsData[maxCons] *)
   LitPush (Intg 0),
   Store 1, (* i ← 0 *)
   Load 1,
   Getstatic tablef consTableC,
   ArrLength,
   Ifcmp GreaterEqual EndFor1, (* i < table.length ? *)
   Getstatic tablef consTableC,
   Load 1,
   New consDataC,
   Dup,
   Invoke-special consDataC Init [],
   ArrStore,
   Load 1,
   LitPush (Intg 1),
   BinOp Add,
```

```

Store 1,          (* i ← i + 1 *)
Goto bucle1,
Return-Void
]

```

## ConsData Class

— ConsData Field Declarations

```

constdefs consDataFdecl :: fdecl list
consDataFdecl ≡ [(tagLf, PrimT Integer),
(nargsf, PrimT Integer),
(tipoArgsf, RefT (ArrayT (PrimT Integer)))]

```

```

constdefs instInitConsData :: instr list
instInitConsData ≡ [Load 0,          (* Load this *)
LitPush (Intg 8),
ArrNew (PrimT Integer), (* tipoArgs ← new int[8] *)
Putfield tipoArgsf consDataC,
Return-Void]

```

## Stack<sub>S</sub>Class

— Stack<sub>S</sub>FieldDeclarations

```

constdefs stackSFdecl :: fdecl list
stackSFdecl ≡ [(Sf, RefT (ArrayT (PrimT Integer))),(topf, PrimT Integer)]

```

```

constdefs instSetMaxSize :: instr list
instSetMaxSize ≡ [Load 0,
ArrNew (PrimT Integer),
Putstatic Sf stackC,   (* S ← new int[maxSize] *)
LitPush (Intg -1),
Putstatic topf stackC, (* top ← -1 *)
Return-Void
]

```

```

constdefs
EndFor2 :: int
EndFor2 ≡ 16
Bucle2 :: int
Bucle2 ≡ -17

```

```

constdefs instSlide :: instr list
instSlide ≡ [Getstatic topf stackC,
Load 0,          (* Load local0 = m *)
BinOp Subtract,
Load 1,          (* Load local1 = n *)
BinOp Subtract,
LitPush (Intg 1),

```

```

BinOp Add,
Store 2,          (* local2 ← top - m - n + 1 *)
Getstatic topf stackC,
Load 0,
BinOp Substract,
LitPush (Intg 1),
BinOp Add,
Store 3,          (* local3 ← top - m + 1 *)
Load 3,          (* Label Bucle2 *)
Getstatic topf stackC,
Ifcmp GreaterThan EndFor2, (* local3 <= top ? *)
Getstatic Sf stackC,
Load 2,
Getstatic Sf stackC,
Load 3,
ArrLoad,
ArrStore,        (* S[local2] ← S[local3] *)
Load 2,
LitPush (Intg 1),
BinOp Add,
Store 2,          (* local2 ← local2 + 1 *)
Load 3,
LitPush (Intg 1),
BinOp Add,
Store 3,          (* local3 ← local3 + 1 *)
Goto Bucle2,
Getstatic topf stackC, (* Label EndFor2 *)
Load 1,
BinOp Substract,
Putstatic topf stackC, (* top ← top - n *)
Return-Void
]

```

DirectoryCell Class

— DirectoryCell Field Declarations

```

constdefs dirCFdecl :: fdecl list
dirCFdecl ≡ [(safeDirf, RefT (ArrayT (RefT (ClassT cellC))))],
             (freef, RefT (ClassT cellC))]

```

```

constdefs
EndFor3 :: int
EndFor3 ≡ 77
If1 :: int
If1 ≡ 25
If2 :: int
If2 ≡ 19

```



```

If3 :: int
If3 ≡ 32
EndIf2 :: int
EndIf2 ≡ 38
Bucle3 :: int
Bucle3 ≡ -82

```

— Instructions List

```

constdefs instMakeDirectory :: instr list
instMakeDirectory ≡ [
  Load 0,
  ArrNew (Class cellC),
  Putstatic safeDirf dirCellC, (* safeDir ← new Cell[maxCell] *)
  LitPush (Intg 0),
  Store 1, (* i ← 0 *)
  Load 1,
  Getstatic safeDirf dirCellC,
  ArrLength,
  Ifcmp GreaterEqual EndFor3, (* i < safeDir.length ? *)
  Getstatic safeDirf dirCellC,
  Load 1,
  New cellC,
  Dup,
  Invoke-special cellC Init [],
  ArrStore, (* safeDir[i] ← new Cell() *)
  Getstatic safeDirf dirCellC,
  Load 1,
  ArrLoad,
  Load 1,
  Putfield idf cellC, (* safeDir[i].id ← i *)
  Load 1,
  LitPush (Intg 0),
  Ifcmp LessEqual If1, (* i > 0 ? *)
  Load 1,
  Getstatic safeDirf dirCellC,
  ArrLength,
  LitPush (Intg 1),
  BinOp Subtract,
  Ifcmp GreaterThan If2, (* i < safeDir.length - 1 *)
  Getstatic safeDirf dirCellC,
  Load 1,
  ArrLoad,
  Getstatic safeDirf dirCellC,
  Load 1,
  LitPush (Intg 1),
  BinOp Subtract,
  ArrLoad,
  Putfield izqf cellC, (* safeDir[i].izq ← safeDir[i - 1] *)
  Getstatic safeDirf dirCellC,

```

```

Load 1,
ArrLoad,
Getfield izqf cellC,
Getstatic safeDirf dirCellC,
Load 1,
ArrLoad,
Putfield derf cellC,      (* safeDir[i].izq.der ← safeDir[i] *)
Goto EndIf2,
Load 1,
Getstatic safeDirf dirCellC,
ArrLength,
LitPush (Intg 1),
BinOp Subtract,
Ifcmp NotEqual If3,      (* i = safeDir.length - 1 ? *)
Getstatic safeDirf dirCellC,
Load 1,
ArrLoad,
Getstatic safeDirf dirCellC,
Load 1,
LitPush (Intg 1),
BinOp Subtract,
ArrLoad,
Putfield izqf cellC,      (* safeDir[i].izq ← safeDir[i - 1] *)
Getstatic safeDirf dirCellC,
Load 1,
ArrLoad,
Getstatic izqf cellC,
Getstatic safeDirf dirCellC,
Load 1,
ArrLoad,
Putfield derf cellC,      (* safeDir[i].izq.der ← safeDir[i] *)
Getstatic safeDirf dirCellC,
Load 1,
ArrLoad,
Getstatic safeDirf dirCellC,
LitPush (Intg 0),
ArrLoad,
Putfield derf cellC,      (* safeDir[i].der ← safeDir[0] *)
Getstatic safeDirf dirCellC,
LitPush (Intg 0),
ArrLoad,
Getstatic safeDirf dirCellC,
Load 1,
ArrLoad,
Putfield izqf cellC,      (* safeDir[0].izq ← safeDir[i] *)
Load 1,
LitPush (Intg 1),
BinOp Add,
Store 1,                  (* i ← i + 1 *)

```

```

    Goto Bucle3,
    Getstatic safeDirf dirCellC,
    LitPush (Intg 0),
    ArrLoad,
    Putstatic freef cellC,    (* free ← safeDir[0] *)
    Return-Void
]

```

**constdefs**

```

If4 :: int
If4 ≡ 18
EndIf4 :: int
EndIf4 ≡ 3

```

**constdefs** instReserveCell :: instr list

```

instReserveCell ≡ [
    Getstatic freef cellC,
    Getfield derf cellC,
    Store 0,                (* p ← free.der *)
    Load 0,
    Getstatic freef cellC,
    Ifcmpeq If4,            (* p ≠ free ? *)
    Load 0,
    Getfield derf cellC,
    Store 1,                (* p2 ← p.der *)
    Getstatic freef cellC,
    Load 1,
    Putfield derf cellC,    (* free.der ← p2 *)
    Load 1,
    Load 0,
    Getfield izqf cellC,
    Putfield izqf cellC,    (* p2.izq ← p.izq *)
    Load 0,
    Load 0,
    LitPush (Null),
    Dup-x1,
    Putfield derf cellC,    (* p.der ← null *)
    Putfield izqf cellC,    (* p.izq ← null *)
    Goto EndIf4,
    LitPush(Intg 1),
    Invoke-static System exitM [PrimT Integer], (* Abort program *)
    Load 0,
    Getfield idf cellC,
    Return                  (* return p.id *)
]

```

**constdefs** instReleaseCell :: instr list

```

instReleaseCell ≡ [
    Getstatic safeDirf dirCellC,

```

```

Load 0,
ArrLoad,
Store 1,          (* cell ← safeDir[index] *)
Load 1,
Getfield izqf cellC,
Store 2,          (* temp1 ← cell.izq   *)
Load 1,
Getfield derf cellC,
Store 3,          (* temp2 ← cell.der    *)
Load 2,
Load 1,
Getfield derf cellC,
Putfield derf cellC, (* temp1.der ← cell.der *)
Load 3,
Load 1,
Getfield izqf cellC,
Putfield izqf cellC, (* temp2.izq ← cell.izq *)
Load 1,
Getstatic freef dirCellC,
Getfield derf cellC,
Putfield derf cellC, (* cell.der ← free.der *)
Getstatic freef dirCellC,
Getfield derf cellC,
Load 1,
Putfield izqf cellC, (* free.der.izq ← cell *)
Getstatic freef cellC,
Load 1,
Putfield derf cellC, (* free.der ← cell   *)
Load 1,
Getstatic freef dirCellC,
Putfield izqf cellC, (* cell.izq ← free   *)
Return-Void
]

```

— Heap Field Declarations

```

constdefs heapFdecl :: fdecl list
  heapFdecl ≡ [(regionsf, RefT (ArrayT (RefT (ClassT cellC))),
               (kf, PrimT Integer),
               (k0f, PrimT Integer)]

```

```

constdefs instMakeHeap :: instr list
  instMakeHeap ≡ [
    Load 0,
    ArrNew (Class cellC),
    Putstatic regionsf heapC,  (* regions ← new Cell[maxRegions] *)
    LitPush (Intg -1),
    Putstatic kf heapC,        (* k ← -1 *)

```

```

LitPush (Intg -1),
Putstatic k0f heapC,      (* k0 ← -1 *)
Return-Void
]

```

### constdefs

```

bucle5 :: int
bucle5 ≡ 3
endbucle5 :: int
endbucle5 ≡ -4

```

### constdefs instDecregion :: instr list

```

instDecregion ≡ [
  Getstatic kf heapC,
  Getstatic k0f heapC,
  Ifcmp LessEqual bucle5,      (* k > k0 ? *)
  Invoke-static heapC popRegion [], (* call PopRegion() method *)
  Goto endbucle5,
  Return-Void
]

```

### constdefs

```

labelEnd :: int
labelEnd ≡ 4

```

### constdefs instPushRegion :: instr list

```

instPushRegion ≡ [
  Getstatic kf heapC,
  LitPush (Intg 1),
  BinOp Add,
  Putstatic kf heapC, (* k ← k + 1 *)
  Invoke-static dirCellC reserveCell [],
  Store 0, (* cell ← reserveCell() *)
  Getstatic regionsf heapC,
  Getstatic kf heapC,
  Getstatic safeDirf dirCellC,
  Load 0,
  ArrLoad,
  ArrStore, (* regions[k] ← safeDir[cell] *)
  Getstatic regionsf heapC,
  Getstatic kf heapC,
  ArrLoad,
  Getstatic regionsf heapC,
  Getstatic kf heapC,
  ArrLoad,

```

```

Getstatic regionsf heapC,
Getstatic kf heapC,
ArrLoad,
Dup-x1,
Putfield izqf cellC, (* regions[k].izq ← regions[k] *)
Putfield derf cellC, (* regions[k].der ← regions[k].izq *)
Goto labelEnd,
Store 0,             (* Exception e *)
LitPush (Intg 1),
Invoke-static System exitM [PrimT Integer], (* Abort program with error *)
Return-Void
]

```

**constdefs** instPopRegion :: instr list

```

instPopRegion ≡ [
  Getstatic freef dirCellC,
  Getfield izqf cellC,
  Getstatic regionsf heapC,
  Getstatic kf heapC,
  ArrLoad,
  Putfield derf cellC, (* free.izq.der ← regions[k] *)
  Getstatic freef dirCellC,
  Getfield izqf cellC,
  Store 0,             (* c1 ← free.izq *)
  Getstatic regionsf heapC,
  Getstatic kf heapC,
  ArrLoad,
  Getfield izqf cellC,
  Store 1,             (* c2 ← regions[k].izq *)
  Getstatic freef cellC,
  Load 1,
  Putfield izqf cellC, (* free.izq ← c2 *)
  Load 1,
  Getstatic freef cellC,
  Putfield derf cellC, (* c2.der ← free *)
  Getstatic regionsf heapC,
  Getstatic kf heapC,
  ArrLoad,
  Load 0,
  Putfield izqf cellC, (* regions[k].izq ← c1 *)
  Getstatic regionsf heapC,
  Getstatic kf heapC,
  LitPush (Null),
  ArrStore,           (* regions[k] ← NULL *)
  Getstatic kf heapC,
  LitPush(Intg 1),
  BinOp Subtract,
  Putstatic kf heapC, (* k ← k - 1 *)
  Return-Void
]

```

```

]

constdefs instInsertCell :: instr list
instInsertCell ≡ [
  Getstatic safeDirf dirCellC,
  Load 1,
  ArrLoad,
  Store 2,          (* cell ← safeDir[p] *)
  Load 2,
  Getstatic regionsf heapC,
  Load 0,
  Getfield izqf cellC,
  Putfield izqf cellC, (* cell.izq ← regions[r].izq *)
  Getstatic regionsf heapC,
  Load 0,
  ArrLoad,
  Getfield izqf cellC,
  Load 2,
  Putfield derf cellC, (* regions[r].izq.der ← cell *)
  Getstatic regionsf heapC,
  Load 0,
  ArrLoad,
  Load 2,
  Putfield izqf cellC, (* regions[r].izq ← cell *)
  Load 2,
  Getstatic regionsf heapC,
  Load 0,
  ArrLoad,
  Putfield derf cellC, (* cell.der ← regions[r] *)
  Return-Void
]

constdefs instCopyCell :: instr list
instCopyCell ≡ [
  Invoke-static dirCellC reserveCell [],
  Store 1,          (* freshCell ← reserveCell() *)
  Getstatic safeDirf dirCellC,
  Load 1,
  ArrLoad,
  Store 2,          (* target ← safeDir[freshCell] *)
  Getstatic safeDirf dirCellC,
  Load 0,
  ArrLoad,
  Store 3,          (* source ← safeDir[p] *)
  Load 3,
  Load 2,
  Invoke-static dirCellC copyCellAux [Class cellC,
                                     Class cellC], (* copyCellAux(source,target) *)
  Load 0,

```

```

Invoke-static dirCellC releaseCell [PrimT Integer], (* releaseCell(p) *)
Load 2,
Getfield derf cellC,
Load 2,
Putfield izqf cellC,      (* target.der.izq ← target *)
Load 2,
Getfield izqf cellC,
Load 2,
Putfield derf cellC,      (* target.izq.der ← target *)
Load 1,
Return                    (* devuelve freshCell *)
]

```

**constdefs** instClone :: instr list

```

instClone ≡ [
  Invoke-static dirCellC reserveCell [],
  Store 2,          (* freshCell ← reserveCell() *)
  Getstatic safeDirf dirCellC,
  Load 2,
  ArrLoad,
  Store 3,          (* target ← safeDir[freshCell] *)
  Getstatic safeDirf dirCellC,
  Load 1,
  ArrLoad,
  Store 4,          (* source ← safeDir[p] *)
  Load 4,
  Load 3,
  Invoke-static heapC copyCellAux [Class cellC,
                                   Class cellC], (* copyCellAux(source,target) *)
  Load 0,
  Load 2,
  Invoke-static heapC insertCell [PrimT Integer,
                                   PrimT Integer], (* insertCell(reg,freshCell) *)
  Load 2,
  Return           (* devuelve entero freshCell *)
]

```

**constdefs** instCopyCellAux :: instr list

```

instCopyCellAux ≡ [
  Load 1,
  Load 0,
  Getfield derf cellC,
  Putfield derf cellC,      (* c2.der ← c1.der *)
  Load 1,
  Load 0,
  Getfield izqf cellC,
  Putfield izqf cellC,      (* c2.izq ← c1.izq *)
  Load 1,
  Load 0,

```



```

    Getfield tagGf cellC,
    Putfield tagGf cellC,      (* c2.tagG ← c1.tagG *)
    Load 1,
    Load 0,
    Getfield arg1 cellC,
    Putfield arg1 cellC,      (* c2.arg1 ← c1.arg1 *)
    Load 1,
    Load 0,
    Getfield arg2 cellC,
    Putfield arg2 cellC,      (* c2.arg2 ← c1.arg2 *)
    Load 1,
    Load 0,
    Getfield arg3 cellC,
    Putfield arg3 cellC,      (* c2.arg3 ← c1.arg3 *)
    Load 1,
    Load 0,
    Getfield arg4 cellC,
    Putfield arg4 cellC,      (* c2.arg4 ← c1.arg4 *)
    Load 1,
    Load 0,
    Getfield arg5 cellC,
    Putfield arg5 cellC,      (* c2.arg5 ← c1.arg5 *)
    Load 1,
    Load 0,
    Getfield arg6 cellC,
    Putfield arg6 cellC,      (* c2.arg6 ← c1.arg6 *)
    Load 1,
    Load 0,
    Getfield arg7 cellC,
    Putfield arg7 cellC,      (* c2.arg7 ← c1.arg7 *)
    Load 1,
    Load 0,
    Getfield arg8 cellC,
    Putfield arg8 cellC      (* c2.arg8 ← c1.arg8 *)
]

```

```

constdefs
  endBucle6 :: int
  endBucle6 ≡ 72
  If5 :: int
  If5 ≡ 62
  lab1 :: int
  lab1 ≡ 1
  endlab1 :: int
  endlab1 ≡ 49
  lab2 :: int
  lab2 ≡ lab1 + 7

```

```

endlab2 :: int
endlab2 ≡ endlab1 - 7
lab3 :: int
lab3 ≡ lab2 + 7
endlab3 :: int
endlab3 ≡ endlab2 - 7
lab4 :: int
lab4 ≡ lab3 + 7
endlab4 :: int
endlab4 ≡ endlab3 - 7
lab5 :: int
lab5 ≡ lab4 + 7
endlab5 :: int
endlab5 ≡ endlab4 - 7
lab6 :: int
lab6 ≡ lab5 + 7
endlab6 :: int
endlab6 ≡ endlab5 - 7
lab7 :: int
lab7 ≡ lab6 + 7
endlab7 :: int
endlab7 ≡ endlab6 - 7
lab8 :: int
lab8 ≡ lab7 + 7
ldef :: int
ldef ≡ lab8 + 6
bucle6 :: int
bucle6 ≡ -77
constdefs instCopy :: instr list
instCopy ≡ [
  Load 1,
  Load 0,
  Invoke-static heapC clone [PrimT Integer,
    PrimT Integer],
  Store 2,          (* target ← clone (j,b) *)
  Getstatic safeDirf dirCellC,
  Load 0,
  ArrLoad,
  Store 3,          (* source ← safeDir[b] *)
  Getstatic tablef consDataC,
  Load 3,
  Getfield tagGf cellC,
  ArrLoad,
  Store 4,          (* info ← table[source.tagG] *)
  LitPush (Intg 0),
  Store 5,          (* i ← 0 *)
  Load 5,
  Load 4,
  Getfield nargsf consDataC,

```

```

Ifcmp GreaterEqual endBucle6,      (* i < info.nargs ?      *)
Load 4,
Getfield tipoArgsf consDataC,
Load 5,
ArrLoad,          (* Integer Array *)
Store 6,          (* tipo ← info.tipoArgs[i] *)
Load 6,
LitPush (Intg 4),
Ifcmp NotEqual If5,
Getstatic safeDirf dirCellC,
Load 2,
ArrLoad,
Store 7,          (* targ ← safeDir[target] *)
Load 5,          (* Load local5 = i *)
Tableswitch 0 7 [lab1,
                lab2,
                lab3,
                lab4,
                lab5,
                lab6,
                lab7,
                lab8,
                ldef],

Load 7,          (* label 1 *)
Load 3,
Getfield arg1 cellC,
Load 1,
Invoke-static heapC copyRTS [PrimT Integer,
                             PrimT Integer],
Putfield arg1 cellC, (* targ.arg1 ← copy(source.arg1,j) *)
Goto endlab1,     (* end label 1 *)

Load 7,          (* label 2 *)
Load 3,
Getfield arg2 cellC,
Load 1,
Invoke-static heapC copyRTS [PrimT Integer,
                             PrimT Integer],
Putfield arg2 cellC, (* targ.arg2 ← copy(source.arg2,j) *)
Goto endlab2,     (* end label 2 *)

Load 7,          (* label 3 *)
Load 3,
Getfield arg3 cellC,
Load 1,
Invoke-static heapC copyRTS [PrimT Integer,

```

```
PrimT Integer],
Putfield arg3 cellC, (* targ.arg3 ← copy(source.arg3,j) *)
Goto endlab3,      (* end label 3 *)
```

```
Load 7,           (* label 4 *)
Load 3,
Getfield arg4 cellC,
Load 1,
Invoke-static heapC copyRTS [PrimT Integer,
PrimT Integer],
Putfield arg4 cellC, (* targ.arg4 ← copy(source.arg4,j) *)
Goto endlab4,      (* end label 4 *)
```

```
Load 7,           (* label 5 *)
Load 3,
Getfield arg5 cellC,
Load 1,
Invoke-static heapC copyRTS [PrimT Integer,
PrimT Integer],
Putfield arg5 cellC, (* targ.arg5 ← copy(source.arg5,j) *)
Goto endlab5,      (* end label 5 *)
```

```
Load 7,           (* label 6 *)
Load 3,
Getfield arg6 cellC,
Load 1,
Invoke-static heapC copyRTS [PrimT Integer,
PrimT Integer],
Putfield arg6 cellC, (* targ.arg6 ← copy(source.arg6,j) *)
Goto endlab6,      (* end label 6 *)
```

```
Load 7,           (* label 7 *)
Load 3,
Getfield arg7 cellC,
Load 1,
Invoke-static heapC copyRTS [PrimT Integer,
PrimT Integer],
Putfield arg7 cellC, (* targ.arg1 ← copy(source.arg1,j) *)
Goto endlab7,      (* end label 7 *)
```

```
Load 7,           (* label 8 *)
Load 3,
Getfield arg8 cellC,
Load 1,
Invoke-static heapC copyRTS [PrimT Integer,
PrimT Integer],
```

```

    Putfield arg8 cellC, (* targ.arg8 ← copy(source.arg8,j) *)
    Load 5,
    LitPush (Intg 1),
    BinOp Add,
    Store 5,          (* i ← i + 1 *)
    Goto bucle6,
    Load 2,
    Return           (* return target *)
]

```

— Program RTS with all Class declarations

```
constdefs Prog-RTS :: jvm-prog
```

```
Prog-RTS ≡
```

```

(let
  classCell = (cellC,objectC,cellFdecl,[]);
  consTableMethod = ((makeConsTable,[PrimT Integer]), PrimT Void,(4,2,instMakeConsTable,[]));
  classConsTable = (consTableC,objectC,consTableFdecl,[consTableMethod]);
  initConsDataMethod = ( (Init, []), PrimT Void, (2,1,instInitConsData,[]));
  classConsData = (consDataC,objectC,consDataFdecl,[initConsDataMethod]);
  setMaxSizeM = ((setMaxSize,[PrimT Integer]), PrimT Void, (1,1,instSetMaxSize,[]));
  slideM = ((slide,[PrimT Integer,PrimT Integer]), PrimT Void, (4,4,instSlide,[]));
  classStackS = (stackC, objectC,stackSFdecl,[setMaxSizeM,slideM]);
  makeDirM = ((makeDirectory, [PrimT Integer]), PrimT Void, (4,2,instMakeDirectory,[]));
  reserM = ((reserveCell,[]), PrimT Integer, (4,2,instReserveCell ,[]));
  releaM = ((releaseCell,[PrimT Integer]), PrimT Void, (2,4,instReleaseCell ,[]));
  classDir = (dirCellC, objectC,dirCFdecl,[makeDirM,reserM,releaM]);
  makeHeapM = ((makeHeap,[PrimT Integer]),PrimT Void, (1,1,instMakeHeap,[]));
  decregionM = ((decregion,[]),PrimT Void,(2,0,instDecregion,[]));
  pushRegionM = ((pushRegion,[]),PrimT Void,(4,1,instPushRegion,[(0,23,26,
    Xcpt ArrayIndexOutOfBounds)]));
  popRegionM = ((popRegion,[]),PrimT Void,(3,2,instPopRegion,[]));
  insertCellM = ((insertCell,[PrimT Integer,PrimT Integer]), PrimT Void,(3,3,instInsertCell,[]));
  copyCellM = ((copyCell,[PrimT Integer]),PrimT Integer,(2,4,instCopyCell,[]));
  cloneM = ((clone,[PrimT Integer,PrimT Integer]),PrimT Integer,(2,5,instClone,[]));
  copyCellAuxM = ((copyCellAux,[RefT (ClassT cellC),RefT (ClassT cellC)],
    PrimT Void,(2,2,instCopyCellAux,[]));
  copyM = ((copyRTS,[PrimT Integer,PrimT Integer]),PrimT Integer,(3,8,instCopy,[]));

  classHeap = (heapC, objectC, heapFdecl,[makeHeapM,decregionM,pushRegionM,popRegionM,
    insertCellM,copyCellM,cloneM,copyCellAuxM,copyM])
  in [classCell,classConsTable,classConsData,classStackS,classDir,classHeap])

```

```
end
```

## 20 Useful functions and theorems from the Haskell Library or Prelude

```
theory HaskellLib
imports Main
begin
```

Function `mapAccumL` is a powerful combination of `map` and `foldl`. Functions `unzip3` and `unzip` are respectively the inverse of `zip3` and `zip`.

**consts**

```
mapAccumL :: ('a => 'b => 'a × 'c) => 'a => 'b list => 'a × 'c list
zipWith   :: ('a => 'b => 'c) => 'a list => 'b list => 'c list
unzip3    :: ('a × 'b × 'c) list => 'a list × 'b list × 'c list
unzip     :: ('a × 'b) list => 'a list × 'b list
```

**primrec**

```
mapAccumL f s [] = (s,[])
mapAccumL f s (x#xs) = (let (s',y) = f s x;
                          (s'',ys) = mapAccumL f s' xs
                        in (s'',y#ys))
```

Some lemmas about `mapAccumL`

**lemma** *mapAccumL-non-empty*:

```
[(s'',ys) = mapAccumL f s xs;
 xs = x#xx]
  => (∃ s' y ys'.
      (s',y) = f s x
      ∧ ys = y # ys')
```

**apply** *clarify*

**apply** (*unfold mapAccumL.simps*)

**apply** (*rule-tac x=fst (f s x) in exI*)

**apply** (*rule-tac x=snd (f s x) in exI*)

**apply** (*rule-tac x=snd (mapAccumL f (fst (f s x)) xx) in exI*)

**apply** (*rule conjI*)

**apply** *simp*

**apply** (*case-tac f s x, simp*)

**by** (*case-tac mapAccumL f a xx, simp*)

**lemma** *mapAccumL-non-empty2*:

```
[(s'',ys) = mapAccumL f s xs;
 xs = x#xx]
  => (∃ s' y ys'.
      (s',y) = f s x
      ∧ (s'',ys') = mapAccumL f s' xx
      ∧ ys = y # ys')
```

**apply** *clarify*

**apply** (*unfold mapAccumL.simps*)

**apply** (*rule-tac x=fst (f s x) in exI*)

**apply** (*rule-tac x=snd (f s x) in exI*)

```

apply (rule-tac x=snd (mapAccumL f (fst (f s x)) xx) in exI)
apply (rule conjI)
apply simp
apply (rule conjI)
apply (case-tac f s x) apply (simp)
apply (case-tac mapAccumL f a xx)
apply (simp)
apply (case-tac f s x) apply (simp)
apply (case-tac mapAccumL f a xx)
apply simp
done

```

```

axioms mapAccumL-non-empty3:
  [(s'',ys) = mapAccumL f s xs;
   0 < length xs]
  ==> (∃ s' y ys'.
        (s',y) = f s (xs!0)
        ∧ (s'',ys') = mapAccumL f s' (tl xs))

```

```

axioms mapAccumL-two-elements:
  [(s3,ys) = mapAccumL f s xs;
   xs = x1 # x2 # xx]
  ==> (∃ s1 s2 y1 y2 ys3.
        (s1,y1) = f s x1
        ∧ (s2,y2) = f s1 x2
        ∧ (s3,ys3) = mapAccumL f s2 xx
        ∧ ys = y1 # y2 # ys3)

```

```

axioms mapAccumL-split:
  [(s2,ys) = mapAccumL f s xs;
   xs1 @ xs2 = xs]
  ==> (∃ s1 ys1 ys2 .
        (s1,ys1) = mapAccumL f s xs1
        ∧ (s2,ys2) = mapAccumL f s1 xs2
        ∧ ys = ys1 @ ys2)

```

```

axioms mapAccumL-one-more:
  [(s1,ys) = mapAccumL f s xs;
   (s2,y) = f s1 x]
  ==> (s2,ys@[y]) = mapAccumL f s (xs@[x])

```

Some integer arithmetic lemmas

```

lemma sum-nat:
  [(x1::nat)=x2;(y1::nat)=y2] ==> x1+y1=x2+y2
apply arith
done

```

```

axioms sum-subtract:
  (x::nat)-y+(z-x)=z-y

```

**axioms additions1:**

$$\begin{aligned} & \llbracket i < m; \text{Suc } m + n \leq l \rrbracket \implies \\ & m - i < \text{nat } (\text{int } l - 1) - n + 1 - (\text{nat } (\text{int } l - 1) - \text{Suc } m - n + 1) \end{aligned}$$

**axioms additions2:**

$$\begin{aligned} & \llbracket i < m; \text{Suc } m + n \leq l \rrbracket \implies \\ & \text{nat } (\text{int } l - 1) - m + (m - \text{Suc } i) = \text{nat } (\text{int } l - 1) - \text{Suc } m + (m - i) \end{aligned}$$

**axioms additions3:**

$$\begin{aligned} & \llbracket i < m; \text{Suc } m + n \leq l \rrbracket \implies \\ & \text{nat } (\text{int } l - 1) - \text{Suc } (m + n) + (m - i) = \text{nat } (\text{int } l - 1) - (m + n) + (m - \text{Suc } i) \end{aligned}$$

**axioms additions4:**

$$\begin{aligned} & \llbracket \text{Suc } m + n \leq l \rrbracket \implies \\ & \text{nat } (\text{int } l - 1) - m = \text{Suc } (\text{nat } (\text{int } l - 1) - \text{Suc } m) \end{aligned}$$

**axioms additions5:**

$$\begin{aligned} & \llbracket \text{Suc } m + n \leq l \rrbracket \implies \\ & \text{Suc } (\text{nat } (\text{int } l - 1) - \text{Suc } (m + n)) = \text{nat } (\text{int } l - 1 - \text{int } n - \text{int } m) \end{aligned}$$

**axioms additions6:**

$$\begin{aligned} & \llbracket \text{Suc } m + n \leq l \rrbracket \implies \\ & n + (\text{nat } (\text{int } l - 1) - \text{Suc } (m + n)) < \text{nat } (\text{int } l - 1) \end{aligned}$$

Some lemmas about lists

**lemma list-non-empty:**

$$0 < \text{length } xs \implies (\exists y \ ys . xs = y \# \ ys)$$

**apply auto**

**apply (insert neq-Nil-conv [of xs])**

**by simp**

**axioms drop-nth:**

$$n < \text{length } xs \implies (\exists y \ ys . \text{drop } n \ xs = y \# \ ys \wedge xs!n = y)$$

**axioms drop-nth3:**

$$n < \text{length } xs \implies \text{drop } n \ xs = (xs!n) \# \ \text{drop } (\text{Suc } n) \ xs$$

**axioms drop-take-Suc:**

$$xs = (\text{take } n \ xs) @ (z \# \ zs) \implies \text{drop } (\text{Suc } n) \ xs = zs$$

**axioms drop-nth2:**

$$\begin{aligned} & \llbracket n < \text{length } xs; \text{drop } n \ xs = ys \rrbracket \\ & \implies ys = xs!n \# \ \text{tl } ys \end{aligned}$$

**axioms drop-append2:**

$$\begin{aligned} & \llbracket \text{drop } n \ xs = zs1 \ @ \ ys1 \ @ \ ys2 \ @ \ zs2 \ @ \ rest; \\ & \text{drop } (m - n) \ (zs1 \ @ \ ys1 \ @ \ ys2 \ @ \ zs2) = ys1 \ @ \ rest' \\ & \rrbracket \implies \\ & \text{drop } (m + \text{length } ys1 - n) \ (zs1 \ @ \ ys1 \ @ \ ys2 \ @ \ zs2) = ys2 \ @ \ zs2 \end{aligned}$$

**axioms drop-append3:**



```

[[ drop n xs = xs1 @ rest;
   drop (m-n) xs1 = ys1 @ ys2
]] ==>
drop m xs = ys1 @ ys2 @ rest

```

**lemma** *nth-via-drop-append*:  $drop\ n\ xs = (y\#\!ys)\@zs \implies xs!\!n = y$   
**apply** (*induct xs arbitrary: n, simp*)  
**by**(*simp add:drop-Cons nth-Cons split:nat.splits*)

**lemma** *drop-Suc-append*:  
 $drop\ n\ xs = (y\#\!ys)\@zs \implies drop\ (Suc\ n)\ xs = ys\@zs$   
**apply** (*induct xs arbitrary: n, simp*)  
**apply** (*simp add:drop-Cons*)  
**by** (*simp split:nat.splits*)

**lemma** *nth-via-drop-append-2*:  $drop\ n\ xs = ((y\#\!ys)\@ws\@zs)\@ms \implies xs!\!n = y$   
**apply** (*induct xs arbitrary: n, simp*)  
**by**(*simp add:drop-Cons nth-Cons split:nat.splits*)

**lemma** *drop-Suc-append-2*:  
 $drop\ n\ xs = ((y\#\!ys)\@ws\@zs)\@ms \implies drop\ (Suc\ n)\ xs = ys\@ws\@zs\@ms$   
**apply** (*induct xs arbitrary: n, simp*)  
**apply** (*simp add:drop-Cons*)  
**by** (*simp split:nat.splits*)

**axioms** *drop-append-length*:  
 $drop\ n\ xs = []\@ys\@zs\@ms \implies drop\ (n + length\ ys)\ xs = zs\@ms$

**axioms** *take-length*:  
 $n \leq length\ xs \implies n = length\ (take\ n\ xs)$

**axioms** *take-append2*:  
 $n < length\ xs \implies x\#\!take\ n\ xs = take\ n\ (x\#\!xs)\@[(x\#\!xs)!\!n]$

**axioms** *take-append3*:  
 $Suc\ n \leq length\ xs \implies take\ (Suc\ n)\ xs = take\ n\ xs\@[xs!\!n]$

**axioms** *concat1*:  
 $xs\@y\#\!ys = (xs\@[y])\@ys$

**axioms** *concat2*:  
 $xs1 = xs2 \implies xs1\@ys = xs2\@ys$

**axioms** *upt-length*:  
 $n \leq m \implies length\ [n..\!m] = m - n$

Some lemmas about finite maps

**axioms** *map-of-distinct*:

```

[[ distinct (map fst xys);
   l < length xys;
   (x,y) = xys ! l
]] => map-of xys x = Some y

```

**axioms** *map-of-distinct2*:

```

map-of xys x = Some y
=> (∃ l . l < length xys ∧ (x,y) = xys ! l)

```

**axioms** *map-upds-nth*:

```

i < m - n => (A([n..<m] [↦] xs)) (n+i) = Some (xs ! i)

```

— The unzip3 function of Haskell library

**primrec**

```

unzip3 [] = ([],[],[])
unzip3 (tup#tups) = (let (xs,ys,zs) = unzip3 tups;
                        (x,y,z) = tup
                        in (x#xs,y#ys,z#zs))

```

**axioms** *unzip3-length*:

```

unzip3 xs = (ys1,ys2,ys3) => length ys1 = length ys2

```

**primrec**

```

unzip [] = ([],[])
unzip (tup#tups) = (let (xs,ys) = unzip tups;
                    (x,y) = tup
                    in (x#xs,y#ys))

```

**primrec**

```

zipWith f (x#xs) yy = (case yy of
  [] => []
  | y#ys => f x y # zipWith f xs ys)
zipWith f [] yy = []

```

**axioms** *zipWith-length*:

```

length (zipWith f xs ys) = min (length xs) (length ys)

```

— The Haskell sum type Either

**datatype** ('a,'b) *Either* = Left 'a | Right 'b

— insertion sort for list of strings

**constdefs**

```

leString :: string => string => bool

```

```
leString s1 s2 == True
```

```
consts
```

```
ins :: string => string list => string list
```

```
primrec
```

```
ins s [] = [s]
```

```
ins s (s'#ss) = (if leString s s' then s#s'#ss  
else s'# ins s ss)
```

```
fun sort :: string list => string list
```

```
where
```

```
sort ss = foldr ins ss []
```

```
fun subList :: 'a list => 'a list => bool
```

```
where
```

```
subList xs ys = ( $\exists$  hs ts. ys = hs @ xs @ ts)
```

```
end
```

## 21 Translation from SVM to JVM

```
theory SVM2JVM
```

```
imports ../JVMSAFE/JVMInstructions SVMState RTSCore HaskellLib
```

```
begin
```

```
types pc = nat
```

```
codeMap = PC  $\rightarrow$  pc
```

```
contMap = CodeLabel  $\rightarrow$  nat
```

```
consMap = Constructor  $\rightarrow$  nat
```

### 21.1 Initialisation code

Initialise cell creation, region stack, and machine stack. The corresponding initialisation method of each class is called.

```
constdefs initSizeTable :: SizesTable => bytecode
```

```
initSizeTable st  $\equiv$  (
```

```
case st of (c,r,s) =>
```

```
[LitPush (Intg (int c)),
```

```
Invoke-static dirCellC makeDirectory [PrimT Integer],
```

```
LitPush (Intg (int r)),
```

```
Invoke-static heapC makeHeap [PrimT Integer],
```

```
LitPush (Intg (int s)),
```

```
Invoke-static stackC setmaxSize [PrimT Integer]])
```

Initialise the runtime constructor table.

```

constdefs auxOneArg :: (ArgType × nat) ⇒ bytecode
auxOneArg arT ≡
  (let j      = snd arT;
      argTyp = (case (fst arT) of
                  IntArg ⇒ 1
                | BoolArg ⇒ 2
                | NonRecursive ⇒ 3
                | Recursive ⇒ 4)
  in [Dup,
      LitPush (Intg (int j)),
      LitPush (Intg argTyp),
      ArrStore])

constdefs fillsOneConstructor :: ((nat × nat × ArgType list) × nat)
  ⇒ instr list
fillsOneConstructor tuple ≡
  (let ((tagl,nargs,argTypes),i) = tuple;
      typePairs      = zip argTypes [0..constdefs initConsTable :: ConstructorTableType ⇒ instr list
initConsTable ct ≡
  (let len      = length ct;
      tuples    = map (% (C,(n,tag,atypes)) . (tag,n,atypes)) ct;
      fillCode  = map fillsOneConstructor (zip tuples [0..

```

## 21.2 Generic translation functions

The translation functions are:

- *trSVM2JVM* translates a complete SafeImp program into a JVM program consisting in a single class "PSafe", with a single method called "PSafeMain", containing the in-line JVM code of the translation.

- *trCodeStore* takes the code store part of the SafeImp program and produces the bytecode part of the JVM program, together with a code map mapping SVM program counters into JVM ones, and a continuation map mapping each SafeImp continuation code label into a distinct natural number global to the program. The first part of the bytecode is a function table solving forward references in the code, and a continuation switch translating the continuation natural numbers into jumps to the corresponding continuation code labels.
- *trSeq* translates a single SafeImp sequence of SVM instructions. It receives the next available JVM pc and returns the next not used pc. It also updates the current code and continuation maps.
- *trInstrAux* translates a single SVM instruction into a bytecode sequence. It updates the next available JVM pc and the codemap. The actual bytecode generation is done by *trInstr*.

**consts**

*trInstr* :: [*pc,codeMap,contMap,consMap,pc,SafeInstr*] => *instr list*

Translate Constructor Table: Each constructor is given a distinct number

**constdefs** *trConsTable* :: *ConstructorTableType* => *consMap*

*trConsTable* *ct* ≡ (  
  *let len = length ct*  
  *in map-of (zip (map fst ct) [0..<len])*)

**constdefs**

*trInstrAux* :: [*CodeLabel,contMap,consMap,pc,pc×nat×codeMap,SafeInstr*]  
  => (*pc×nat×codeMap*) × *instr list*

*trInstrAux* *p ctm map com pcc state safeinstr* == (  
  *case state of (pc,i,cdmap) =>*  
  *let cdmap' = cdmap ((p,i) ↦ pc);*  
  *instrs = trInstr pc cdmap' ctm map com pcc safeinstr;*  
  *n = length instrs*  
  *in ((pc+n,i+1,cdmap<sup>^</sup>),instrs)*)

**constdefs**

*trSeq* :: [*contMap,consMap,pc,pc×codeMap,CodeLabel×CodeSequence×FunName*]  
  => (*pc×codeMap*) × *instr list*

*trSeq* *ctmap com pcc state seq* == (  
  *let (pc,cdmap) = state;*  
  *(p,svms,f) = seq;*  
  *((pc',n,cdmap<sup>^</sup>),instss) =*  
  *mapAccumL (trInstrAux p ctm map com pcc) (pc,0,cdmap) svms*  
  *in ((pc',cdmap<sup>^</sup>),concat instss)*)

**axioms** *svms-good*:  
*distinct (map fst (svms::SVMCode))*

**constdefs**

```

trCodeStore :: [CodeLabel,pc,ContinuationMap,consMap,SVMCode]
              ⇒ instr list × codeMap × contMap
trCodeStore inip inipc ctmap com svms == (
  let (fs,ps,contss) = unzip3 ctmap;
      conts          = concat contss;
      nc             = length conts;
      nf             = length fs;
      cdini          = map-of (zip (zip ps (replicate nf 0)) [inipc+1..<inipc+nf+1]);
      ctm            = map-of (zip conts [1..<nc+1]);
      pcc            = inipc + nf + 1;
      ((pc,cdm),iss) = mapAccumL (trSeq ctm com pcc) (pcc + 1,cdini) svms;
      funJumps       = zipWith (% p n. int(the(cdm (p,0)))-int n)
                          ps [inipc+1..<inipc+nf+1];
      funTable       = map Goto funJumps;
      contSwitch     = Tableswitch 1 (int nc) (map
        ((% n. n-int pcc) ∘ int ∘ the ∘ cdm ∘ (% p.(p,0)))
        conts)
  in (Goto (int (the (cdm (inip,0)))-int inipc) # funTable
     @ (contSwitch # concat iss),cdm,ctm))

```

— Names and signature of the translated program single class and method

**constdefs**

```

safeP :: cname
safeP ≡ Cname "SafeP"
safeMain :: mname
safeMain ≡ "PSafeMain"
sigSafeMain :: sig
sigSafeMain ≡ (safeMain,[])

```

It generates code to initialize the runtime system by creating the constructor table, the cells, the heap and the stack. Then, it translates the SVM instructions.

**constdefs**

```

trSVM2JVM :: SafeImpProg ⇒ jvm-prog × codeMap × contMap
              × consMap
trSVM2JVM prog == (
  let ((svms,ctmap),ini,constable,sizeTable) = prog;
      instrConsTable = initConsTable constable;
      lenCT           = length instrConsTable;
      comap           = trConsTable constable;
      instrSizeTable = initSizeTable sizeTable;
      lenST           = length instrSizeTable;
      (instrs,cdm,ctm) = trCodeStore ini (lenCT+lenST) ctmap comap svms;

```

```

method      = (sigSafeMain, PrimT Void,
              (10,10,instrConsTable @ instrSizeTable @ instrs,[]));
classes     = [(safeP,objectC,[],[method])]
in (classes @ Prog-RTS,cdm,ctm,comap)

```

### constdefs

```

extractBytecode :: jvm-prog ⇒ bytecode
extractBytecode P ≡ (fst (snd (snd (snd (the (method' (P, safeP) sigSafe-
Main)))))))

```

## 21.3 Specific translation from each SafeImp instruction to bytecode

Auxiliary functions of trInstr

*trAddr* translates two absolute addresses into a relative one. Used in JVM instructions Goto and Tableswitch.

### constdefs

```

trAddr :: [nat,nat] ⇒ int
trAddr addr1 addr2 == int addr1 - int addr2

```

*nat2Str* translates a nat into a string

### constdefs

```

nat2Str :: nat ⇒ string
nat2Str v == (if v = 1 then "1"
             else if v = 2 then "2"
             else if v = 3 then "3"
             else if v = 4 then "4"
             else if v = 5 then "5"
             else if v = 6 then "6"
             else if v = 7 then "7"
             else "8")

```

PC Increases

### constdefs

```

incCall :: nat
incCall ≡ 1
incPop  :: nat
incPop  ≡ 24
incMatchN :: nat
incMatchN ≡ 7

```

These are fragments of MATCH and MATCHD translations.

### consts

```

endlabel1 :: int
endlabel2 :: int
endlabel3 :: int

```

```

endlabel4 :: int
endlabel5 :: int
endlabel6 :: int
endlabel7 :: int
endlabel8 :: int
label1 :: int
label2 :: int
label3 :: int
label4 :: int
label5 :: int
label6 :: int
label7 :: int
label8 :: int
labelEndLoop :: int
labelLoop :: int

```

### constdefs

```

InstLabel1 :: instr list
InstLabel1 ≡ [Getstatic Sf stackC, (* S[top-i] ← o.arg1 *)
  Getstatic topf stackC,
  Load 5,
  BinOp Subtract,
  Load 6,
  Getfield arg1 cellC,
  ArrStore,
  Goto endlabel1]

nlab1 :: nat
nlab1 ≡ length InstLabel1

```

```

InstLabel2 :: instr list
InstLabel2 ≡ [Getstatic Sf stackC, (* S[top-i] ← o.arg2 *)
  Getstatic topf stackC,
  Load 5,
  BinOp Subtract,
  Load 6,
  Getfield arg2 cellC,
  ArrStore,
  Goto endlabel2]

nlab2 :: nat
nlab2 ≡ length InstLabel2

```

```

InstLabel3 :: instr list
InstLabel3 ≡ [Getstatic Sf stackC, (* S[top-i] ← o.arg3 *)
  Getstatic topf stackC,
  Load 5,
  BinOp Subtract,
  Load 6,
  Getfield arg3 cellC,
  ArrStore,

```



```

        Goto endlabel3]
nlab3 :: nat
nlab3 ≡ length InstLabel3

InstLabel4 :: instr list
InstLabel4 ≡ [Getstatic Sf stackC, (* S[top-i] ← o.arg4 *)
  Getstatic topf stackC,
  Load 5,
  BinOp Subtract,
  Load 6,
  Getfield arg4 cellC,
  ArrStore,
  Goto endlabel4]
nlab4 :: nat
nlab4 ≡ length InstLabel4

InstLabel5 :: instr list
InstLabel5 ≡ [Getstatic Sf stackC, (* S[top-i] ← o.arg5 *)
  Getstatic topf stackC,
  Load 5,
  BinOp Subtract,
  Load 6,
  Getfield arg5 cellC,
  ArrStore,
  Goto endlabel5]
nlab5 :: nat
nlab5 ≡ length InstLabel5

InstLabel6 :: instr list
InstLabel6 ≡ [Getstatic Sf stackC, (* S[top-i] ← o.arg6 *)
  Getstatic topf stackC,
  Load 5,
  BinOp Subtract,
  Load 6,
  Getfield arg6 cellC,
  ArrStore,
  Goto endlabel6]
nlab6 :: nat
nlab6 ≡ length InstLabel6

InstLabel7 :: instr list
InstLabel7 ≡ [Getstatic Sf stackC, (* S[top-i] ← o.arg7 *)
  Getstatic topf stackC,
  Load 5,
  BinOp Subtract,
  Load 6,
  Getfield arg7 cellC,
  ArrStore,
  Goto endlabel7]

```

```

nlab7 :: nat
nlab7 ≡ length InstLabel7

InstLabel8 :: instr list
InstLabel8 ≡ [Getstatic Sf stackC, (* S[top-i] ← o.arg8 *)
  Getstatic topf stackC,
  Load 5,
  BinOp Subtract,
  Load 6,
  Getfield arg8 cellC,
  ArrStore,
  Goto endlabel8]

nlab8 :: nat
nlab8 ≡ length InstLabel8

Match11 :: nat ⇒ instr list
Match11 l ≡ [Getstatic Sf stackC, (* load S!l *)
  Getstatic topf stackC,
  LitPush (Intg (int l)),
  BinOp Subtract,
  ArrLoad,
  Store 1, (* local1 ← b *)
  Getstatic safeDirf dirCellC,
  Load 1,
  ArrLoad,
  Store 6, (* local6 ← o *)
  Load 6,
  Getfield tagGf cellC,
  Store 2, (* local2 ← TagG *)
  Getstatic tablef consTableC,
  Load 2,
  ArrLoad,
  Dup,
  Getfield tagLf consDataC,
  Store 3, (* local3 ← tagL *)
  Getfield nargsf consDataC,
  Store 4, (* local4 ← nargs *)
  LitPush (Intg 0), (* i ← 0 *)
  Store 5, (* local5 ← i *)
  Getstatic topf stackC,
  Load 4,
  BinOp Add,
  Putstatic topf stackC (* top ← top + nargs *)]

nMatch11 :: nat
nMatch11 ≡ length (Match11 0)

Match12 :: instr list
Match12 ≡ [Load 4, (* nargs *)
  Load 5, (* i *)]

```

```

    Ifcmp GreaterEqual labelEndLoop, (* i ≥ nargs ? *)
    Load 5, (* no, load argument i *)
    Tableswitch 0 7 [label1,label2,label3,label4,
                    label5,label6,label7,label8]]
nMatch12 :: nat
nMatch12 ≡ length Match12

Match2 :: instr list
Match2 ≡ [Load 5,
          LitPush (Intg 1),
          BinOp Add,
          Store 5, (* i = i + 1 *)
          Goto labelLoop]
nMatch2 :: nat
nMatch2 ≡ 5

```

PC increases for MATCH and MATCHD

**defs**

```

endlab1-def: endlab1 ≡ int (nlab2 + nlab3 + nlab4 + nlab5 + nlab6 + nlab7
+ nlab8 + 1)
endlab2-def: endlab2 ≡ int (nlab3 + nlab4 + nlab5 + nlab6 + nlab7 + nlab8
+ 1)
endlab3-def: endlab3 ≡ int (nlab4 + nlab5 + nlab6 + nlab7 + nlab8 + 1)
endlab4-def: endlab4 ≡ int (nlab5 + nlab6 + nlab7 + nlab8 + 1)
endlab5-def: endlab5 ≡ int (nlab6 + nlab7 + nlab8 + 1)
endlab6-def: endlab6 ≡ int (nlab7 + nlab8 + 1)
endlab7-def: endlab7 ≡ int (nlab8 + 1)
endlab8-def: endlab8 ≡ 1

```

**defs**

```

label1-def: label1 ≡ 1
label2-def: label2 ≡ int (nlab1 + 1)
label3-def: label3 ≡ int (nlab1 + nlab2 + 1)
label4-def: label4 ≡ int (nlab1 + nlab2 + nlab3 + 1)
label5-def: label5 ≡ int (nlab1 + nlab2 + nlab3 + nlab4 + 1)
label6-def: label6 ≡ int (nlab1 + nlab2 + nlab3 + nlab4 + nlab5 + 1)
label7-def: label7 ≡ int (nlab1 + nlab2 + nlab3 + nlab4 + nlab5 + nlab6 + 1)
label8-def: label8 ≡ int (nlab1 + nlab2 + nlab3 + nlab4 + nlab5 + nlab6 +
nlab7 + 1)

```

**defs**

```

labelEndLoop-def: labelEndLoop ≡ int (nlab1 + nlab2 + nlab3 + nlab4 + nlab5
+ nlab6 + nlab7
+ nlab8 + 3 + nMatch2)

```

**defs**

```

labelLoop-def: labelLoop ≡ -int (nlab1 + nlab2 + nlab3 + nlab4 + nlab5 + nlab6
+ nlab7
+ nlab8 + nMatch2 + nMatch12) + 1

```

**constdefs**

```

incMatch :: nat
incMatch ≡ nMatch11 + nMatch12 + nlab1 + nlab2 + nlab3 + nlab4 + nlab5
+ nlab6 + nlab7
          + nlab8 + nMatch2 + 1

incMatchD :: nat
incMatchD ≡ nMatch11 + nMatch12 + nlab1 + nlab2 + nlab3 + nlab4 + nlab5
+ nlab6 + nlab7
          + nlab8 + nMatch2 + 3

```

These are fragments of BUILDENV and BUILDCLS translations

**consts**

```

pushAux' :: Item ⇒ nat ⇒ instr list

```

**primrec**

```

pushAux' (ItemConst v) i =
  [Getstatic Sf stackC,
   Load 1,                (* load top+n *)
   LitPush (Intg (int i)),
   BinOp Subtract,
   (if (isBool v = True)
    then LitPush (Bool (the-BoolT v))
    else LitPush (Intg (the-IntT v))),
   ArrStore                (* S[top+n-i] ← v *)
  ]

```

```

pushAux' (ItemVar l) i =
  [Getstatic Sf stackC,
   Getstatic topf stackC,
   LitPush (Intg (int l)),
   BinOp Subtract,
   ArrLoad,
   Store 2,                (* v ← S!l *)
   Getstatic Sf stackC,
   Load 1,                (* load top+n *)
   LitPush (Intg (int i)),
   BinOp Subtract,
   Load 2,                (* load v *)
   ArrStore                (* S[top+n-i] ← v *)
  ]

```

```

pushAux' (ItemRegSelf) i =
  [Getstatic Sf stackC,
   Load 1,                (* load top+n *)
   LitPush (Intg (int i)),
   BinOp Subtract,
   Getstatic kf heapC,
   ArrStore                (* S[top+n-i] ← k *)
  ]

```

]

**constdefs**

*pushAux* :: *Item* × *nat* ⇒ *instr list*  
*pushAux pair* == *case pair of (it,i) => pushAux' it i*

*regAux* selects the region where to insert the fresh cell

**consts**

*regAux* :: *Item* ⇒ *instr list*

**primrec**

*regAux* (*ItemRegSelf*) =  
[*Getstatic kf heapC*,  
*Store 3*]  
*regAux* (*ItemVar l*) =  
[*Getstatic Sf stackC*,  
*Getstatic topf stackC*,  
*LitPush (Intg (int l))*,  
*BinOp Subtract*,  
*ArrLoad*,  
*Store 3*]

*fillAux'* fills the cell

**consts**

*fillAux'* :: *Item* ⇒ *nat* ⇒ *instr list*

**primrec**

*fillAux'* (*ItemVar l*) *i* =  
[*Load 2*, (\* *Load object Cell* \*)  
*Getstatic Sf stackC*, (\* *load S!l* \*)  
*Getstatic topf stackC*,  
*LitPush (Intg (int l))*,  
*BinOp Subtract*,  
*ArrLoad*,  
*Putfield (VName ("arg"@ nat2Str i)) cellC*]  
  
*fillAux'* (*ItemConst v*) *i* =  
[*Load 2*, (\* *load object Cell* \*)  
(*if (isBool v = True) then LitPush (Bool (the-BoolT v))*  
*else LitPush (Intg (the-IntT v))*),  
*Putfield (VName ("arg"@ nat2Str i)) cellC*]

**constdefs**

*fillAux* :: *Item* × *nat* ⇒ *instr list*  
*fillAux pair* == *case pair of (it,i) => fillAux' it i*

Translation to bytecode of each SafeImp Instruction

**primrec**

*trInstr pc cdm ctm com pcc DECREGION* =  
[*Invoke-static heapC decregion* []]

```

trInstr pc cdm ctm com pcc (SLIDE m n) =
  [LitPush (Intg (int m)),
   LitPush (Intg (int n)),
   Invoke-static stackC slide [PrimT Integer,
                                PrimT Integer]]

```

```

trInstr pc cdm ctm com pcc (CALL p) = (
  let pc1 = the(cdm(p,0));
      offset = trAddr pc1 (pc + incCall)
  in [Invoke-static heapC pushRegion [],
      Goto offset] )

```

```

trInstr pc cdm ctm com pcc (PRIMOP oper) =
  [Getstatic Sf stackC, (* load S[top - 1] *)
   Getstatic topf stackC,
   LitPush (Intg 1),
   BinOp Subtract,
   Dup2, (* Dup 2 top opstack *)
   ArrLoad,
   Store 1, (* save v2 *)
   Getstatic Sf stackC,
   Getstatic topf stackC,
   ArrLoad,
   Load 1, (* push v2 on top of v1 *)
   BinOp oper, (* compute v1 op v2 *)
   ArrStore, (* store it at S[top - 1] *)
   Getstatic topf stackC, (* top <- top - 1 *)
   LitPush (Intg 1),
   BinOp Subtract,
   Putstatic topf stackC]

```

```

trInstr pc cdm ctm com pcc REUSE =
  [Getstatic Sf stackC,
   Getstatic topf stackC,
   ArrLoad,
   Invoke-static heapC copyCell [PrimT Integer],
   Store 1, (* local1 <- p *)
   Getstatic Sf stackC,
   Getstatic topf stackC,
   Load 1, (* S[top] <- p *)
   ArrStore]

```

```

trInstr pc cdm ctm com pcc COPY =
  [Getstatic Sf stackC,
   Getstatic topf stackC,
   ArrLoad,
   Store 1, (* local1 <- b *)
   Getstatic Sf stackC,

```

```

Getstatic topf stackC,
LitPush (Intg 1),
BinOp Subtract,
ArrLoad,
Store 2,
Load 1,          (* local2 <- j *)
Load 2,
Invoke-static heapC copyRTS [PrimT Integer,
                             PrimT Integer],
Store 3,          (* local3 <- b' *)
Getstatic topf stackC,
LitPush (Intg 1),
BinOp Subtract,
Putstatic topf stackC,  (* top <- top - 1 *)
Getstatic Sf stackC,
Getstatic topf stackC,
Load 3,
ArrStore]

```

```

trInstr pc cdm ctm com pcc POPCONT =
[Getstatic Sf stackC,
Getstatic topf stackC,
Dup2,
Dup2,
Dup2,
ArrLoad,
Store 1,          (* local1 <- b *)
LitPush (Intg 1),
BinOp Subtract,
ArrLoad,
Store 2,          (* local2 <- k' *)
LitPush (Intg 2),
BinOp Subtract,
ArrLoad,
Store 3,          (* local3 <- p *)
LitPush (Intg 2),
BinOp Subtract,
Dup,
Putstatic topf stackC,  (* top <- top - 2 *)
Load 1,
ArrStore,          (* S[top] <- b *)
Load 2,
Putstatic kOf heapC,  (* k0 <- k' *)
Load 3,            (* jump to continuation *)
Goto (trAddr pcc (pc + incPop))]

```

```

trInstr pc cdm ctm com pcc (PUSHCONT p) =
( let n = the(ctm(p))
  in

```

```

[Getstatic topf stackC,
 LitPush (Intg 1),
 BinOp Add,
 Putstatic topf stackC,      (* top ← top + 1 *)
 Getstatic Sf stackC,
 Getstatic topf stackC,
 LitPush (Intg (int n)),
 ArrStore,                   (* S[top] ← p' *)
 Getstatic topf stackC,
 LitPush (Intg 1),
 BinOp Add,
 Putstatic topf stackC,      (* top ← top + 1 *)
 Getstatic Sf stackC,
 Getstatic topf stackC,
 Getstatic k0f heapC,
 ArrStore,                   (* S[top] ← k0 *)
 Getstatic kf heapC,         (* k0 ← k *)
 Putstatic k0f heapC]

```

```

trInstr pc cdm ctm com pcc (MATCHN l v m ps) =
  (let pcs = map (%p.(the(cdm(p,0)))) ps;
   pcs' = map (%n.(trAddr n (pc + incMatchN))) pcs
  in [Getstatic Sf stackC,      (* load S!l *)
     Getstatic topf stackC,
     LitPush (Intg (int l)),
     BinOp Subtract,
     ArrLoad,
     LitPush (Intg (int v)),    (* subtract v *)
     BinOp Subtract,
     Tableswitch 0 (int (m + 1)) pcs'])

```

```

trInstr pc cdm ctm com pcc (MATCH l ps) =
  (let len = length ps;
   pcs = map (%p.(the(cdm(p,0)))) ps;
   pcs' = map (%n.(trAddr n (pc + incMatch))) pcs
  in Match11 l @ Match12 @
     InstLabel1 @ InstLabel2 @ InstLabel3 @ InstLabel4 @
     InstLabel5 @ InstLabel6 @ InstLabel7 @ InstLabel8 @
     Match2 @
     [Load 3,                  (* Load tagL *)
      Tableswitch 0 (int (len - 1)) pcs'])

```

```

trInstr pc cdm ctm com pcc (MATCHD l ps) =
  (let len = length ps;
   pcs = map (%p.(the(cdm(p,0)))) ps;
   pcs' = map (%n.(trAddr n (pc + incMatchD))) pcs
  in Match11 l @ Match12 @
     InstLabel1 @ InstLabel2 @ InstLabel3 @ InstLabel4 @
     InstLabel5 @ InstLabel6 @ InstLabel7 @ InstLabel8 @

```



```

Match2 @
[Load 1, (* release Cell b *)
Invoke-static dirCellC releaseCell [PrimT Integer],
Load 3, (* Load tagL *)
Tableswitch 0 (int (len - 1)) pcs']

trInstr pc cdm ctm com pcc (BUILDENV its) =
(let n = length its;
 genCode = map pushAux (zip its [0..<n])
 in [Getstatic topf stackC,
 LitPush (Intg (int n)),
 BinOp Add,
 Store 1] (* local1 ← top + n *)
 @ concat genCode
 @ [Load 1, Putstatic topf stackC])

trInstr pc cdm ctm com pcc (BUILDCLS c its item) =
(let tagg = the(com c);
 n = length its;
 codRes = [Invoke-static dirCellC reserveCell [],
 Store 1, (* local1 ← p *)
 Getstatic safeDirf dirCellC,
 Load 1,
 ArrLoad,
 Store 2, (* local2 ← o cell *)
 Load 2,
 LitPush (Intg (int tagg)),
 Putfield tagGf cellC]; (* o.tagG = tagg *)
 codFil = map fillAux (zip its [1..<n+1]);
 codIns = [Load 3,
 Load 1,
 Invoke-static heapC insertCell [PrimT Integer,
 PrimT Integer],
 Getstatic topf stackC,
 LitPush (Intg 1),
 BinOp Add,
 Putstatic topf stackC, (* top ← top + 1 *)
 Getstatic Sf stackC,
 Getstatic topf stackC,
 Load 1,
 ArrStore] (* S[top] ← b *)
 in codRes @ regAux item @ concat codFil @ codIns)

```

end

## 22 Semantics of the SVM instructions

```
theory SVMSemantics
imports SVMState HaskellLib
begin
```

- 'execSVMInst' executes a single SafeImp instruction in a state and gives another state.
- 'execSVM' executes the next SafeImp instruction in a SVM program.
- 'execOp' operates two basic values with an operator.

```
consts
  execSVMInst :: [SafeInstr, ConstructorTableFun, Heap, Region, PC, Stack]
               => (SVMState, SVMState) Either
```

```
consts
  item2Stack :: Region => Stack => Item => StackObject
  item2Val   :: Stack => Item => Val
```

Auxiliary functions for BUILDENV and BUILDCLS

```
fun stackOffset :: Stack => nat => StackObject (infix !+ 110 ) where
  (so # S) !+ 0 = so
| (Val v # S) !+ n = S !+ (n - 1)
| (Reg r # S) !+ n = S !+ (n - 1)
| (Cont c # S) !+ n = S !+ (n - 2)
```

```
primrec
  item2Stack k S (ItemConst v) = Val v
  item2Stack k S (ItemVar off) = S !+ off
  item2Stack k S ItemRegSelf   = Reg k
```

```
primrec
  item2Val S (ItemConst v) = v
  item2Val S (ItemVar off) = (case S !+ off of Val v => v)
```

— We use fun here for full pattern matching

```
constdefs
  execSVM :: SafeImpProg => SVMState => (SVMState, SVMState) Either
  execSVM prog state == (
    case prog of ((code, cnt), q, ct, st) =>
```

```

case state of (h,k0,(l,i),S) =>
let codef = map-of code
in execSVMInst (fst(the (codef l))!i) (map-of ct) h k0 (l,i) S

```

**consts**

```

execSVM-N :: SafeImpProg => SVMState => nat => SVMState

```

**primrec**

```

execSVM-N prog s 0 = s
execSVM-N prog s (Suc n) = (case execSVM prog s of
  Right s' => execSVM-N prog s' n)

```

Function `execSVMInst` executes a single SVM instruction. If it is POP-CONT and there is no continuation in the stack, it returns the current state with the constructor `Left`. Otherwise, it returns the next state with the constructor `Right`.

**primrec**

```

execSVMInst DECREGION ct h k0 pc S = Right (h ↓ k0,k0,incrPC pc,S)

execSVMInst POPCONT ct h k0 pc S = (case S of
  Val v#[] => Left (h,k0,pc,S)
| Val v#Cont (k0,l)#S' => Right (h,k0,(l,0),Val v#S'))

execSVMInst (PUSHCONT p)ct h k0 pc S = Right (h,snd h,incrPC pc,Cont
(k0,p)#S)

execSVMInst COPY ct h k0 pc S = (case S of
  Val (Loc b)#Reg j#S' =>
    if j ≤ snd h then
      let pair = copy h j b in
      case pair of (h',b') =>
        Right (h',k0,incrPC pc,Val (Loc b')#S')
    else Left (h,k0,pc,S))

execSVMInst REUSE ct h k0 pc S = (case S of
  Val (Loc b)#S' =>
    case h of (hm,k) =>
      let cell = the (hm b);
          b' = getFresh hm in
      Right ((hm(b:=None)(b'↦ cell),k),
        k0,incrPC pc,Val (Loc b')#S'))

execSVMInst (CALL p) ct h k0 pc S = (case h of
  (hm,k) => Right ((hm,k+1),k0,(p,0),S))

execSVMInst (PRIMOP oper) ct h k0 pc S = (case S of
  Val (IntT i)#Val (IntT i')#S' =>
    let v = execOp oper (IntT i) (IntT i') in

```

*Right (h,k0,incrPC pc, Val v#S')*

*execSVMInst (MATCH l ps) ct h k0 pc S = (case S!+l of Val (Loc b) =>*  
*case (fst h) b of*  
*Some (j,C,vs) =>*  
*case ct C of*  
*Some (-,r,-) =>*  
*Right (h,k0,(ps!r,0),(map Val vs) @ S))*

*execSVMInst (MATCHD l ps) ct h k0 pc S = (case S!+l of Val (Loc b) =>*  
*case (fst h) b of*  
*Some (j,C,vs) =>*  
*case ct C of*  
*Some (-,r,-) =>*  
*let h' = ((fst h)(b:= None), snd h) in*  
*Right (h',k0,(ps!r,0),(map Val vs) @ S))*

*execSVMInst (MATCHN l v m ps) ct h k0 pc S = (case S!+l of*  
*Val (IntT i) =>*  
*let r = (nat i) - v;*  
*p = if r ≥ 0 ∧ r ≤ m - 1*  
*then ps!r*  
*else ps!m*  
*in Right (h,k0,(p,0),S)*  
*| Val (BoolT b) =>*  
*let p = if ¬ b then ps!0 else ps!1*  
*in Right (h,k0,(p,0),S))*

*execSVMInst (BUILDENV is) ct h k0 pc S =*  
*(let bs = map (item2Stack (snd h) S) is*  
*in Right (h,k0,incrPC pc,bs @ S))*

*execSVMInst (BUILDCLS C is i) ct h k0 pc S =*  
*(case item2Stack (snd h) S i of*  
*Reg j =>*  
*case h of*  
*(hm,k) =>*  
*if j ≤ snd h then*  
*let vs = map (item2Val S) is;*  
*b = getFresh hm;*  
*h' = (hm (b ↦ (j,C,vs)),k)*  
*in Right (h',k0,incrPC pc, Val (Loc b)#S)*  
*else Left (h,k0,pc,S))*

*execSVMInst (SLIDE m n) ct h k0 pc S = Right (h,k0,incrPC pc,take m S @*  
*drop (m+n) S)*

We convert function `execSVM` into a reflexive transitive relation. Also, we define an inductive relation `execSVMBalanced` giving us the list of all the states reachable from an initial one so that the stack does not decrease more than a given amount `n`. The list is given in reverse order.

**constdefs**

```

execSVMAll :: [SafeImpProg,SVMState,SVMState] => bool
            (-+ - -svm→ - [61,61,61]60)
P ⊢ s -svm→ s' ≡ (s,s') ∈ {(s,t) . execSVM P s = Right t} ^*

```

**fun**

```
diffStack :: SVMState => SVMState => nat => int
```

**where**

```

diffStack (h',k0',pc',S') (h,k0,pc,S) m = (
  let n = length S;
      n' = length S'
  in int m + (int n' - int n))

```

**fun instrSVM :: SafeImpProg => SVMState => SafeInstr**

**where**

```
instrSVM ((code,cnt),q,ct,st) (h,k0,(l,i),S) = fst(the (map-of code l))!i
```

**inductive**

```

execSVMBalanced :: [SafeImpProg,SVMState,nat list,SVMState list,nat list] =>
bool
            (-+- , - -svm→ - , - [61,61,61,61,61]60)

```

**where**

```

init: P ⊢ s, n#ns -svm→ [s], n#ns
| step: [[ P ⊢ s, n#ns -svm→ s'#ss, m#ms;
  execSVM P s' = Right s'';
  m' = nat (diffStack s'' s' m);
  m' ≥ 0;
  ms' = (if pushcont (instrSVM P s') then 0#m#ms
    else if popcont (instrSVM P s') ∧ ms=m''#ms'' then (Suc m')#ms''
    else m'#ms)]]
⇒
P ⊢ s, n#ns -svm→ s''#s'#ss, ms'

```

**lemma popcont-neq-pushcont:**

```
popcont (instrSVM P s) ⇒ ¬pushcont (instrSVM P s)
```

**by** (case-tac instrSVM P s,simp-all)

**lemma execSVMBalanced-n-step-aux [rule-format]:**

```
P ⊢ s'', n''#ns'' -svm→ sss , n'''#ns''' ⇒ (∀ s''' ss'' ss. sss = s'''#ss'' ⇒
```

```

P ⊢ s', n'#ns' -svm→ s'' # ss', n''#ns'' ⇒
ss = ss'' @ ss' ⇒

```

```

  P ⊢ s', n' # ns' -svm → s''' # ss, n''' # ns'''
apply (rule impI)

apply (erule execSVMBalanced.induct) apply simp
apply (rule allI)+
apply (erule-tac x=s'a in allE)
apply (erule-tac x=ss in allE)
apply (erule-tac x=ss @ ss' in allE)
apply (rule impI)+
apply simp

apply (elim conjE)

apply (rule conjI) apply (rule impI) apply (rule conjI) apply (rule impI)
apply simp

apply (erule conjE) apply (erule popcont-neq-pushcont) apply simp

apply (rule impI)
apply simp
apply (erule-tac t=ss'' in sym) apply simp
apply (rule execSVMBalanced.step) apply simp apply simp apply simp apply
simp apply simp
apply (rule conjI) apply simp apply simp

apply (rule impI) apply (rule conjI) apply (rule impI)
apply simp
apply (erule-tac t=ss'' in sym) apply simp
apply (rule execSVMBalanced.step) apply simp apply simp apply simp apply
simp apply simp

apply (rule impI)
apply simp
apply (erule-tac t=ss'' in sym) apply simp
apply (rule execSVMBalanced.step) apply simp apply simp apply simp apply
simp
apply (split split-if-asm) apply simp apply simp
apply (rule impI) apply (elim conjE)
apply (rule conjI) apply simp apply auto
done

lemma execSVMBalanced-n-step:
  [[P ⊢ s', n' # ns' -svm → s'' # ss', n'' # ns'';
    P ⊢ s'', n'' # ns'' -svm → s''' # ss'', n''' # ns''';
    ss = ss'' @ ss']]
  ⇒ P ⊢ s', n' # ns' -svm → s''' # ss, n''' # ns'''
apply (rule execSVMBalanced-n-step-aux)
apply simp-all

```

done

end

## 23 Certification of the translation from SVM to JVM

```
theory CertifSVM2JVM
imports ../JVMSAFE/JVMExec SVM2JVM SVMSemantics
begin
```

### 23.1 Equivalence relations between SVM and JVM states

```
types injection = Location  $\rightarrow$  loc
```

— Equivalence of values under a cell directory and an injection

```
fun equivV :: Val => val => entries- => injection => bool
where
  equivV (IntT i) v d g = ( $\exists j. v = Intg j \wedge (i = j)$ )
| equivV (BoolT b) v d g = ( $\exists i. v = Intg i \wedge$ 
   $((i \neq 0 \wedge b) \vee (i = 0 \wedge \neg b))$ )
| equivV (Val.Loc l) v d g = ( $\exists i l'. v = Intg i$ 
   $\wedge d (nat i) = Some (Addr l') \wedge g l = Some l'$ )
```

— Cell to the right of a given one in JVM heap

```
constdefs
  nextCell :: aheap  $\Rightarrow$  loc  $\Rightarrow$  loc
  nextCell h l == (case h l of
    Some (Obj cls flds) =>
      case flds (derf, cellC) of
        Some (Addr next) => next)
```

— Cells belonging to a region in the JVM heap, starting from a given one

```
inductive-set
  cellsReg :: aheap  $\Rightarrow$  loc  $\Rightarrow$  loc set
  for h :: aheap and l :: loc
where
  cellReg-basic: l  $\in$  cellsReg h l
| cellReg-step:  $\llbracket l \in \text{cellsReg } h \ l; l' = \text{nextCell } h \ l \rrbracket$ 
   $\Longrightarrow l' \in \text{cellsReg } h \ l$ 
```

— this gives us the set of references belonging to a given region

```
constdefs
  region :: entries- => aheap => nat => loc set
  region rs h j == (
    let header = the-Addr (the (rs j));
```

$allcells = cellsReg\ h\ header$   
in  $allcells - \{header\}$ )

— this gives us the set of references belonging to all regions

**constdefs**

$activeCells :: entries- \Rightarrow aheap \Rightarrow nat \Rightarrow loc\ set$   
 $activeCells\ rs\ h\ k \equiv \{p . \exists j . j \leq k \wedge p \in region\ rs\ h\ j\}$

**constdefs**

$passiveCells :: entries- \Rightarrow nat \Rightarrow loc\ set$   
 $passiveCells\ rs\ k \equiv \{p . \exists j . j \leq k \wedge p = the-Addr\ (the\ (rs\ j))\}$

— This gives us the i-th argument of a cell

**constdefs**

$argCell :: fields- \Rightarrow nat \Rightarrow val$   
 $argCell\ fds\ i == (\$   
   $let\ arg-i = if\ i=1\ then\ arg1\ else$   
   $if\ i=2\ then\ arg2\ else$   
   $if\ i=3\ then\ arg3\ else$   
   $if\ i=4\ then\ arg4\ else$   
   $if\ i=5\ then\ arg5\ else$   
   $if\ i=6\ then\ arg6\ else$   
   $if\ i=7\ then\ arg7\ else\ arg8$   
 $in\ the\ (fds\ (arg-i, cellC)))$

A SVM cell is equivalent to a JVM cell if the JVM object is an instance of the class Cell, if they live in the same region j, the constructor name and the global tag are made equivalent by the constructor map cm, and all the arguments are equivalent.

— Equivalence of cells under a constructor map, a region stack, a cell directory

— and an injection

**fun**

$equivC :: [Region \times Cell, aheap, loc, heap-entry, nat, consMap, entries-, entries-, injection] \Rightarrow bool$

**where**

$equivC\ cell\ h\ ref\ cell'\ k\ cm\ regStack\ d\ g = (\$   
   $\exists j\ C\ vs\ Obj\ cname\ fds\ tag\ n\ reg-j\ tag' .$   
   $cell = (j, C, vs) \wedge$   
   $j \leq k \wedge$   
   $cell' = Obj\ cname\ fds \wedge$   
   $cm\ C = Some\ tag \wedge$   
   $n = length\ vs \wedge$   
   $reg-j = region\ regStack\ h\ j \wedge$   
   $tag' = nat\ (the-Intg\ (the\ (fds\ (tagGf, cname)))) \wedge$   
   $cname = cellC \wedge$   
   $ref \in reg-j \wedge$   
   $tag = tag' \wedge$   
   $(\forall i \in \{k . 0 \leq k \ \& \ k < n\} . equivV\ (vs!i)\ (argCell\ fds\ i)\ d\ g))$



- Equivalence between SVM and JVM heaps under a constructor map,
- a region stack, a cell directory and an injection

**fun**

*equivH* :: [*Heap*,*ahheap*,*nat*,*consMap*,*entries-*,*entries-*,*injection*] => *bool*

**where**

*equivH* (*H*,*k1*) *h k2 cm regStack d g* = (  
*k1=k2* ∧  
*ran g* = *activeCells regStack h k2* ∧  
*finite (dom H)* ∧  
(∀ *l* ∈ *dom H* . ∃ *l'*. *l' = the (g l)*)  
∧ *equivC (the (H l)) h l' (the (h l')) k2 cm regStack d g*)

- Equivalence between SVM and JVM stacks under a continuation map, a cell
- directory and an injection

**fun**

*equivS* :: [*Stack*,*entries-*,*int*,*contMap*,*entries-*,*injection*] => *bool*

**where**

*equivS* [] *S' n ctm d g* = (*n = (-1)*)  
| *equivS (Cont (k,p)#S) S' n ctm d g* = (*n >= 1*  
∧ *k = nat (the-Intg (the (S' (nat n))))*  
∧ *the (ctm p) = nat (the-Intg (the (S' ((nat n - 1))))*)  
∧ *equivS S S' (n - 2) ctm d g*)  
| *equivS (Reg j#S) S' n ctm d g* = (*n >= 0*  
∧ *j = nat (the-Intg (the (S' (nat n))))*  
∧ *equivS S S' (n - 1) ctm d g*)  
| *equivS (Val v#S) S' n ctm d g* = (*n >= 0*  
∧ *equivV v (the (S' (nat n))) d g*  
∧ *equivS S S' (n - 1) ctm d g*)

- Equivalence between SVM and JVM states

**constdefs**

*equivState* :: [*codeMap*,*contMap*,*consMap*,*SVMState*,*jvm-state*] => *bool*

(- , -, - ⊢ -  $\triangleq$  - [*71*,*71*,*71*,*71*,*71*] *70*)

*equivState cdm ctm com st1 st2* == (

∃ *H k k0 PC S sh h inih vs pc ref k' k0' l ty m S' n l' regS l''*  
*m' m'' d g* .

*st1* = ((*H,k,k0,PC,S*) ∧  
*st2* = (*None,sh,h,inh,([],vs, safeP, sigSafeMain,pc,ref)#[]*) ∧  
*k'* = *nat (the-Intg (the (sh (heapC,kf))))* ∧  
*k0'* = *nat (the-Intg (the (sh (heapC,k0f))))* ∧  
*sh (stackC,Sf) = Some (Addr l)* ∧  
*distinct [l,l',l'']* ∧  
*activeCells regS h k' ∩ {l,l',l''} = {}* ∧  
*h l = Some (Arr ty m S')* ∧  
*sh (stackC,topf) = Some (Intg n)* ∧

$$\begin{aligned}
& sh \text{ (heapC,regionsf)} = \text{Some (Addr } l') \wedge \\
& h \text{ } l' = \text{Some (Arr ty } m' \text{ regS)} \wedge \\
& sh \text{ (dirCellC,safeDirf)} = \text{Some (Addr } l'') \wedge \\
& h \text{ } l'' = \text{Some (Arr ty } m'' \text{ d)} \wedge \\
& dom \text{ } g = dom \text{ } H \wedge \\
& inj\text{-on } g \text{ (dom } H) \wedge \\
& equivH \text{ (H,k)} \text{ } h \text{ } k' \text{ com regS } d \text{ } g \wedge \\
& n < int \text{ } m \wedge \\
& equivS \text{ } S \text{ } S' \text{ } n \text{ ctm } d \text{ } g \wedge \\
& k0 = k0' \wedge \\
& pc = the \text{ (cdm PC)}
\end{aligned}$$

**axioms** *PC-good*:

$$\begin{aligned}
& \llbracket (P', cdm, ctm, com) = trSVM2JVM ((svms, ctmap), ini, ct, st); \\
& \quad equivState \text{ cdm ctm com ((H,k),k0,(l,i),S) st2} \\
& \rrbracket \implies l : dom \text{ (map-of svms)} \\
& \quad \wedge i < length \text{ (fst (the (map-of svms l)))}
\end{aligned}$$

**axioms** *RightNotUndefined* :  $\forall x . \neg (undefined = Right \ x)$

**axioms** *execSVMInstr-COPY* :

$$\begin{aligned}
& \llbracket (P', cdm, ctm, com) = trSVM2JVM ((svms, ctmap), ini, ct, ah, ai, bc); \\
& \quad cdm , ctm, com \vdash ((hm, k), k0, (l, i), S) \hat{=} S1'; \\
& \quad (fst \text{ (the (map-of svms l)) } ! i) = COPY; \\
& \quad execSVMInst \text{ COPY (map-of ct) (hm, k) } k0 \text{ (l, i) } S = \text{Either.Right } S2; \\
& \quad drop \text{ (the (cdm (l, i))) (extractBytecode } P') = \\
& \quad trInstr \text{ (the (cdm (l, i))) } cdm' \text{ ctm}' \text{ com } pcc \text{ COPY @ bytecode}' \\
& \rrbracket \implies \exists v' sh' dh' ih' fms' . \\
& \quad P' \vdash S1' \text{ -jvm} \rightarrow (v', sh', dh', ih', fms') \wedge \\
& \quad cdm , ctm, com \vdash S2 \hat{=} (v', sh', dh', ih', fms')
\end{aligned}$$

**axioms** *execSVMInstr-MATCHD* :

$$\begin{aligned}
& \llbracket (P', cdm, ctm, com) = trSVM2JVM ((svms, ctmap), ini, ct, ah, ai, bc); \\
& \quad cdm , ctm, com \vdash ((hm, k), k0, (l, i), S) \hat{=} S1'; \\
& \quad (fst \text{ (the (map-of svms l)) } ! i) = MATCHD \text{ off ps}; \\
& \quad execSVMInst \text{ (MATCHD off ps) (map-of ct) (hm, k) } k0 \text{ (l, i) } S = \\
& \quad \quad \text{Either.Right } S2; \\
& \quad drop \text{ (the (cdm (l, i))) (extractBytecode } P') = \\
& \quad trInstr \text{ (the (cdm (l, i))) } cdm' \text{ ctm}' \text{ com } pcc \text{ (MATCHD off ps) @ bytecode}' \\
& \rrbracket \implies \exists v' sh' dh' ih' fms' . \\
& \quad P' \vdash S1' \text{ -jvm} \rightarrow (v', sh', dh', ih', fms') \wedge \\
& \quad cdm , ctm, com \vdash S2 \hat{=} (v', sh', dh', ih', fms')
\end{aligned}$$

**axioms** *equivS-ge-n*:

*equivS S S' n ctm d g*  $\implies$   
 $(\forall p > 0. \text{equivS } S (S'(\text{nat } (p + n) \mapsto A)) n \text{ ctm } d g)$

**axioms** *equivS-length*:

*equivS S S' n ctm d g*  $\implies n = \text{int } (\text{length } S) - 1$

**axioms** *length-vs*:

*length vs = 10*

**lemma** *l-not-in-cellReg*:

$\llbracket l \notin \text{activeCells } \text{regS } h k; l \notin \text{passiveCells } \text{regS } k \rrbracket$   
 $\implies \forall j \leq k. l \notin \text{cellsReg } h (\text{the-Addr } (\text{the } (\text{regS } j)))$

**apply** (*simp add: activeCells-def*)

**apply** (*simp add: passiveCells-def*)

**apply** (*simp add: region-def*)

**done**

**axioms** *safeDir-bounds*:

*the-Intg (the (S' i)) < int m  $\wedge$  0  $\leq$  the-Intg (the (S' i))*

**axioms** *l-not-in-cellReg*:

$l \notin \text{activeCells } \text{regS } h k$   
 $\implies \forall j . l \notin \text{cellsReg } h (\text{the-Addr } (\text{the } (\text{regS } j)))$

**axioms** *activeCells*:

$l \notin \text{activeCells } \text{regS } h k$   
 $\implies l \notin \text{activeCells } \text{regS } (h(l \mapsto A)) k$

**axioms** *activeCells-2*:

$\llbracket l' \notin \text{activeCells } \text{regS } h k; l \neq l' \rrbracket$   
 $\implies l' \notin \text{activeCells } \text{regS } (h(l \mapsto A)) k$

**axioms** *method-safeMain*:

$(P', \text{cdm}, \text{ctm}, \text{com}) = \text{trSVM2JVM } ((\text{svms}, \text{ctmap}), \text{ini}, \text{ct}, \text{ah}, \text{ai}, \text{bc})$   
 $\implies (\text{safeP}, \text{PrimT } \text{Void}, 10, 10, (\text{instrConsTable } @ \text{instrSizeTable } @ \text{instrs}, [])) =$   
 $(\text{the } (\text{method}' (P', \text{safeP}) (\text{safeMain}, [])))$

**declare** *equivH.simps* [*simp del*]

**lemma** *cellsReg-monotone-1* [*rule-format*]:

$x \in \text{cellsReg } h l$   
 $\longrightarrow x \neq l$   
 $\longrightarrow l \neq l'$   
 $\longrightarrow x = \text{nextCell } (h(l' \mapsto A)) l$

```

apply (rule impI)
apply (erule cellsReg.induct,simp)
by (simp add: nextCell-def)

```

```

lemma cellsReg-monotone-2 [rule-format]:

```

```

   $x \in \text{cellsReg } h(l' \mapsto A) \ l$ 
   $\longrightarrow x \neq l$ 
   $\longrightarrow l \neq l'$ 
   $\longrightarrow x = \text{nextCell } h \ l$ 

```

```

apply (rule impI)
apply (erule cellsReg.induct,simp)
by (simp add: nextCell-def)

```

```

lemma l-not-in-regs:

```

```

   $l \notin \text{cellsReg } h \ (\text{the-Addr } (\text{the } (\text{regS } j)))$ 
   $\implies \text{the-Addr } (\text{the } (\text{regS } j)) \neq l$ 
apply (subgoal-tac the-Addr (the (regS j))  $\in$  cellsReg h (the-Addr (the (regS j))))

```

```

apply blast
by (rule cellReg-basic)

```

```

end

```

## 24 Lemmas on the static properties of the translation SVM to JVM

```

theory dem-translation

```

```

imports ../JVMSAFE/JVMExec SVM2JVM SVMSemantics CertifSVM2JVM
begin

```

```

declare trInstr.simps [simp del]
declare equivS.simps [simp del]
declare drop-append [simp del]

```

```

lemma fun-trInstr1-aux [rule-format]:

```

```

   $(\forall pc \ i \ cdm \ pc' \ i' \ cdm' \ insts \ .$ 
   $(\forall j \geq i \ . (p,j) \notin \text{dom } cdm) \longrightarrow$ 
   $((pc',i',cdm'),insts) = \text{mapAccumL } (\text{trInstrAux } p \ \text{ctm } \ \text{com } \ \text{pcc}) \ (pc,i,cdm) \ \text{svms}$ 
   $\longrightarrow cdm \subseteq_m cdm' \wedge (\forall j \geq i'. (p,j) \notin \text{dom } cdm')$ 

```

```

apply (induct svms)

```

```

apply simp

```

```

apply (rule allI)+
apply (rename-tac inst svms pc1 i1 cdm1 pc2 i2 cdm2 insts2)
apply (unfold mapAccumL.simps)
apply (case-tac trInstrAux p ctm com pcc (pc1, i1, cdm1) inst)
apply (rename-tac state bc)
apply (case-tac state, rename-tac pc3 rest)
apply (case-tac rest, rename-tac i3 cdm3, simp)
apply (case-tac mapAccumL (trInstrAux p ctm com pcc) (pc3, i3, cdm3) svms)
apply (rename-tac state4 insts4)
apply (case-tac state4, rename-tac pc4 rest4)
apply (case-tac rest4, rename-tac i4 cdm4, clarsimp)
apply (subgoal-tac ((pc4, i4, cdm4), insts4)
  = mapAccumL (trInstrAux p ctm com pcc) (pc3, i3, cdm3) svms)
prefer 2 apply simp

```

```

apply (erule-tac x=pc3 in allE)
apply (rotate-tac -1)
apply (erule-tac x=i3 in allE)
apply (rotate-tac -1)
apply (erule-tac x=cdm3 in allE)

```

```

apply (erule impE)
prefer 2
apply (rotate-tac -1)
apply (erule-tac x=pc4 in allE)
apply (rotate-tac -1)
apply (erule-tac x=i4 in allE)
apply (rotate-tac -1)
apply (erule-tac x=cdm4 in allE)
apply (erule impE)
apply (erule-tac x=insts4 in exI, assumption)

```

```

apply (elim conjE)
apply (rule conjI)
apply (unfold trInstrAux-def) apply (simp)
apply (elim conjE)
apply (subgoal-tac cdm1  $\subseteq_m$  cdm1 ((p, i1)  $\mapsto$  pc1))
apply (erule map-le-trans, simp)

```

```

apply (erule-tac x=i1 in allE)
apply (erule impE, simp)
apply (unfold map-le-def) apply (force)
apply assumption
apply simp
apply (elim conjE)

```

**apply** *clarify*  
**apply** (*erule-tac*  $x=j$  **in** *allE*)  
**apply** (*drule* *mp*)  
**by** *auto*

**lemma** *fun-trInstr1* [*rule-format*]:  
 $(\forall pc\ i\ cdm\ pc'\ i'\ cdm'\ insts .$   
 $(\forall j \geq i . (p,j) \notin \text{dom } cdm) \longrightarrow$   
 $((pc',i',cdm'),insts) = \text{mapAccumL } (trInstrAux\ p\ ctm\ \text{com } pcc) (pc,i,cdm)\ svms$   
 $\longrightarrow cdm \subseteq_m cdm')$   
**apply** (*rule allI*)+  
**apply** (*rule impI*)+  
**thm** *fun-trInstr1-aux*  
**apply** (*subgoal-tac*  $cdm \subseteq_m cdm' \wedge (\forall j \geq i'. (p,j) \notin \text{dom } cdm')$ )  
**prefer** 2 **apply** (*rule fun-trInstr1-aux*)  
**apply** (*erule-tac*  $x=j$  **in** *allE*) **apply** (*drule* *mp*)  
**by** *auto*

**lemma** *fun-trInstrAux*:  
 $((pc',cdm'),bc') = trSeq\ ctm\ \text{com } pcc\ (pc,cdm)\ (p,svms,fn) \longrightarrow$   
 $svms = is \# iss \longrightarrow$   
 $(\forall j \geq 0 . (p,j) \notin \text{dom } cdm) \longrightarrow$   
 $drop\ pc\ bc = bc' @ rest \longrightarrow$   
 $(\forall i . i < \text{length } svms \longrightarrow$   
 $(\exists pc''\ cdm''\ bc''\ rest'\ insts1.$   
 $((pc''+\text{length } bc'',i+1,cdm''((p,i) \mapsto pc'')),bc'') =$   
 $trInstrAux\ p\ ctm\ \text{com } pcc\ (pc'',i,cdm'')\ (svms!i)$   
 $\wedge drop\ (pc''-pc)\ bc' = bc'' @ rest'$   
 $\wedge ((pc'',i,cdm''),insts1) =$   
 $mapAccumL\ (trInstrAux\ p\ ctm\ \text{com } pcc)\ (pc,0,cdm)\ (take\ i\ svms)$   
 $\wedge cdm''((p,i) \mapsto pc'') \subseteq_m cdm')$

**apply** (*rule impI*)+  
**apply** (*rule allI*)

**apply** (*induct-tac* *i*)

**apply** (*rule impI*)  
**apply** (*unfold trSeq-def*)  
**apply** *clarify*  
**apply** (*case-tac*  $(pc,cdm)$ )  
**apply** (*case-tac*  $(p, is \# iss, fn)$ )  
**apply** (*rename-tac*  $pc2\ cdm2\ p2\ rest2$ )  
**apply** (*case-tac*  $rest2$ )  
**apply** (*rename-tac*  $instr\ ff$ )  
**apply** (*simp del: mapAccumL.simps*)  
**apply** (*case-tac*  $mapAccumL\ (trInstrAux\ p2\ ctm\ \text{com } pcc)\ (pc2, 0, cdm2)\ instr$ )

```

apply (rename-tac result instss2)
apply (case-tac result)
apply (rename-tac pc3 rest2)
apply (case-tac rest2)
apply (simp del: mapAccumL.simps)
apply clarify

```

```

apply (rename-tac nn bca)
apply (subgoal-tac ( $\exists$  s' y ys'.
  (s',y) = trInstrAux p ctm com pcc (pc, 0, cdm) is
   $\wedge$  ((pc',nn,cdm'),ys')=mapAccumL (trInstrAux p ctm com pcc) s' iss
   $\wedge$  instss2 = y # ys'))
prefer 2
apply (subgoal-tac ((pc', nn, cdm'), instss2)=
  mapAccumL (trInstrAux p ctm com pcc) (pc, 0, cdm) (is # iss))
prefer 2 apply simp
apply (erule mapAccumL-non-empty2,simp)
apply (elim exE,elim conjE)
apply clarify
apply (rename-tac pc4 nn4 cdm4 bc0 ys')

```

```

apply (rule-tac x=pc in exI)
apply (rule-tac x=cdm in exI)
apply (rule-tac x=bc0 in exI)

```

```

apply (subgoal-tac (pc4, nn4, cdm4)=
  (pc + length bc0, Suc 0,  $\lambda$  a. if a = (p, 0) then Some pc else cdm a))
apply (rule conjI)
apply simp
apply (rule conjI)
apply (rule-tac x=concat ys' in exI)
apply force
apply (rule conjI)
apply (rule-tac x=[] in exI,simp)
apply (simp del: mapAccumL.simps add: trInstrAux-def)
apply (fold fun-upd-apply)
apply (elim conjE)
apply clarify
thm fun-trInstr1
apply (rule-tac i=Suc 0 and p=p in fun-trInstr1)
apply (erule-tac x=j in allE)
apply force
apply simp
apply (simp del: mapAccumL.simps add: trInstrAux-def)
apply (fold fun-upd-apply)
apply force

```

```

apply clarify
apply (case-tac (pc,cdm))
apply (case-tac (p, is # iss, fn))
apply (rename-tac pc2 cdm2 p2 rest2)
apply (case-tac rest2)
apply (rename-tac instr ff)
apply (simp del: mapAccumL.simps)
apply (case-tac mapAccumL (trInstrAux p2 ctm com pcc) (pc2, 0, cdm2) instr)
apply (rename-tac result instss)
apply (case-tac result)
apply (rename-tac pc3 rest3)
apply (case-tac rest3)
apply (rename-tac n3 cdm3)
apply (simp del: mapAccumL.simps)
apply clarify

```

```

apply (subgoal-tac (take n (is#iss)) @ (drop n (is#iss))=is#iss)
prefer 2
apply (rule append-take-drop-id)
apply (subgoal-tac  $\exists s1\ insts1\ insts2 . (s1,insts1) =$ 
  mapAccumL (trInstrAux p ctm com pcc) (pc, 0, cdm) (take n (is # iss))
   $\wedge ((pc', n3, cdm'), insts2)=$ 
  mapAccumL (trInstrAux p ctm com pcc) s1 (drop n (is # iss))
   $\wedge instss = insts1 @ insts2$ )
prefer 2
apply (rule mapAccumL-split [where xs=is # iss]) apply(simp,simp)
apply (elim exE, elim conjE)
apply (case-tac s1)
apply (rename-tac pc1 rest1)
apply (case-tac rest1)
apply (rename-tac nn1 cdm1)
apply clarify
apply (subgoal-tac ((pc'', n, cdm''), insts1)=((pc1, nn1, cdm1), insts1a))
prefer 2
apply simp
apply clarify
apply (thin-tac ?x=mapAccumL ?f (pc,0,cdm) ?xs)

```

```

apply (subgoal-tac nn1 < length (is#iss))
apply (subgoal-tac ( $\exists is2\ iss2 . drop\ nn1\ (is\ \#\ iss) = is2\ \#\ iss2$ 
   $\wedge (is\ \#\ iss) ! nn1 = is2$ ))
prefer 2
apply (erule drop-nth)
prefer 2
apply simp

```



```

apply (elim exE)
apply clarify

apply (subgoal-tac ( $\exists s' y ys'$ .
  ( $s',y = trInstrAux p ctm com pcc (pc1, nn1, cdm1) ((is \# iss) ! nn1)$ 
   $\wedge ((pc', n3, cdm'),ys') = mapAccumL (trInstrAux p ctm com pcc) s' iss2$ 
   $\wedge insts2 = y \# ys')$ )
  prefer 2
  apply (erule mapAccumL-non-empty2,simp)
apply (elim exE)
apply (elim conjE)
apply (simp del: mapAccumL.simps drop-append)
apply (rename-tac s' bc2 ys')
apply (case-tac s')
apply (rename-tac pc2 rest2)
apply (case-tac rest2, rename-tac nn2 cdm2,clarify)
apply (subgoal-tac ( $(pc1 + length bc'', Suc nn1,$ 
   $\% a. if a = (p, nn1) then Some pc1 else cdm1 a), bc''$ 
   $=((pc2, nn2, cdm2), bc2)$ )
  prefer 2
  apply simp
apply (clarify, thin-tac ?x=trInstrAux p ctm com pcc ?y ?z)

apply (subgoal-tac  $\exists iss3. iss2=(iss ! nn1)\#iss3$ )
apply (elim exE,clarify)
apply (subgoal-tac ( $\exists s1 s2 y1 y2 ys3 .$ 
  ( $s1,y1=trInstrAux p ctm com pcc (pc1, nn1, cdm1) ((is \# iss) ! nn1)$ 
   $\wedge (s2,y2)=trInstrAux p ctm com pcc s1 (iss!nn1)$ 
   $\wedge ((pc', n3, cdm'), ys3)=mapAccumL (trInstrAux p ctm com pcc) s2 iss3$ 
   $\wedge bc2\#ys'=y1\#y2\#ys3$ ))
  prefer 2
  apply (erule mapAccumL-two-elements)
  apply simp
apply (elim exE)
apply (case-tac s1, rename-tac pc3 rest3,clarify)
apply (rename-tac nn3 cdm3)
apply (subgoal-tac ( $(pc3, nn3, cdm3), y1)=$ 
  ( $pc1 + length y1, Suc nn1,$ 
   $\% a. if a = (p, nn1) then Some pc1 else cdm1 a), y1$ ))
  prefer 2
  apply simp
apply (clarify,thin-tac ?x=trInstrAux p ctm com pcc ?y ?z)
apply (rename-tac pc4 nn4 cdm4 y1 y2 ys3 pc3 nn3 cdm3)

prefer 2
apply (rule-tac x=tl iss2 in exI)
apply (subgoal-tac drop (Suc nn1) (is\#iss) = iss2)

```

```

prefer 2
apply (subgoal-tac drop (Suc nn1) (is # iss)=drop nn1 (tl (is#iss)))
prefer 2
apply (rule drop-Suc)
apply (subgoal-tac is # iss=take nn1 (is # iss) @ (is # iss) ! nn1 # iss2)
prefer 2
apply simp
apply (erule drop-take-Suc)
apply (rule drop-nth2,assumption,simp)

```

```

apply (rule-tac x=pc1 + length y1 in exI)
apply (rule-tac x= % a. if a = (p, nn1) then Some pc1 else cdm1 a in exI)
apply (rule-tac x=y2 in exI)
apply (rule conjI)

```

```

apply (thin-tac mapAccumL ?x ?y ?z = ?w)
apply (thin-tac ?w=mapAccumL ?x ?y ?z)
apply (thin-tac ?w=mapAccumL ?x ?y ?z)
apply (thin-tac ?x=trInstrAux ?x1 ?x2 ?x3 ?x4 ?x5 ?x6)
apply (simp add: trInstrAux-def)
apply clarify
apply (fold fun-upd-apply)
apply force

```

```

apply (rule conjI)
apply (rule-tac x=concat ys3 in exI)
apply (simp only: concat.simps)
apply (rule drop-append2,simp,simp)

```

```

apply (rule conjI)
apply (rule-tac x=insts1a @ [y1] in exI)
apply (subgoal-tac is # take nn1 iss=(take nn1 (is # iss)) @ [(is # iss) ! nn1])
apply (erule-tac t=is # take nn1 iss in ssubst)
prefer 2
apply (erule take-append2)
apply (rule mapAccumL-one-more,simp,assumption)

```

```

apply (thin-tac drop ?n ?x = ?y)
apply (thin-tac drop ?n ?x = ?y)
apply (thin-tac drop ?n ?x = ?y)
apply (thin-tac (take ?n ?x)@?z = ?y)
apply (thin-tac ?x < ?y)
apply (subgoal-tac cdm4 = cdm1((p, nn1) ↦ pc1,
                               (p, Suc nn1) ↦ pc1 + length y1))

```

```

prefer 2 apply (unfold trInstrAux-def)
apply force
apply (simp del:mapAccumL.simps trInstr.simps)
apply (fold fun-upd-apply)
apply (elim conjE)
apply (thin-tac mapAccumL ?f ?s ?x = ?y)

apply (subgoal-tac cdm1 ((p, nn1) ↦ pc1, (p, Suc nn1) ↦ pc1 + length
  (trInstr pc1 (cdm1((p, nn1) ↦ pc1)) ctm com pcc ((is # iss) ! nn1)))
  ⊆m cdm' ∧ (∀ j ≥ n3.(p,j) ∉ dom cdm'))
prefer 2
apply (rule fun-trInstr1-aux)
prefer 2 apply simp
prefer 2 apply (elim conjE,assumption)

apply (subgoal-tac cdm ⊆m cdm1 ∧ (∀ j ≥ nn1.(p,j) ∉ dom cdm1))
prefer 2
apply (rule fun-trInstr1-aux)
prefer 2 apply simp
apply (erule-tac x=ja in allE,assumption)
apply (thin-tac ?y = mapAccumL ?f ?s ?x)
apply (thin-tac ?y = mapAccumL ?f ?s ?x)
apply (thin-tac ?y = mapAccumL ?f ?s ?x)
apply (thin-tac ?mapAccumL ?f ?s ?x = ?y)
apply (elim conjE)
apply clarify
apply (unfold fun-upd-apply)
apply (rotate-tac 1)
apply (erule-tac x=j in allE)
apply (drule mp,simp)
apply force
done

```

**lemma fun-trInstr :**

```

[[ ((pc',cdm'),bc') = trSeq ctm com pcc (pc,cdm) (p,svms,fn);
  svms = is1 # iss;
  (∀ j ≥ 0 . (p,j) ∉ dom cdm);
  drop pc bc = bc' @ rest;
  svms = svms ! i;
  i < length svms
]] ⇒ (∃ pc'' cdm'' bc'' rest .
  bc'' = trInstr pc'' cdm'' ctm com pcc svms
  ∧ drop pc'' bc = bc'' @ rest
  ∧ cdm'' (p,i) = Some pc''
  ∧ cdm'' ⊆m cdm')

```

```

apply (insert fun-trInstrAux [of pc' cdm' bc' ctm com pcc pc cdm
  p svms fn is1 iss bc rest])

```

**apply** (*drule mp,assumption*)+  
**apply** (*rotate-tac -1*)  
**apply** (*erule-tac x=i in allE*)  
**apply** (*drule mp,assumption*)  
**apply** (*elim exE, elim conjE*)

**apply** (*rule-tac x=pc'' in exI*)  
**apply** (*rule-tac x=cdm''((p, i) ↦ pc'') in exI*)  
**apply** (*rule-tac x=bc'' in exI*)  
**apply** (*rule-tac x=rest' @ rest in exI*)

**apply** (*rule conjI*)  
**apply** (*unfold trInstrAux-def*)  
**apply** *simp*

**apply** (*rule conjI*)  
**apply** *clarsimp*  
**apply** (*erule drop-append3,assumption*)  
**apply** (*rule conjI, simp*)

**by** *assumption*

**axioms** *fun-trSeq1*:  
 $(\exists pc2\ cdm2\ bc2 . ((pc2, cdm2), bc2) = trSeq\ ctm\ com\ pcc\ (pc1, cdm1)\ svms$   
 $\wedge\ pc2 - pc1 = length\ bc2)$

**axioms** *fun-trSeq2*:  
 $((pc2, cdm2), bc2) = trSeq\ ctm\ com\ pcc\ (pc1, cdm1)\ seq$   
 $\implies\ cdm1 \subseteq_m\ cdm2$

**lemma** *fun-trSeq3* [*rule-format*]:  
 $(\forall\ pc1\ cdm1\ pc2\ cdm2\ bcs .$   
 $((pc2, cdm2), bcs) = mapAccumL\ (trSeq\ ctm\ com\ pcc)\ (pc1, cdm1)\ svms \implies$   
 $cdm1 \subseteq_m\ cdm2)$

**apply** (*induct svms*)

**apply** *simp*

**apply** (*rule allI*)+  
**apply** (*rename-tac seq svms pc1 cdm1 pc2 cdm2 bcs*)  
**apply** (*unfold mapAccumL.simps*)  
**apply** (*case-tac trSeq ctm com pcc (pc1, cdm1) seq*)  
**apply** (*rename-tac state bc*)  
**apply** (*case-tac state, rename-tac pc3 cdm3, simp*)

**apply** (*case-tac mapAccumL (trSeq ctm com pcc) (pc3, cdm3) svms, rename-tac state4 bcs4*)  
**apply** (*case-tac state4, rename-tac pc4 cdm4, clarsimp*)  
**apply** (*subgoal-tac ((pc4, cdm4), bcs4)*  
= *mapAccumL (trSeq ctm com pcc) (pc3, cdm3) svms*)  
**prefer 2 apply simp**  
**apply** (*erule-tac x=pc3 in allE*)  
**apply** (*erule-tac x=cdm3 in allE*)  
**apply** (*erule-tac x=pc4 in allE*)  
**apply** (*erule-tac x=cdm4 in allE*)  
**apply** (*erule impE*)  
**apply** (*rule-tac x=bcs4 in exI, assumption*)  
**apply** (*subgoal-tac cdm1  $\subseteq_m$  cdm3*)  
**prefer 2 apply** (*rule fun-trSeq2, rule sym, simp*)  
**by** (*erule map-le-trans, assumption*)

**lemma fun-trSeq-aux :**

*((pc, cdm), bcs) = mapAccumL (trSeq ctm com pcc) (pcini, cdini) svms  $\longrightarrow$*   
*drop pcini bc = concat bcs  $\longrightarrow$*   
*( $\forall i . i < \text{length } svms \longrightarrow$*   
*( $\exists pc1 cdm1 pc2 cdm2 bcs1 bcs2 bc2 .$*   
*((pc1, cdm1), bcs1) = mapAccumL (trSeq ctm com pcc) (pcini, cdini) (take i svms)*  
 *$\wedge ((pc2, cdm2), bc2) = \text{trSeq ctm com pcc } (pc1, cdm1) (svms ! i)$*   
 *$\wedge ((pc, cdm), bcs2) = \text{mapAccumL (trSeq ctm com pcc) (pc1, cdm1) (drop i svms)$*   
 *$\wedge bcs = bcs1 @ bcs2$*   
 *$\wedge (pc1 - pcini) = \text{length } (\text{concat } bcs1)$*   
 *$\wedge cdm2 \subseteq_m cdm$ )*

**apply** (*rule impI*)  
**apply** (*rule allI*)

**apply** (*induct-tac i*)

**apply** (*rule impI*)  
**apply** (*rule-tac x=pcini in exI*)  
**apply** (*rule-tac x=cdini in exI*)  
**apply** (*subgoal-tac ( $\exists pc2 cdm2 bc2 .$*   
*((pc2, cdm2), bc2) = trSeq ctm com pcc (pcini, cdini) (svms ! 0)*  
 *$\wedge pc2 - pcini = \text{length } bc2$ )*)  
**prefer 2**  
**apply** (*rule fun-trSeq1*)  
**apply** (*elim exE, elim conjE*)  
**apply** (*rule-tac x=pc2 in exI*)  
**apply** (*rule-tac x=cdm2 in exI*)  
**apply** (*rule-tac x=[] in exI*)  
**apply** (*rule-tac x=bcs in exI*)  
**apply** (*rule-tac x=bc2 in exI*)

```

apply (rule conjI,simp)
apply (rule conjI,assumption)
apply (rule conjI,simp)
apply (rule conjI,simp)
apply (rule conjI,simp)
apply (rule conjI,simp)
apply (subgoal-tac  $\exists s' y ys'$ .
  (s',y) = trSeq ctm com pcc (pcini, cdini) (svms!0)
   $\wedge ((pc, cdm),ys') = \text{mapAccumL} (\text{trSeq ctm com pcc}) s' (\text{tl svms}))$ 
  prefer 2 apply (erule mapAccumL-non-empty3,assumption)
apply (elim exE,elim conjE)
apply clarify apply (rename-tac pc2' cdm2' bc2' bcs')
apply (subgoal-tac ((pc2', cdm2'), bc2')=((pc2, cdm2), bc2))
  prefer 2 apply simp
apply clarify
apply (erule fun-trSeq3)

```

```

apply (rule impI, drule mp, simp)
apply (elim exE,elim conjE)

```

```

apply (subgoal-tac drop n svms = (svms ! n) # drop (Suc n) svms)
apply clarify
  prefer 2 apply (rule drop-nth3, simp)
apply (subgoal-tac  $\exists s' y ys'$ .
  (s',y) = trSeq ctm com pcc (pc1, cdm1) (svms ! n)
   $\wedge ((pc, cdm),ys') = \text{mapAccumL} (\text{trSeq ctm com pcc}) s' (\text{drop} (\text{Suc } n) \text{ svms})$ 
   $\wedge \text{bcs2} = y \# ys'$ )
  prefer 2
  apply (rule mapAccumL-non-empty2 [where f=trSeq ctm com pcc])
  apply (blast,simp)
apply (elim exE,elim conjE)
apply (case-tac s')
apply clarify
apply (rename-tac bc2' bcs3 pc2' cdm2')
apply (subgoal-tac ((pc2', cdm2'), bc2')=((pc2, cdm2), bc2))
  prefer 2 apply simp
apply clarify
apply (thin-tac ?x=trSeq ?x1 ?x2 ?x3 ?x4 ?x5)

```

```

apply (rule-tac x=pc2 in exI)
apply (rule-tac x=cdm2 in exI)
apply (subgoal-tac  $\exists pc2a cdm2a bc2a$  .
  ((pc2a, cdm2a), bc2a) = trSeq ctm com pcc (pc2, cdm2) (svms ! Suc n)
   $\wedge pc2a - pc2 = \text{length } bc2a$ )
  prefer 2 apply (rule fun-trSeq1)
apply (elim exE,elim conjE)

```

```

apply (rule-tac x=pc2a in exI)
apply (rule-tac x=cdm2a in exI)
apply (rule-tac x=bcs1 @ [bc2] in exI)
apply (rule-tac x=bcs3 in exI)
apply (rule-tac x=bc2a in exI)

```

```

apply (rule conjI)
apply (subgoal-tac take (Suc n) svms=(take n svms)@[svms ! n])
  prefer 2
  apply (rule take-append3,simp)
apply (erule-tac t=take (Suc n) svms in ssubst)
apply (erule mapAccumL-one-more,assumption)

```

```

apply (rule conjI,assumption)

```

```

apply (rule conjI,assumption)

```

```

apply (rule conjI, simp)

```

```

apply (subgoal-tac  $\exists$  pc2b cdm2b bc2b .
  ((pc2b, cdm2b), bc2b) = trSeq ctm com pcc (pc1, cdm1) (svms ! n)
   $\wedge$  pc2b - pc1 = length bc2b)
  prefer 2
  apply (rule fun-trSeq1)
  apply (elim exE,elim conjE)
apply (subgoal-tac ((pc2, cdm2), bc2)=((pc2b, cdm2b), bc2b))
  prefer 2 apply simp
apply clarify
apply simp
apply (subgoal-tac (pc1 - pcini)+(pc2b - pc1)=
  (length (concat bcs1))+length bc2b))
  prefer 2
  apply (erule sum-nat,assumption)
apply (rule conjI)
apply (rule sym)
apply (subgoal-tac length (concat bcs1) + length bc2b
  = pc1 - pcini + (pc2b - pc1))
  prefer 2 apply simp
apply (erule-tac t=length (concat bcs1) + length bc2b in ssubst)
apply (rule sum-subtract)
apply (rename-tac cdm3 bc3 pc2 cdm2 bc2)

```

```

apply (subgoal-tac  $\exists s' y ys'$ .
  (s',y) = trSeq ctm com pcc (pc2, cdm2) ((drop (Suc n) svms)!0)
   $\wedge ((pc, cdm),ys') =$ 
  mapAccumL (trSeq ctm com pcc) s' (tl (drop (Suc n) svms)))
prefer 2 apply (erule mapAccumL-non-empty3)
apply force
apply (elim exE, elim conjE)
apply clarify apply (rename-tac pc2' cdm3' bc3' bcs')
apply (subgoal-tac ((pc2', cdm3'), bc3')=((pc2a, cdm3), bc3))
prefer 2 apply simp
apply clarify
apply (erule fun-trSeq3)
done

```

**lemma** fun-trSeq :

```

 $\llbracket ((pc, cdm), bcs) = \text{mapAccumL } (trSeq \text{ ctm com pcc}) (pcini, cdini) \text{ svms};$ 
  drop pcini bc = concat bcs;
  l < length svms;
  seq = svms ! l
 $\rrbracket \implies (\exists pc' cdm' pc'' cdm'' bc' rest .$ 
  ((pc'', cdm''), bc') = trSeq ctm com pcc (pc', cdm') seq
   $\wedge$  drop pc' bc = bc' @ rest
   $\wedge$  cdm''  $\subseteq_m$  cdm)

```

```

apply (insert fun-trSeq-aux [of pc cdm bcs ctm com pcc pcini cdini svms bc])
apply (erule mp, assumption)+
apply (erule-tac x=l in allE)
apply (erule mp, assumption)
apply (elim exE, elim conjE)

```

```

apply (rule-tac x=pc1 in exI)
apply (rule-tac x=cdm1 in exI)
apply (rule-tac x=pc2 in exI)
apply (rule-tac x=cdm2 in exI)
apply (rule-tac x=bc2 in exI)
apply (rule-tac x=drop (pc2-pc1) (concat bcs2) in exI)

```

```

apply (rule conjI, simp)

```

```

apply (rule conjI)
apply (thin-tac ?x=mapAccumL ?y ?z ?u)
apply (thin-tac ?x=mapAccumL ?y ?z ?u)
apply (subgoal-tac drop l svms=(svms!l) # drop (Suc l) svms)
prefer 2
apply (erule drop-nth3, clarify)

```



```

apply (case-tac trSeq ctm com pcc (pc1, cdm1) (svms ! l), clarsimp)
apply (rename-tac pc2 cdm2 bc2)
apply (case-tac mapAccumL (trSeq ctm com pcc) (pc2, cdm2) (drop (Suc l) svms), clarsimp)
apply (rename-tac bcs3)
apply (subgoal-tac  $\exists$  pc2' cdm2' bc2' .
  ((pc2', cdm2'), bc2') = trSeq ctm com pcc (pc1, cdm1) (svms ! l)
   $\wedge$  pc2' - pc1 = length bc2')
  prefer 2 apply (rule fun-trSeq1, clarify)
apply (subgoal-tac ((pc2', cdm2'), bc2') = ((pc2, cdm2), bc2))
  prefer 2 apply simp
apply clarify
apply auto
apply (subgoal-tac drop pc1 bc = [] @ [] @ (bc2 @ concat bcs3))
  prefer 2 apply (rule drop-append3)
apply (simp, simp, simp add: drop-append)
done

```

**lemma** fun-trCodeStore:

$$\begin{aligned}
 (P', cdm, ctm, com) = trSVM2JVM ((svms, ctm), ini, ct, st) \longrightarrow \\
 bc = extractBytecode P' \longrightarrow \\
 (\exists cdm' ctm' pcc pcini cdini pc bcs . \\
 ((pc, cdm), bcs) = mapAccumL (trSeq ctm' com pcc) (pcini, cdini) svms \\
 \wedge drop pcini bc = concat bcs)
 \end{aligned}$$

```

apply (rule impI)+
apply (unfold trSVM2JVM-def)
apply clarsimp
apply (case-tac trCodeStore ini (length (initConsTable ct) +
  length (initSizeTable st)) ctm (trConsTable ct) svms)
apply (rename-tac instrs1 rest)
apply (case-tac rest)
apply (rename-tac cdm1 ctm1)
apply clarsimp
apply (subgoal-tac extractBytecode
  ((safeP, objectC, [], [(sigSafeMain, PrimT Void, 10, 10,
  initConsTable ct @ initSizeTable st @ instrs1, [])]) # Prog-RTS) =
  initConsTable ct @ initSizeTable st @ instrs1)
apply simp
apply (thin-tac extractBytecode ?x=?y)
  prefer 2
  apply (unfold extractBytecode-def)
  apply (unfold method'-def)
  apply simp
apply (subgoal-tac class ((safeP, objectC, [], [(sigSafeMain, PrimT Void,
  10, 10, initConsTable ct @ initSizeTable st @ instrs1, [])]) # Prog-RTS)
  safeP = Some (objectC, [], [(sigSafeMain, PrimT Void, 10, 10,
  initConsTable ct @ initSizeTable st @ instrs1, [])]))
apply simp

```

**apply** (*unfold Prog-RTS-def*)  
**apply** (*simp add:class-def*)

**apply** (*unfold trCodeStore-def*)  
**apply** (*case-tac unzip3 ctnmap, rename-tac fs1 rest,*  
*case-tac rest, rename-tac ps1 contss1*)  
**apply** *clarify* **apply** (*simp del: mapAccumL.simps trConsTable-def zip.simps upt-Suc*)  
**apply** (*case-tac mapAccumL (trSeq (map-of (zip (concat contss1)*  
*[Suc 0..<Suc (length (concat contss1))])) (trConsTable ct)*  
*(Suc (length (initConsTable ct) + length (initSizeTable st) + length fs1)))*  
*(Suc (Suc (length (initConsTable ct) + length (initSizeTable st)*  
*+ length fs1)), map-of (zip (zip ps1 (replicate (length fs1) 0))*  
*[Suc (length (initConsTable ct) + length (initSizeTable st))..<*  
*Suc (length (initConsTable ct) + length (initSizeTable st) +*  
*length fs1])) svms*)  
**apply** (*rename-tac rest iss1, case-tac rest, rename-tac pc1 cdm1*)  
**apply** (*clarify*) **apply** (*simp del: mapAccumL.simps trConsTable-def zip.simps*  
*upt-Suc*)  
**apply** (*elim conjE, clarify*)

**apply** (*rule-tac x=map-of (zip (concat contss1)*  
*[Suc 0..<Suc (length (concat contss1))]) in exI*)  
**apply** (*rule-tac x=Suc (length (initConsTable ct) + length (initSizeTable st)*  
*+ length fs1) in exI*)  
**apply** (*rule-tac x=Suc (Suc (length (initConsTable ct)*  
*+ length (initSizeTable st) + length fs1)) in exI*)  
**apply** (*rule-tac x=map-of (zip (zip ps1 (replicate (length fs1) 0))*  
*[Suc (length (initConsTable ct) + length (initSizeTable st))..<*  
*Suc (length (initConsTable ct) + length (initSizeTable st)*  
*+ length fs1]) in exI*)  
**apply** (*rule-tac x=pc1 in exI*)  
**apply** (*rule-tac x=iss1 in exI*)

**apply** (*rule conjI*)  
**apply** *simp*

**apply** (*subgoal-tac length fs1 =length ps1*)  
**prefer 2** **apply** (*rule unzip3-length, simp*)  
**apply** (*subgoal-tac length (map Goto*  
*(zipWith (% p n. int (the (cdm (p, 0))) - int n) ps1*  
*[Suc (length (initConsTable ct) + length (initSizeTable st))..<*  
*Suc (length (initConsTable ct) + length (initSizeTable st) + length fs1]))*  
*= length fs1*)  
**prefer 2**  
**apply** (*subgoal-tac length [Suc (length (initConsTable ct) +*

```

length (initSizeTable st)..<
Suc (length (initConsTable ct) + length (initSizeTable st)
+ length fs1] =
(Suc (length (initConsTable ct) + length (initSizeTable st)
+ length fs1) -
(Suc (length (initConsTable ct) + length (initSizeTable st))))
prefer 2 apply (rule upt-length,simp)
apply (subgoal-tac length (zipWith (% p n. int (the (cdm (p, 0)))
- int n) ps1
[Suc (length (initConsTable ct) + length (initSizeTable st)..<
Suc (length (initConsTable ct) + length (initSizeTable st)
+ length fs1)]) =
min (length ps1) (length [Suc (length (initConsTable ct) +
length (initSizeTable st)..<
Suc (length (initConsTable ct) + length (initSizeTable st)
+ length fs1)])))
prefer 2 apply (rule zipWith-length)
apply simp
apply (simp add: drop-append)
done

```

**axioms** fun-SVM2JVM-aux:

```

[[ l < length svms; svms ! l = (p, seq, fn);
((pc, cdm), bcs) = mapAccumL (trSeq ctm' com pcc) (pcini, cdini) svms;
((pc2, cdm2), bcSeq) = trSeq ctm' com pcc (pcSeq, cdmSeq) (svms ! l)
]] ==> (forall j >= 0. (p, j) not in dom cdmSeq)

```

**lemma** fun-SVM2JVM [rule-format]:

```

(P', cdm, ctm, com) = trSVM2JVM ((svms, ctmap), ini, ct, st) ->
l < length svms ->
svms ! l = (p, seq, fn) ->
i < length seq ->
svm = fst (the (map-of svms p)) ! i ->
pc = the (cdm (p, i)) ->
bytecode = extractBytecode P' ->
(exists cdm' ctm' pcc inss bytecode' .
inss = trInstr pc cdm' ctm' com pcc svm ^
drop pc bytecode = inss @ bytecode' ^
cdm' subset_m cdm)

```

**apply** (rule impI)+

```

apply (insert fun-trCodeStore [of P' cdm ctm com svms ctmap ini ct st bytecode])
apply (drule mp, assumption)+
apply (erule exE)+

```

**apply**(*elim conjE*)

**apply** (*subgoal-tac* ( $\exists pc' cdm' pc'' cdm'' bc' rest .$   
(( $pc'', cdm''$ ),  $bc'$ ) = *trSeq*  $ctm' com pcc (pc', cdm')$  (*svms ! l*)  
 $\wedge drop pc' bytecode = bc' @ rest$   
 $\wedge cdm'' \subseteq_m cdm$ ))  
**prefer** 2 **apply** (*rule fun-trSeq*) **apply**(*simp, assumption, simp, simp*)  
**apply** (*elim exE, elim conjE*)  
**apply** *clarify*  
**apply** (*rename-tac pcSeq cdmSeq pc2 cdm2 bcSeq restSeq*)

**apply** (*subgoal-tac*  $\exists is iss . seq = is\#iss$ )  
**prefer** 2 **apply** (*rule list-non-empty, force*)  
**apply** (*elim exE*)  
**apply** (*subgoal-tac seq=fst (the (map-of svms p))*)  
**prefer** 2 **apply** (*subgoal-tac map-of svms p= Some (seq, fn)*)  
**prefer** 2 **apply** (*rule map-of-distinct*)  
**apply** (*rule svms-good, simp, simp*)  
**apply** *simp*  
**apply** (*insert svms-good [of svms]*)  
**apply** (*subgoal-tac* ( $\exists pc'' cdm'' bc'' rest .$   
 $bc'' = trInstr pc'' cdm'' ctm' com pcc (seq ! i)$   
 $\wedge drop pc'' (extractBytecode P') = bc'' @ rest$   
 $\wedge cdm'' (p, i) = Some pc''$   
 $\wedge cdm'' \subseteq_m cdm2$ ))  
**prefer** 2 **apply** (*rule fun-trInstr*)  
**apply** (*simp*)  
**apply** (*rule sym*)  
**apply** *force*  
**apply** (*erule fun-SVM2JVM-aux [of l svms p seq fn], assumption, simp, simp*)  
**apply** *simp*  
**apply** *simp*  
**apply** *simp*  
**apply** (*elim exE, elim conjE*)

**apply** (*rule-tac x=cdm'' in exI*)  
**apply** (*rule-tac x=ctm' in exI*)  
**apply** (*rule-tac x=pcc in exI*)  
**apply** (*rule-tac x=bc'' in exI*)  
**apply** (*rule-tac x=rest in exI*)  
**apply** (*subgoal-tac cdm''  $\subseteq_m$  cdm*)  
**prefer** 2 **apply** (*erule map-le-trans, assumption*)  
**apply** (*subgoal-tac the (cdm (p, i))=pc''*)  
**prefer** 2 **apply** (*unfold map-le-def*)  
**apply** *force*  
**apply** *clarify*

**apply** *simp*  
**done**

**consts**

*nonJumping* :: *SafeInstr*  $\Rightarrow$  *bool*

**primrec**

*nonJumping* *DECREGION* = *True*  
*nonJumping* (*SLIDE* *m n*) = *True*  
*nonJumping* (*PRIMOP* *oper*) = *True*  
*nonJumping* *REUSE* = *True*  
*nonJumping* *COPY* = *True*  
*nonJumping* (*PUSHCONT* *p*) = *True*  
*nonJumping* (*BUILDENV* *its*) = *True*  
*nonJumping* (*BUILDCLS* *c is i*) = *True*  
*nonJumping* (*CALL* *p*) = *False*  
*nonJumping* *POPCONT* = *False*  
*nonJumping* (*MATCHN* *l v m ps*) = *False*  
*nonJumping* (*MATCH* *l ps*) = *False*  
*nonJumping* (*MATCHD* *l ps*) = *False*

**axioms** *nonJumping-Suc-pc* :

$\llbracket (P', cdm, ctm, com) = trSVM2JVM ((svms, ctm), ini, ct, st);$   
*svm* = *fst* (*the* (*map-of* *svms l*)) ! *i*;  
*nonJumping* *svm*;  
*iss* = *trInstr* (*the* (*cdm* (*l, i*))) *cdm'* *ctm'* *com* *pcc* *svm*  
 $\rrbracket \Rightarrow$  *the* (*cdm* (*l, i+1*)) = *the* (*cdm* (*l, i*)) + *length* *iss*

**end**

## 25 Lemmas on the dynamic properties of the translation SVM to JVM

**theory** *dem-PUSHCONT*

**imports** *../JVMSAFE/JVMExec SVM2JVM SVMSemantics CertifSVM2JVM*  
*dem-translation*

**begin**

**no-translations** *Norm* *s* == (*None*, *s*)

**no-translations** *ex-table-of* *m* == *snd* (*snd* (*snd* *m*))

**declare** *trInstr.simps* [*simp del*]

**declare** *equivS.simps* [*simp del*]

**axioms** *maxPush-PUSHCONT*:

$2 + n < int\ m$

**lemma** *equivS-PUSHCONT-aux*:

*equivS (v1 # S) S' n ctm d g*  
 $\implies n + 1 \geq 0$   
**apply** (*induct S*)  
**apply** (*case-tac v1, simp-all*)  
**apply** (*unfold equivS.simps*)  
**apply** (*case-tac (nat n), simp, simp*)  
**apply** (*case-tac nata, simp, simp*)  
**apply** (*case-tac x, simp*)  
**apply** (*simp add: equivS.simps*)  
**apply** (*case-tac v1, simp-all*)  
**apply** (*simp add: equivS.simps*)  
**apply** (*simp add: equivS.simps*)  
**apply** (*case-tac x*)  
**by** (*simp add: equivS.simps*)

**lemma** *equivS-PUSHCONT*:

*equivS (drop (Suc 0) S) S' n ctm d g  $\implies n + 1 \geq 0$*   
**apply** (*induct S*)  
**apply** (*simp add: equivS.simps, simp*)  
**apply** (*case-tac S, simp*)  
**by** (*simp, erule equivS-PUSHCONT-aux*)

**lemma** *equivS-good*:

*equivS S S' n ctm d g  $\implies 1 \leq 2 + n$*   
**apply** (*induct S*)  
**apply** (*unfold equivS.simps, simp, clarsimp*)  
**apply** (*case-tac a, simp-all*)  
**apply** (*simp add: equivS.simps*)  
**apply** (*simp add: equivS.simps*)  
**apply** (*case-tac x*)  
**by** (*simp add: equivS.simps*)

**axioms** *equiv-ctm*:

*the (ctm p) = the (ctm' p)*

**lemma** *equivS-PUSHCONT-S2*:

$\llbracket \text{equivS } S \text{ } S' \text{ } n \text{ } \text{ctm } d \text{ } g \rrbracket$   
 $\implies \text{equivS } (\text{Cont } (\text{nat } (\text{the-Intg } (\text{the } (\text{sha } (\text{heapC}, \text{k0f}))))), p) \# S$   
 $(S'(\text{nat } (n + 1) \mapsto \text{Intg } (\text{int } (\text{the } (\text{ctm}' p))), \text{nat } (2 + n) \mapsto \text{the } (\text{sha}$   
 $(\text{heapC}, \text{k0f})))) (2 + n) \text{ctm } d \text{ } g$   
**apply** (*frule equivS-good*)  
**apply** (*unfold equivS.simps*)  
**apply** (*rule conjI, simp*)  
**apply** (*rule conjI, clarsimp*)  
**apply** (*rule conjI, auto*)  
**apply** (*rule equiv-ctm*)  
**apply** (*drule equivS-ge-n, erule-tac x=1 in allE, simp*)  
**apply** (*subgoal-tac 1 + n = n + 1, simp*) **prefer 2** **apply** *simp*

by (drule equivS-ge-n, erule-tac x=2 in allE,simp)

**lemma activeCells-PUSHCONT:**

$\llbracket la \notin \text{activeCells regS } h \text{ (nat (the-Intg (the (sha (heapC, kf)))))) \rrbracket$   
 $\implies \text{activeCells regS } h \text{ (nat (the-Intg (the (sha (heapC, kf)))))) =$   
 $\text{activeCells regS } (h(la \mapsto \text{Arr ty } m \text{ (S'(nat (n + 1) \mapsto \text{Intg (int (the (ctm$   
 $p))))), \text{nat } (2 + n) \mapsto \text{the (sha (heapC, k0f))}))$   
 $\text{(nat (the-Intg (the (sha (heapC, kf))))))$   
**apply** (frule l-not-in-cellReg)  
**apply** (unfold activeCells-def,auto)  
**apply** (unfold region-def, simp add: Let-def, elim conjE)  
**apply** (rule-tac x=j in exI,simp)  
**apply** (rule cellReg-step)  
**apply** (rule cellReg-basic)  
**apply** (erule-tac x=j in allE)  
**apply** (rule cellsReg-monotone-1,assumption+)  
**apply** (erule l-not-in-regs)  
**apply** (simp add: Let-def, elim conjE)  
**apply** (rule-tac x=j in exI, simp)  
**apply** (rule cellReg-step)  
**apply** (rule cellReg-basic)  
**apply** (erule-tac x=j in allE)  
**apply** (erule cellsReg-monotone-2,assumption+)  
**by** (erule l-not-in-regs)

**lemma domH-PUSHCONT:**

$\llbracket la \notin \text{activeCells regS } h \text{ (nat (the-Intg (the (sha (heapC, kf)))))) \rrbracket;$   
 $\forall l \in \text{dom } H. \exists l'. l' = \text{the (g } l) \wedge \text{equivC (the (H } l)) h l' \text{ (the (h } l')) \text{ (nat$   
 $\text{(the-Intg (the (sha (heapC, kf)))))) com regS } d \text{ g} \rrbracket$   
 $\implies \forall l \in \text{dom } H.$   
 $\exists l'. l' = \text{the (g } l) \wedge$   
 $\text{equivC (the (H } l)) (h(la \mapsto \text{Arr ty } m \text{ (S'(nat (n + 1) \mapsto \text{Intg (int$   
 $\text{(the (ctm } p))))),$   
 $\text{nat } (2 + n) \mapsto \text{the (sha (heapC,$   
 $\text{k0f))}))$   
 $l' \text{ (the ((h(la \mapsto \text{Arr ty } m \text{ (S'(nat (n + 1) \mapsto \text{Intg (int (the (ctm$   
 $p))))),$   
 $\text{nat } (2 + n) \mapsto \text{the (sha (heapC, k0f))})) l'))$   
 $\text{(nat (the-Intg (the (sha (heapC, kf))))))$   
 $\text{com regS } d \text{ g}$   
**apply** (frule l-not-in-cellReg)  
**apply** (rule ballI)

```

apply (erule-tac x=l in ballE)
  prefer 2 apply simp
apply (erule exE)
apply (erule-tac x=l' in exI) apply (elim conjE)
  apply (rule conjI, assumption)
apply clarsimp
apply (rule conjI) apply (rule impI)
apply clarsimp apply (simp add: region-def)
apply clarsimp
apply (erule-tac x=Obj in exI)
apply (erule-tac x=flds in exI) apply simp
apply (simp add: region-def) apply (elim conjE)
apply (rule cellReg-step)
  apply (rule cellReg-basic)
  apply (erule-tac x=j in allE)
  apply (rule cellsReg-monotone-1, assumption+)
by (erule l-not-in-regs)

```

```

declare equivH.simps [simp del]

```

```

lemma equivH-PUSHCONT:

```

```

  [[ (hm, k), k0, (l, i), S) = ((H, ka), k0a, PC, Sa);
    k' = nat (the-Intg (the (sha (heapC, kf))));
    equivH (H, ka) h k' com regS d g;
    activeCells regS h k' ∩ {la, l', l''} = {}]]
  ⇒ equivH (hm, k) (h(la ↦ Arr ty m (S'(nat (n + 1) ↦ Intg (int (the (ctm
p))))),
    nat (2 + n) ↦ the (sha (heapC, k0f))))))
  (nat (the-Intg (the ((sha((stackC, topf) ↦ Intg (n + 1 + 1),
(heapC, k0f) ↦ the (sha (heapC, kf))))) (heapC,
kf)))))) com regS d g
apply simp apply (elim conjE)
apply (erule l-not-in-cellReg)
apply (simp add: heapC-def add: stackC-def)
apply (simp add: kf-def add: k0f-def)
apply (fold heapC-def, fold kf-def, fold k0f-def)
apply (unfold equivH.simps)
apply (elim conjE)
apply (rule conjI, simp)
apply (rule conjI) apply (erule activeCells-PUSHCONT, simp)
apply (rule conjI, assumption)
by (rule domH-PUSHCONT, assumption+)

```

```

declare exec-instr.simps [simp]

```



**lemma** *execSVMInstr-PUSHCONT* :

$$\begin{aligned} & \llbracket (P', cdm, ctm, com) = trSVM2JVM ((svms, ctmap), ini, ct, ah, ai, bc); \\ & \quad cdm, ctm, com \vdash ((hm, k), k0, (l, i), S) \triangleq S1'; \\ & \quad (fst (the (map-of svms l)) ! i) = PUSHCONT p; \\ & \quad execSVMInst (PUSHCONT p) (map-of ct) (hm, k) k0 (l, i) S = Either.Right \\ & \quad S2; \\ & \quad drop (the (cdm (l, i))) (extractBytecode P') = \\ & \quad trInstr (the (cdm (l, i))) cdm' ctm' com pcc (PUSHCONT p) @ bytecode' \\ & \rrbracket \implies \exists v' sh' dh' ih' fms'. \\ & P' \vdash S1' -jvm \rightarrow (v', sh', dh', ih', fms') \wedge \\ & cdm, ctm, com \vdash S2 \triangleq (v', sh', dh', ih', fms') \end{aligned}$$

**apply** (*case-tac*  $S1'$ )  
**apply** (*rename-tac*  $v \text{ tup}$ )  
**apply** (*case-tac*  $\text{tup}$ )  
**apply** (*rename-tac*  $sh \text{ tup}$ )  
**apply** (*case-tac*  $\text{tup}$ )  
**apply** (*rename-tac*  $dh \text{ tup}$ )  
**apply** (*case-tac*  $\text{tup}$ )  
**apply** (*rename-tac*  $ih \text{ fms}$ )  
**apply** (*simp*)

**apply** (*subgoal-tac*  $\exists S'. S = drop (Suc 0) S'$ )  
**prefer** 2 **apply** (*rule-tac*  $x = snd (snd (snd S2))$ ) **in**  $exI, clarsimp$ )  
**apply** (*unfold equivState-def*)  
**apply** (*elim exE, elim conjE*)

**apply** (*rule-tac*  $x = None$ ) **in**  $exI$ )  
**apply** (*rule-tac*  $x = sha((stackC, topf) \mapsto Intg (n + 1 + 1), (heapC, k0f) \mapsto the (sha (heapC, kf)))$ ) **in**  $exI$ )  
**apply** (*rule-tac*  $x = h(la \mapsto Arr \text{ ty } m (S'a(nat (n + 1) \mapsto Intg (int (the (ctm' p))), nat (2 + n) \mapsto the (sha (heapC, k0f))))$ ) **in**  $exI$ )  
**apply** (*rule-tac*  $x = inih$ ) **in**  $exI$ )  
**apply** (*rule-tac*  $x = [([], vs, safeP, sigSafeMain, the (cdm (l, i)) + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1, ref)]$ ) **in**  $exI$ )

**apply** (*subgoal-tac*  $\exists vs \text{ pc} . fms = ([(), vs, safeP, sigSafeMain, pc, ref]) \# []$ )  
**prefer** 2 **apply** *simp*  
**apply** (*erule exE*) +  
**apply** (*unfold exec-all-def*)  
**apply** (*rule conjI*)

```

apply (subgoal-tac PC = (l,i)) prefer 2 apply simp
apply simp apply (elim conjE) apply clarsimp
apply (subgoal-tac P⊢ (None,sha,h,inih,[[[], vs, safeP, sigSafeMain, the (cdm
(l,i)), ag,bd]]) -jvm→
      (None,sha,h,inih,[[the (sha (stackC,topf))], vs, safeP, sigSafeMain,
the (cdm (l,i)) + 1, ag,bd]]))
apply (unfold exec-all-def)
apply (erule rtrancl-trans)

```

```

prefer 2
apply (rule r-into-rtrancl)
apply (clarify)
apply (unfold JVMEexec.exec.simps)
apply (unfold Let-def)
apply (unfold extractBytecode-def)
apply (subgoal-tac (fst (snd (snd (snd (snd (the (method' (P', safeP) sigSafe-
Main))))))) !
      the (cdm (l, i))) = Getstatic topf stackC,simp)
apply (unfold trInstr.simps, unfold Let-def)
apply (erule nth-via-drop-append)

```

```

apply (drule drop-Suc-append)
apply (subgoal-tac P⊢ (None, sha, h, iniH,
      [[the (sha (stackC, topf))],
      vs, safeP, sigSafeMain, the (cdm (l, i)) + 1, ag, bd]])
      -jvm→
      (None,sha,h,inih,
      [[Intg 1,
      the (sha (stackC,topf))],
      vs, safeP, sigSafeMain, the (cdm (l, i)) + 1 + 1, ag, bd]]))

```

```

apply (unfold exec-all-def)
apply (erule rtrancl-trans)
prefer 2
apply (rule r-into-rtrancl)
apply (clarify)
apply (unfold JVMEexec.exec.simps)
apply (unfold Let-def)
apply (subgoal-tac (fst (snd (snd (snd (snd (the (method' (P', safeP) sigSafe-
Main))))))) ! Suc (the (cdm (l, i))) =
      LitPush (Intg 1),simp)
apply (erule nth-via-drop-append)

```

```

apply (drule drop-Suc-append)
apply (subgoal-tac P⊢ (None,sha,h,inih,
      [[Intg 1,
      the (sha (stackC,topf))],

```

```

      vs, safeP, sigSafeMain, the (cdm (l, i)) + 1 + 1, ag, bd))]
    -jvm→
      (None, sha, h, inih,
       [[Intg (the-Intg (the (sha (stackC, topf))) + 1)],
        vs, safeP, sigSafeMain, the (cdm (l, i)) + 1 + 1 + 1, ag, bd]]))
  apply (unfold exec-all-def)
  apply (erule rtrancl-trans)
  prefer 2
  apply (rule r-into-rtrancl)
  apply (clarify)
  apply (unfold JVMExec.exec.simps)
  apply (unfold Let-def)
  apply (subgoal-tac (fst (snd (snd (snd (snd (the (method' (P', safeP) sigSafe-
Main)))))) ! (Suc (Suc (the (cdm (l, i)))))) =
              BinOp Add, simp))
  apply (erule nth-via-drop-append)

```

```

  apply (drule drop-Suc-append)
  apply (subgoal-tac P⊢ (None, sha, h, inih,
    [[Intg (the-Intg (the (sha (stackC, topf))) + 1)],
     vs, safeP, sigSafeMain, the (cdm (l, i)) + 1 + 1 + 1, ag, bd]]
    -jvm→
      (None, sha((stackC, topf) ↦ Intg (n + 1)), h, inih,
       [([],
        vs, safeP, sigSafeMain, the (cdm (l, i)) + 1 + 1 + 1 + 1, ag,
        bd]]))
  apply (unfold exec-all-def)
  apply (erule rtrancl-trans)
  prefer 2
  apply (rule r-into-rtrancl)
  apply (clarify)
  apply (unfold JVMExec.exec.simps)
  apply (unfold Let-def)
  apply (subgoal-tac (fst (snd (snd (snd (snd (the (method' (P', safeP) sigSafe-
Main)))))) ! (Suc (Suc (Suc (the (cdm (l, i)))))) =
              Putstatic topf stackC, simp))
  apply (erule nth-via-drop-append)

```

```

  apply (drule drop-Suc-append)
  apply (subgoal-tac P⊢ (None, sha((stackC, topf) ↦ Intg (n + 1)), h, inih,
    [([],
     vs, safeP, sigSafeMain, the (cdm (l, i)) + 1 + 1 + 1 + 1,
     ag, bd]]
    -jvm→

```

```

      (None, sha((stackC, topf) ↦ Intg (n + 1)), h, inih,
      [[the (sha (stackC, Sf))],
      vs, safeP, sigSafeMain, the (cdm (l, i)) + 1 + 1 + 1 + 1 +
1, ag, bd]]))
apply (unfold exec-all-def)
apply (erule rtrancl-trans)
prefer 2
  apply (rule r-into-rtrancl)
  apply (clarify)
  apply (unfold JVMExec.exec.simps)
  apply (unfold Let-def)
  apply (subgoal-tac (fst (snd (snd (snd (snd (the (method' (P', safeP) sigSafe-
Main))))))) !
      (Suc (Suc (Suc (Suc (the (cdm (l, i))))))) =
      Getstatic Sf stackC, simp)
  apply (simp add: Sf-def add: topf-def)
apply (erule nth-via-drop-append)

apply (drule drop-Suc-append)
apply (subgoal-tac P+ (None, sha((stackC, topf) ↦ Intg (n + 1)), h, inih,
      [[the (sha (stackC, Sf))],
      vs, safeP, sigSafeMain, the (cdm (l, i)) + 1 + 1 + 1 + 1 +
1, ag, bd]))
      -jvm→
      (None, sha((stackC, topf) ↦ Intg (n + 1)), h, inih,
      [[Intg (the-Intg (the (sha (stackC, topf))) + 1), the (sha
(stackC, Sf))],
      vs, safeP, sigSafeMain, the (cdm (l, i)) + 1 + 1 + 1 + 1 +
1 + 1, ag, bd]]))
apply (unfold exec-all-def)
apply (erule rtrancl-trans)
prefer 2
  apply (rule r-into-rtrancl)
  apply (clarify)
  apply (unfold JVMExec.exec.simps)
  apply (unfold Let-def)
  apply (subgoal-tac (fst (snd (snd (snd (snd (the (method' (P', safeP) sigSafe-
Main))))))) !
      (Suc (Suc (Suc (Suc (Suc (the (cdm (l, i))))))) =
      Getstatic topf stackC, simp)
apply (erule nth-via-drop-append)

apply (drule drop-Suc-append)
apply (subgoal-tac P+ (None, sha((stackC, topf) ↦ Intg (n + 1)), h, inih,
      [[Intg (the-Intg (the (sha (stackC, topf))) + 1), the (sha
(stackC, Sf))],
      vs, safeP, sigSafeMain, the (cdm (l, i)) + 1 + 1 + 1 + 1 +

```

```

1 + 1, ag, bd)])
  -jvm→
    (None,sha((stackC, topf) ↦ Intg (n + 1)),h,inih,
    [([Intg (int (the (ctm' p))),
      Intg (the-Intg (the (sha (stackC, topf))) + 1),
      the (sha (stackC, Sf))],
      vs, safeP, sigSafeMain, the (cdm (l, i)) + 1 + 1 + 1 + 1 +
1 + 1 + 1, ag, bd)))]
apply (unfold exec-all-def)
apply (erule rtrancl-trans)
prefer 2
apply (rule r-into-rtrancl)
apply (clarify)
apply (unfold JVMExec.exec.simps)
apply (unfold Let-def)
apply (subgoal-tac (fst (snd (snd (snd (snd (the (method' (P', safeP) sigSafe-
Main))))))) !
      (Suc (Suc (Suc (Suc (Suc (Suc (the (cdm (l, i)))))))))) =
      LitPush (Intg (int (the (ctm' p))),simp)
apply (erule nth-via-drop-append)

apply (drule drop-Suc-append)
apply (subgoal-tac P ⊢ (None,sha((stackC, topf) ↦ Intg (n + 1)),h,inih,
    [([Intg (int (the (ctm' p))),
      Intg (the-Intg (the (sha (stackC, topf))) + 1),
      the (sha (stackC, Sf))],
      vs, safeP, sigSafeMain, the (cdm (l, i)) + 1 + 1 + 1 + 1 +
1 + 1 + 1, ag, bd)])
  -jvm→
    (None,sha((stackC, topf) ↦ Intg (n + 1)),
      h(la ↦ Arr ty m (S'a(nat (n + 1) ↦ Intg (int (the (ctm'
p)))))),inih,
    [([],
      vs, safeP, sigSafeMain, the (cdm (l, i)) + 1 + 1 + 1 + 1 +
1 + 1 + 1 + 1, ag, bd)))]
apply (unfold exec-all-def)
apply (erule rtrancl-trans)
prefer 2
apply (rule r-into-rtrancl)
apply (clarify)
apply (unfold JVMExec.exec.simps)
apply (unfold Let-def)
apply (subgoal-tac (fst (snd (snd (snd (snd (the (method' (P', safeP) sigSafe-
Main))))))) !
      (Suc (Suc (Suc (Suc (Suc (Suc (Suc (the (cdm (l, i)))))))))) =
      ArrStore,simp)
apply (simp add: raise-system-xcpt-def)

```

```

apply (frule equivS-PUSHCONT,simp)
apply (subgoal-tac 2 + n < int m,simp)
apply (rule maxPush-PUSHCONT)
apply (erule nth-via-drop-append)

apply (drule drop-Suc-append)
apply (subgoal-tac P⊢ (None,sha((stackC, topf) ↦ Intg (n + 1)),
      h(la ↦ Arr ty m (S'a(nat (n + 1) ↦ Intg (int (the (ctm'
p))))))),inih,
      [([],
      vs, safeP, sigSafeMain, the (cdm (l, i)) + 1 + 1 + 1 + 1 +
1 + 1 + 1 + 1, ag, bd)])
      -jvm→
      (None,sha((stackC, topf) ↦ Intg (n + 1)),
      h(la ↦ Arr ty m (S'a(nat (n + 1) ↦ Intg (int (the (ctm'
p))))))),inih,
      ([ Intg (n + 1)],
      vs, safeP, sigSafeMain, the (cdm (l, i)) + 1 + 1 + 1 + 1 +
1 + 1 + 1 + 1 + 1, ag, bd))))
apply (unfold exec-all-def)
apply (erule rtrancl-trans)
prefer 2
apply (rule r-into-rtrancl)
apply (clarify)
apply (unfold JVMEexec.exec.simps)
apply (unfold Let-def)
apply (subgoal-tac (fst (snd (snd (snd (snd (the (method' (P', safeP) sigSafe-
Main))))))) !
      (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (the (cdm (l, i))))))))))))))
=
      Getstatic topf stackC,simp)
apply (erule nth-via-drop-append)

```

```

apply (drule drop-Suc-append)
apply (subgoal-tac P⊢ (None,sha((stackC, topf) ↦ Intg (n + 1)),
      h(la ↦ Arr ty m (S'a(nat (n + 1) ↦ Intg (int (the (ctm'
p))))))),inih,
      [([Intg (n + 1)],
      vs, safeP, sigSafeMain, the (cdm (l, i)) + 1 + 1 + 1 + 1 +
1 + 1 + 1 + 1 + 1, ag, bd)])
      -jvm→
      (None,sha((stackC, topf) ↦ Intg (n + 1)),
      h(la ↦ Arr ty m (S'a(nat (n + 1) ↦ Intg (int (the (ctm'
p))))))),inih,
      [([Intg 1,

```

$$\begin{aligned} & \text{Intg } (n + 1)], \\ & \text{vs, safeP, sigSafeMain, the (cdm (l, i)) + 1 + 1 + 1 + 1 +} \\ & 1 + 1 + 1 + 1 + 1 + 1, \text{ ag, bd}]])) \\ \text{apply (unfold exec-all-def)} \\ \text{apply (erule rtrancl-trans)} \\ \text{prefer } 2 \\ \text{apply (rule r-into-rtrancl)} \\ \text{apply (clarify)} \\ \text{apply (unfold JVMExec.exec.simps)} \\ \text{apply (unfold Let-def)} \\ \text{apply (subgoal-tac (fst (snd (snd (snd (snd (the (method' (P', safeP) sigSafe-} \\ \text{Main))))))))) !} \\ & \text{(Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (the (cdm (l,} \\ \text{i))))))))))))) =} \\ & \text{LitPush (Intg 1),simp)} \\ \text{apply (erule nth-via-drop-append)} \\ \\ \text{apply (drule drop-Suc-append)} \\ \text{apply (subgoal-tac } P \vdash \text{ (None, sha((stackC, topf) } \mapsto \text{Intg } (n + 1)), \\ & \text{h(la } \mapsto \text{Arr ty m (S'a(nat (n + 1) } \mapsto \text{Intg (int (the (ctm' } \\ \text{p)))))), inih,} \\ & \text{[[[Intg 1,} \\ & \text{Intg } (n + 1)],} \\ & \text{vs, safeP, sigSafeMain, the (cdm (l, i)) + 1 + 1 + 1 + 1 +} \\ & 1 + 1 + 1 + 1 + 1 + 1, \text{ ag, bd}]])) \\ & \text{-jvm} \mapsto \\ & \text{(None, sha((stackC, topf) } \mapsto \text{Intg } (n + 1)), \\ & \text{h(la } \mapsto \text{Arr ty m (S'a(nat (n + 1) } \mapsto \text{Intg (int (the (ctm' } \\ \text{p)))))), inih,} \\ & \text{[[[Intg (n + 1 + 1)],} \\ & \text{vs, safeP, sigSafeMain, the (cdm (l, i)) + 1 + 1 + 1 + 1 +} \\ & 1 + 1 + 1 + 1 + 1 + 1, \text{ ag, bd}]])) \\ \text{apply (unfold exec-all-def)} \\ \text{apply (erule rtrancl-trans)} \\ \text{prefer } 2 \\ \text{apply (rule r-into-rtrancl)} \\ \text{apply (clarify)} \\ \text{apply (unfold JVMExec.exec.simps)} \\ \text{apply (unfold Let-def)} \\ \text{apply (subgoal-tac (fst (snd (snd (snd (snd (the (method' (P', safeP) sigSafe-} \\ \text{Main))))))))) !} \\ & \text{(Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (the (cdm} \\ \text{(l, i))))))))))))) =} \\ & \text{BinOp Add, simp)} \\ \text{apply (erule nth-via-drop-append)} \end{aligned}$$

**apply** (*drule drop-Suc-append*)  
**apply** (*subgoal-tac*  $P \vdash (\text{None}, \text{sha}((\text{stackC}, \text{topf}) \mapsto \text{Intg } (n + 1)),$   
 $h(\text{la} \mapsto \text{Arr ty m } (S'a(\text{nat } (n + 1) \mapsto \text{Intg } (\text{int } (\text{the } (\text{ctm}'$   
 $p))))))$ ), *inih*,  
 $[[[\text{Intg } (n + 1 + 1)],$   
 $vs, \text{safeP}, \text{sigSafeMain}, \text{the } (\text{cdm } (l, i)) + 1 + 1 + 1 + 1 +$   
 $1 + 1 + 1 + 1 + 1 + 1 + 1, \text{ag}, \text{bd}]]$ )  
 $-jvm \rightarrow$   
 $(\text{None}, \text{sha}((\text{stackC}, \text{topf}) \mapsto \text{Intg } (n + 1 + 1)),$   
 $h(\text{la} \mapsto \text{Arr ty m } (S'a(\text{nat } (n + 1) \mapsto \text{Intg } (\text{int } (\text{the } (\text{ctm}'$   
 $p))))))$ ), *inih*,  
 $[[[],$   
 $vs, \text{safeP}, \text{sigSafeMain}, \text{the } (\text{cdm } (l, i)) + 1 + 1 + 1 + 1 +$   
 $1 + 1 + 1 + 1 + 1 + 1 + 1 + 1, \text{ag}, \text{bd}]]$ )  
**apply** (*unfold exec-all-def*)  
**apply** (*erule rtrancl-trans*)  
**prefer** 2  
**apply** (*rule r-into-rtrancl*)  
**apply** (*clarify*)  
**apply** (*unfold JVMExec.exec.simps*)  
**apply** (*unfold Let-def*)  
**apply** (*subgoal-tac* ( $\text{fst } (\text{snd } (\text{snd } (\text{snd } (\text{snd } (\text{the } (\text{method}' (P', \text{safeP}) \text{sigSafeMain}))))))$ ) !  
 $(\text{Suc } (\text{Suc } (\text{Suc } (\text{Suc } (\text{Suc } (\text{Suc } (\text{Suc } (\text{Suc } (\text{Suc } (\text{Suc } (\text{Suc } (\text{the } (\text{cdm } (l, i)))))))))))))) =$   
 $\text{Putstatic topf stackC, simp}$ )  
**apply** (*erule nth-via-drop-append*)

**apply** (*drule drop-Suc-append*)  
**apply** (*subgoal-tac*  $P \vdash (\text{None}, \text{sha}((\text{stackC}, \text{topf}) \mapsto \text{Intg } (n + 1 + 1)),$   
 $h(\text{la} \mapsto \text{Arr ty m } (S'a(\text{nat } (n + 1) \mapsto \text{Intg } (\text{int } (\text{the } (\text{ctm}'$   
 $p))))))$ ), *inih*,  
 $[[[],$   
 $vs, \text{safeP}, \text{sigSafeMain}, \text{the } (\text{cdm } (l, i)) + 1 + 1 + 1 + 1 +$   
 $1 + 1 + 1 + 1 + 1 + 1 + 1 + 1, \text{ag}, \text{bd}]]$ )  
 $-jvm \rightarrow$   
 $(\text{None}, \text{sha}((\text{stackC}, \text{topf}) \mapsto \text{Intg } (n + 1 + 1)),$   
 $h(\text{la} \mapsto \text{Arr ty m } (S'a(\text{nat } (n + 1) \mapsto \text{Intg } (\text{int } (\text{the } (\text{ctm}'$   
 $p))))))$ ), *inih*,  
 $[[[\text{Addr } \text{la}],$   
 $vs, \text{safeP}, \text{sigSafeMain}, \text{the } (\text{cdm } (l, i)) + 1 + 1 + 1 + 1 +$   
 $1 + 1 + 1 + 1 + 1 + 1 + 1 + 1, \text{ag}, \text{bd}]]$ )  
**apply** (*unfold exec-all-def*)  
**apply** (*erule rtrancl-trans*)  
**prefer** 2  
**apply** (*rule r-into-rtrancl*)  
**apply** (*clarify*)



**apply** (*unfold JVMExec.exec.simps*)  
**apply** (*unfold Let-def*)  
**apply** (*subgoal-tac (fst (snd (snd (snd (snd (the (method' (P', safeP) sigSafeMain)))))) !*  
*(Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (the (cdm (l, i)))))))))))))) =*  
*Getstatic Sf stackC,simp)*)  
**apply** (*simp add: Sf-def add: topf-def*)  
**apply** (*erule nth-via-drop-append*)

**apply** (*drule drop-Suc-append*)  
**apply** (*subgoal-tac P<sup>⊢</sup> (None,sha((stackC, topf) ↦ Intg (n + 1 + 1)),*  
*h(la ↦ Arr ty m (S'a(nat (n + 1) ↦ Intg (int (the (ctm'*  
*p))))),inih,*  
*[[[Addr la],*  
*vs, safeP, sigSafeMain, the (cdm (l, i)) + 1 + 1 + 1 + 1 +*  
*1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1, ag, bd]]]*  
*-jvm→*  
*(None,sha((stackC, topf) ↦ Intg (n + 1 + 1)),*  
*h(la ↦ Arr ty m (S'a(nat (n + 1) ↦ Intg (int (the (ctm'*  
*p))))),inih,*  
*[[[Intg (n + 1 + 1),*  
*Addr la],*  
*vs, safeP, sigSafeMain, the (cdm (l, i)) + 1 + 1 + 1 + 1 +*  
*1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1, ag, bd]]])*)  
**apply** (*unfold exec-all-def*)  
**apply** (*erule rtrancl-trans*)  
**prefer** 2  
**apply** (*rule r-into-rtrancl*)  
**apply** (*clarify*)  
**apply** (*unfold JVMExec.exec.simps*)  
**apply** (*unfold Let-def*)  
**apply** (*subgoal-tac (fst (snd (snd (snd (snd (the (method' (P', safeP) sigSafeMain)))))) !*  
*(Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (the (cdm (l, i)))))))))))))) =*  
*Getstatic topf stackC,simp)*)  
**apply** (*erule nth-via-drop-append*)

**apply** (*drule drop-Suc-append*)  
**apply** (*subgoal-tac P<sup>⊢</sup> (None,sha((stackC, topf) ↦ Intg (n + 1 + 1)),*  
*h(la ↦ Arr ty m (S'a(nat (n + 1) ↦ Intg (int (the (ctm'*  
*p))))),inih,*  
*[[[Intg (n + 1 + 1),*  
*Addr la],*  
*vs, safeP, sigSafeMain, the (cdm (l, i)) + 1 + 1 + 1 + 1 +*

```

1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1, ag, bd]])
  -jvm→
    (None,sha((stackC, topf) ↦ Intg (n + 1 + 1)),
      h(la ↦ Arr ty m (S'a(nat (n + 1) ↦ Intg (int (the (ctm'
p)))))),inih,
      [([the (sha (heapC,k0f)),
        Intg (n + 1 + 1),
        Addr la],
        vs, safeP, sigSafeMain, the (cdm (l, i)) + 1 + 1 + 1 + 1 +
1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1, ag, bd]])
apply (unfold exec-all-def)
apply (erule rtrancl-trans)
prefer 2
apply (rule r-into-rtrancl)
apply (clarify)
apply (unfold JVMExec.exec.simps)
apply (unfold Let-def)
apply (subgoal-tac (fst (snd (snd (snd (snd (the (method' (P', safeP) sigSafe-
Main))))))) !
      (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc
(Suc (Suc (the (cdm (l, i)))))))))))))))))) =
      Getstatic k0f heapC,simp)
apply (simp add: heapC-def add: stackC-def add: k0f-def add: topf-def)
apply (erule nth-via-drop-append)

apply (drule drop-Suc-append)
apply (subgoal-tac P⊢ (None,sha((stackC, topf) ↦ Intg (n + 1 + 1)),
      h(la ↦ Arr ty m (S'a(nat (n + 1) ↦ Intg (int (the (ctm'
p)))))),inih,
      [([the (sha (heapC,k0f)),
        Intg (n + 1 + 1),
        Addr la],
        vs, safeP, sigSafeMain, the (cdm (l, i)) + 1 + 1 + 1 + 1 +
1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1, ag, bd]])
  -jvm→
    (None,sha((stackC, topf) ↦ Intg (n + 1 + 1)),
      h(la ↦ Arr ty m (S'a(nat (n + 1) ↦ Intg (int (the (ctm'
p))))), nat (2 + n) ↦ the (sha (heapC, k0f))))),
      inih,
      [([],
        vs, safeP, sigSafeMain,
        the (cdm (l, i)) + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 +
1 + 1 + 1 + 1 + 1 + 1 + 1, ag, bd]])
apply (unfold exec-all-def)
apply (erule rtrancl-trans)
prefer 2
apply (rule r-into-rtrancl)
apply (clarify)

```

```

apply (unfold JVMEExec.exec.simps)
apply (unfold Let-def)
apply (subgoal-tac (fst (snd (snd (snd (snd (the (method' (P', safeP) sigSafe-
Main))))))) !
      (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc
(Suc (Suc (Suc (the (cdm (l, i)))))))))))))))))) =
      ArrStore,simp)
apply (simp add: raise-system-xcpt-def)
apply (subgoal-tac int m > 2 + n,simp)
prefer 2 apply (rule maxPush-PUSHCONT)
apply (rule conjI, rule impI)
apply (frule equivS-PUSHCONT,simp)
apply (rule impI)
apply (frule equivS-PUSHCONT,simp)
apply (erule nth-via-drop-append)

apply (drule drop-Suc-append)
apply (subgoal-tac P⊢ (None,sha((stackC, topf) ↦ Intg (n + 1 + 1)),
      h(la ↦ Arr ty m (S'a(nat (n + 1) ↦ Intg (int (the (ctm'
p))), nat (2 + n) ↦ the (sha (heapC, kOf))))),
      inih,
      [([],
      vs, safeP, sigSafeMain,
      the (cdm (l, i)) + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 +
1 + 1 + 1 + 1 + 1 + 1 + 1 + 1, ag, bd)])
      -jvm→
      (None,sha((stackC, topf) ↦ Intg (n + 1 + 1)),
      h(la ↦ Arr ty m (S'a(nat (n + 1) ↦ Intg (int (the (ctm'
p))), nat (2 + n) ↦ the (sha (heapC, kOf))))),
      inih,
      [([the (sha (heapC, kf))],
      vs, safeP, sigSafeMain,
      the (cdm (l, i)) + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 +
1 + 1 + 1 + 1 + 1 + 1 + 1 + 1, ag, bd)]))
apply (unfold exec-all-def)
apply (erule rtrancl-trans)
prefer 2
apply (rule r-into-rtrancl)
apply (clarify)
apply (unfold JVMEExec.exec.simps)
apply (unfold Let-def)
apply (subgoal-tac (fst (snd (snd (snd (snd (the (method' (P', safeP) sigSafe-
Main))))))) !
      (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc
(Suc (Suc (Suc (the (cdm (l, i)))))))))))))))))) =
      Getstatic kf heapC,simp)
apply (simp add: heapC-def add: stackC-def add: kf-def add: topf-def)

```

**apply** (*erule nth-via-drop-append*)

**apply** (*drule drop-Suc-append*)

**apply** (*subgoal-tac*  $P \vdash (None, sha((stackC, topf) \mapsto Intg (n + 1 + 1)),$   
 $h(la \mapsto Arr\ ty\ m\ (S'a(nat\ (n + 1) \mapsto Intg\ (int\ (the\ (ctm'$   
 $p))))), nat\ (2 + n) \mapsto the\ (sha\ (heapC, k0f))))))$ ),

*inih,*

$[[the\ (sha\ (heapC, kf))],$

*vs, safeP, sigSafeMain,*

$the\ (cdm\ (l, i) + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 +$

$1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1, ag, bd)]])$

*-jvm→*

$(None, sha((stackC, topf) \mapsto Intg (n + 1 + 1), (heapC, k0f)$

$\mapsto the\ (sha\ (heapC, kf))),$

$h(la \mapsto Arr\ ty\ m\ (S'a(nat\ (n + 1) \mapsto Intg\ (int\ (the\ (ctm'$   
 $p))))), nat\ (2 + n) \mapsto the\ (sha\ (heapC, k0f))))))$ ),

*inih,*

$[[[],$

*vs, safeP, sigSafeMain,*

$the\ (cdm\ (l, i) + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 +$

$1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1, ag, bd)]])$

**apply** (*unfold exec-all-def*)

**apply** (*erule rtrancl-trans*)

**prefer** 2

**apply** (*rule r-into-rtrancl*)

**apply** (*clarify*)

**apply** (*unfold JVMExec.exec.simps*)

**apply** (*unfold Let-def*)

**apply** (*subgoal-tac* (*fst* (*snd* (*snd* (*snd* (*snd* (*the* (*method'* (*P'*, *safeP*) *sigSafe-*  
*Main*))))))))) !

$(Suc\ (Suc\ (Suc\ (Suc\ (Suc\ (Suc\ (Suc\ (Suc\ (Suc$

$(Suc\ (Suc\ (Suc\ (Suc\ (Suc\ (Suc\ (Suc\ (Suc\ (the\ (cdm\ (l,$

$i))))))))))))))))) =$

*Putstatic k0f heapC, simp*)

**apply** (*erule nth-via-drop-append*)

**apply** *simp*

**apply** (*frule nonJumping-Suc-pc*)

**apply** (*erule-tac sym* [**where**  $t = PUSHCONT\ p$ ])

**apply** (*simp add: nonJumping.simps*)

**apply** (*simp add: trInstr.simps*)

**apply** (*rule-tac*  $x = hm$  **in** *exI*)

**apply** (*rule-tac*  $x = k$  **in** *exI*)

```

apply (rule-tac x=k in exI)
apply (rule-tac x=(l,Suc i) in exI)
apply (rule-tac x=Cont (k0, p) # S in exI)

apply (rule-tac x= sha((stackC, topf) ↦ Intg (n + 1 + 1), (heapC, k0f) ↦ the
  (sha (heapC, kf))) in exI)
apply (rule-tac x= h(la ↦ Arr ty m (S'a(nat (n + 1) ↦ Intg (int (the (ctm'
  p))),
  nat (2 + n) ↦ the (sha (heapC, k0f)))))) in exI)
apply (rule-tac x=inih in exI)
apply (rule-tac x=vs in exI)
apply (rule-tac x=the (cdm (l, i)) + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 +
  1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 in exI)
apply (rule-tac x=ref in exI)

apply (rule-tac x=nat (the-Intg (the ((sha((stackC, topf) ↦ Intg (n + 1 + 1),
  (heapC, k0f) ↦ the (sha (heapC, kf)))))) in exI)

apply (rule-tac x=nat (the-Intg (the ((sha((stackC, topf) ↦ Intg (n + 1 + 1),
  (heapC, k0f) ↦ the (sha (heapC, kf)))))) in exI)

apply (rule-tac x=la in exI)
apply (rule-tac x=ty in exI)
apply (rule-tac x=m in exI)
apply (rule-tac x=S'a(nat (n + 1) ↦ Intg (int (the (ctm' p))),
  nat (2 + n) ↦ the (sha (heapC, k0f)))) in exI)
apply (rule-tac x=(2 + n) in exI)
apply (rule-tac x=l' in exI)
apply (rule-tac x=regS in exI)
apply (rule-tac x=l'' in exI)
apply (rule-tac x=m' in exI)
apply (rule-tac x=m'' in exI)
apply (rule-tac x=d in exI)
apply (rule-tac x=g in exI)

apply (rule conjI, drule-tac t=S2 in sym, assumption)
apply (rule conjI, rule refl)
apply (rule conjI, rule refl)
apply (rule conjI, rule refl)
apply (rule conjI) apply clarsimp
  apply (simp add: Sf-def add: topf-def)
  apply (simp add: stackC-def add: heapC-def)
apply (rule conjI, assumption)
apply (rule conjI) apply clarsimp
  apply (simp add: stackC-def add: heapC-def)
  apply (simp add: kf-def add: k0f-def)
apply (rule conjI) apply (erule activeCells)

```

```

apply (rule conjI) apply (erule activeCells-2,assumption)
      apply (erule activeCells-2,assumption)
apply (rule conjI, simp)
apply (rule conjI) apply clarsimp
      apply (simp add: stackC-def add: heapC-def)
apply (rule conjI) apply clarsimp
      apply (simp add: stackC-def add: heapC-def)
      apply (simp add: regionsf-def add: k0f-def)
apply (rule conjI) apply clarsimp
apply (rule conjI) apply clarsimp
      apply (simp add: dirCellC-def add: stackC-def)
      apply (simp add: heapC-def add: safeDirf-def)
apply (rule conjI) apply clarsimp
apply (rule conjI) apply clarsimp
apply (rule conjI) apply clarsimp
apply (rule conjI) apply (rule equivH-PUSHCONT,assumption+)
apply (rule conjI) apply (rule maxPush-PUSHCONT)
apply (rule conjI) apply (simp, rule equivS-PUSHCONT-S2,simp)
apply (rule conjI) apply (unfold equivH.simps,simp)
by simp

```

**end**

```

theory dem-POPCONT
imports ../JVMSAFE/JVMEExec SVM2JVM SVMSemantics CertifSVM2JVM
begin

```

```

no-translations Norm s == (None,s)

```

```

no-translations ex-table-of m == snd (snd (snd m))

```

```

declare trInstr.simps [simp del]
declare equivS.simps [simp del]

```

```

axioms maxPush-POPCONT:
  2 + n < int m

```

**lemma** *equivS-POPCONT*:  
 $equivS (v1 \# Cont (k1, l1) \# S) S' n ctm d g \implies n \geq 2$   
**apply** (*induct S*)  
**apply** (*case-tac v1,simp-all*)  
**apply** (*unfold equivS.simps*)  
**apply** (*case-tac (nat n), simp, simp*)  
**apply** (*case-tac nata, simp, simp*)  
**apply** (*case-tac x,simp*)  
**apply** (*simp add: equivS.simps*)  
**apply** (*case-tac v1,simp-all*)  
**apply** (*simp add: equivS.simps*)  
**apply** (*simp add: equivS.simps*)  
**apply** (*case-tac x*)  
**by** (*simp add: equivS.simps*)

**lemma** *equivS2-POPCONT*:  
 $\llbracket 2 < length\ vs; equivS (Val\ v1 \# Cont (k1, l1) \# S') S'a\ n\ ctm\ d\ g \rrbracket$   
 $\implies equivS (Val\ v1 \# S') (S'a(nat\ (n - 2) \mapsto vs[Suc\ 0 := the\ (S'a\ (nat\ n))]) !$   
 $Suc\ 0)) (n - 2) ctm\ d\ g$   
**apply** (*simp add: equivS.simps, clarsimp*)  
**apply** (*drule equivS-ge-n*)  
**apply** (*erule-tac x=1 in allE,simp*)  
**by** (*subgoal-tac -2 + n = n - 2,simp,simp*)

**lemma** *equivS-k1*:  
 $\llbracket equivS (Val\ v1 \# Cont (k1, l1) \# S') S'a\ n\ ctm\ d\ g \rrbracket$   
 $\implies k1 = nat (the-Intg (the (S'a (nat (n - 1))))))$   
**by** (*simp add: equivS.simps*)

**declare** *equivH.simps* [*simp del*]

**lemma** *activeCells-POPCONT*:  
 $\llbracket la \notin activeCells\ regS\ h\ (nat\ (the-Intg\ (the\ (sha\ (heapC,\ kf)))))) \rrbracket$   
 $\implies activeCells\ regS\ h\ (nat\ (the-Intg\ (the\ (sha\ (heapC,\ kf)))) =$   
 $activeCells\ regS\ (h(la \mapsto Arr\ ty\ m\ (S'a(nat\ (n - 2) \mapsto the\ (S'a\ (nat\ n))))))$   
 $(nat\ (the-Intg\ (the\ (sha\ (heapC,\ kf))))))$   
**apply** (*frule l-not-in-cellReg*)  
**apply** (*unfold activeCells-def,auto*)  
**apply** (*unfold region-def, simp add: Let-def, elim conjE*)  
**apply** (*rule-tac x=j in exI,simp*)  
**apply** (*rule cellReg-step*)  
**apply** (*rule cellReg-basic*)  
**apply** (*erule-tac x=j in allE*)  
**apply** (*rule cellsReg-monotone-1,assumption+*)  
**apply** (*erule l-not-in-regs*)  
**apply** (*simp add: Let-def, elim conjE*)

**apply** (*rule-tac*  $x=j$  **in**  $exI$ , *simp*)  
**apply** (*rule* *cellReg-step*)  
**apply** (*rule* *cellReg-basic*)  
**apply** (*erule-tac*  $x=j$  **in**  $allE$ )  
**apply** (*erule* *cellsReg-monotone-2,assumption+*)  
**by** (*erule* *l-not-in-regs*)

**lemma** *domH-POPCONT*:

$\llbracket la \notin activeCells\ regS\ h\ (nat\ (the-Intg\ (the\ (sha\ (heapC,\ kf)))));$   
 $\forall l \in dom\ H.\ \exists l'.\ l' = the\ (g\ l) \wedge equivC\ (the\ (H\ l))\ h\ l'\ (the\ (h\ l'))\ (nat$   
 $(the-Intg\ (the\ (sha\ (heapC,\ kf))))\ com\ regS\ d\ g \rrbracket$   
 $\implies \forall l \in dom\ H.$   
 $\exists l'.\ l' = the\ (g\ l) \wedge$   
 $equivC\ (the\ (H\ l))\ (h(la \mapsto Arr\ ty\ m\ (S'a(nat\ (n - 2) \mapsto the\ (S'a$   
 $(nat\ n))))))\ l'$   
 $(the\ ((h(la \mapsto Arr\ ty\ m\ (S'a(nat\ (n - 2) \mapsto the\ (S'a\ (nat\ n))))))$   
 $l'))$   
 $(nat\ (the-Intg\ (the\ (sha\ (heapC,\ kf))))))$   
 $com\ regS\ d\ g$

**apply** (*erule* *l-not-in-cellReg*)  
**apply** (*rule* *ballI*)  
**apply** (*erule-tac*  $x=l$  **in**  $ballE$ )  
**prefer** 2 **apply** *simp*  
**apply** (*erule*  $exE$ )  
**apply** (*rule-tac*  $x=l'$  **in**  $exI$ ) **apply** (*elim* *conjE*)  
**apply** (*rule* *conjI*, *assumption*)  
**apply** *clarsimp*  
**apply** (*rule* *conjI*, *rule* *impI*)  
**apply** (*simp* *add: region-def*)  
**apply** (*rule* *impI*)  
**apply** (*rule-tac*  $x=Obj$  **in**  $exI$ )  
**apply** (*rule-tac*  $x=flds$  **in**  $exI$ ) **apply** *simp*  
**apply** (*simp* *add: region-def*)  
**apply** *clarsimp*  
**apply** (*rule* *cellReg-step*)  
**apply** (*rule* *cellReg-basic*)  
**apply** (*erule-tac*  $x=j$  **in**  $allE$ )  
**apply** (*rule* *cellsReg-monotone-1,assumption+*)  
**by** (*erule* *l-not-in-regs*)

**lemma** *equivH-POPCONT*:

$\llbracket ((hm,\ k),\ k0,\ (l,\ i),\ v1 \# Cont\ (k1,\ l1) \# S') = ((H,\ ka),\ k0a,\ PC,\ S);$   
 $k' = nat\ (the-Intg\ (the\ (sha\ (heapC,\ kf))));$   $k0' = nat\ (the-Intg\ (the\ (sha$   
 $(heapC,\ k0f))));$   $sha\ (stackC,\ Sf) = Some\ (Addr\ la);$



$distinct [la, l', l'']; h\ la = Some (Arr\ ty\ m\ S'a); sha\ (stackC, topf) = Some$   
 $(Intg\ n);$   
 $sha\ (heapC, regionsf) = Some (Addr\ l'); h\ l' = Some (Arr\ ty\ m'\ regS); sha$   
 $(dirCellC, safeDirf) = Some (Addr\ l'');$   
 $h\ l'' = Some (Arr\ ty\ m''\ d); inj-on\ g\ (dom\ H); equivH\ (H, ka)\ h\ k'\ com\ regS$   
 $d\ g; k0a = k0';$   
 $activeCells\ regS\ h\ k' \cap \{la, l', l''\} = \{\}; 2 < length\ vs]$   
 $\implies equivH\ (hm, k)\ (h(la \mapsto Arr\ ty\ m\ (S'a(nat\ (n - 2)) \mapsto vs[Suc\ 0 := the$   
 $(S'a\ (nat\ n))]) ! Suc\ 0))$   
 $(nat\ (the-Intg$   
 $(the\ ((sha((stackC, topf) \mapsto Intg\ (n - 2), (heapC, k0f) \mapsto$   
 $vs[Suc\ 0 := the\ (S'a\ (nat\ n)), 2 := the\ (S'a\ (nat\ (n - 1)))] !$   
 $2))$   
 $(heapC, kf))))))$   
 $com\ regS\ d\ g$   
**apply**  $(simp\ add: heapC-def\ add: stackC-def)$   
**apply**  $(simp\ add: kf-def\ add: k0f-def)$   
**apply**  $(fold\ heapC-def, fold\ kf-def)$   
**apply**  $(unfold\ equivH.simps, elim\ conjE)$   
**apply**  $(rule\ conjI, assumption)$   
**apply**  $(rule\ conjI, frule\ activeCells-POPCONT, simp)$   
**apply**  $(rule\ conjI, assumption)$   
**by**  $(rule\ domH-POPCONT, assumption+)$

**declare**  $extractBytecode-def [simp\ del]$   
**declare**  $initConsTable-def [simp\ del]$   
**declare**  $equivS.simps [simp\ del]$

**axioms**  $Tableswitch-POPCONT:$

$drop\ pcc\ (extractBytecode\ P') =$   
 $[Tableswitch\ 1\ (int\ (length\ ps))\ ps] @ bytecode' \wedge ctm\ (nat\ (ps!i)) = Some\ i$

**axioms**  $equivpc-POPCONT:$

$if'\ (the-Intg\ (vs[Suc\ 0 := the\ (S'a\ (nat\ n)),$   
 $2 := the\ (S'a\ (nat\ (n - 1))),$   
 $3 := the\ (S'a\ (nat\ (n - 2)))] ! 3) < 1 \vee$   
 $int\ (length\ ps) < the-Intg\ (vs[Suc\ 0 := the\ (S'a\ (nat\ n),$   
 $2 := the\ (S'a\ (nat\ (n - 1))),$   
 $3 := the\ (S'a\ (nat\ (n - 2)))] ! 3))$   
 $(nat\ (24 + int\ (the\ (cdm\ (l, i))) + trAddr\ pcc\ (the\ (cdm\ (l, i)) + incPop)) +$   
 $nat\ (ps ! nat\ (int\ (length\ ps))))$   
 $(nat\ (24 + int\ (the\ (cdm\ (l, i))) + trAddr\ pcc\ (the\ (cdm\ (l, i)) + incPop)) +$   
 $nat\ (ps ! nat\ (the-Intg\ (vs[Suc\ 0 := the\ (S'a\ (nat\ n),$   
 $2 := the\ (S'a\ (nat\ (n - 1))),$   
 $3 := the\ (S'a\ (nat\ (n - 2)))] ! 3) - 1))) =$   
 $the\ (cdm\ (l1, 0))$

```

constdefs if' :: [bool, 'a', 'a'] => 'a'
  if' c a b ≡ if c then a else b

declare if'-def [simp del]

lemma execSVMInstr-POPCONT :
  [| (P', cdm, ctm, com) = trSVM2JVM ((svms, ctmap), ini, ct, ah, ai, bc);
   cdm , ctm, com ⊢ ((hm, k), k0, (l, i), S) ≐ S1';
   (fst (the (map-of svms l)) ! i) = POPCONT;
   execSVMInst POPCONT (map-of ct) (hm, k) k0 (l, i) S = Either.Right
S2;
   drop (the (cdm l i)) (extractBytecode P') =
   trInstr (the (cdm l i)) cdm' ctm' com pcc POPCONT @ bytecode'
  |] ⇒ ∃ v' sh' dh' ih' fms' .
  P' ⊢ S1' -jvm→ (v',sh',dh',ih', fms') ∧
  cdm , ctm, com ⊢ S2 ≐ (v',sh',dh',ih', fms')

apply (case-tac S1')
apply (rename-tac v tup)
apply (case-tac tup)
apply (rename-tac sh tup)
apply (case-tac tup)
apply (rename-tac dh tup)
apply (case-tac tup)
apply (rename-tac ih fms)
apply (simp)
apply (subgoal-tac ∃ v1 k1 l1 S'. S = Val v1 # Cont (k1,l1) # S', clarsimp)

prefer 2
apply (case-tac S)
apply simp
apply (insert RightNotUndefined)
apply (erule-tac x = S2 in allE, force)
apply simp
apply (case-tac a)
apply simp
apply (case-tac list)
apply simp
apply simp
apply (case-tac aa)
apply simp
apply (erule-tac x = S2 in allE, force)
apply simp
apply simp
apply simp
apply simp

```

**apply** (*unfold equivState-def*)  
**apply** (*elim exE,elim conjE*)

**apply** (*rule-tac x=None in exI*)  
**apply** (*rule-tac x=sha((stackC, topf) ↦ Intg (n - 2),*  
*(heapC, kOf) ↦ vs[Suc 0 := the (S'a (nat n)), 2 := the (S'a*  
*(nat (n - 1)))] ! 2) in exI*)  
**apply** (*rule-tac x=h(la ↦ Arr ty m (S'a(nat (n - 2) ↦ vs[Suc 0 := the (S'a*  
*(nat n)] ! Suc 0))) in exI*)  
**apply** (*rule-tac x=inih in exI*)  
**apply** (*rule-tac x=[[],*  
*vs[Suc 0 := the (S'a (nat n)), 2 := the (S'a (nat (n - 1))), 3*  
*:= the (S'a (nat (n - 2)))],*  
*safeP, sigSafeMain,*  
*if' (((the-Intg (vs[Suc 0 := the (S'a (nat n)),*  
*2 := the (S'a (nat (n - 1))),*  
*3 := the (S'a (nat (n - 2)))] ! 3) < 1) ∨*  
*(int (length ps) < the-Intg (vs[Suc 0 := the (S'a (nat n)),*  
*2 := the (S'a (nat (n - 1))),*  
*3 := the (S'a (nat (n - 2)))] ! 3))))*  
*(nat (24 + int (the (cdm (l, i))) + trAddr pcc (the (cdm (l, i))*  
*+ incPop)) +*  
*nat (ps ! nat (int (length ps))))*  
*(nat (24 + int (the (cdm (l, i))) + trAddr pcc (the (cdm (l, i))*  
*+ incPop)) +*  
*nat (ps ! nat (the-Intg (vs[Suc 0 := the (S'a (nat n)),*  
*2 := the (S'a (nat (n - 1))),*  
*3 := the (S'a (nat (n - 2)))] !*  
*3) - 1))),*  
*ref)] in exI*)

**apply** (*subgoal-tac ∃ vs pc . fms=([],vs,safeP,sigSafeMain,pc,ref)#[]*)  
**prefer 2 apply simp**  
**apply** (*erule exE*)+  
**apply** (*unfold exec-all-def*)  
**apply** (*rule conjI*)

**apply** (*subgoal-tac PC = (l,i)*) **prefer 2 apply simp**  
**apply simp apply** (*elim conjE*) **apply clarsimp**  
**apply** (*subgoal-tac P<sup>↑</sup> (None,sha,h,inih,[[], vs, safeP, sigSafeMain, the (cdm*  
*(l,i)), aa,ba)] -jvm→*  
*(None,sha,h,inih,[[(the (sha (stackC,Sf))], vs, safeP, sigSafeMain,*  
*the (cdm (l,i)) + 1, aa,ba)]])*)  
**apply** (*unfold exec-all-def*)

```

apply (erule rtrancl-trans)

prefer 2
apply (rule r-into-rtrancl)
apply (clarify)
apply (unfold JVMExec.exec.simps)
apply (unfold Let-def)
apply (unfold extractBytecode-def)
apply (subgoal-tac (fst (snd (snd (snd (snd (the (method' (P', safeP) sigSafeMain)))))) !
      the (cdm (l, i))) = Getstatic Sf stackC,simp)
apply (unfold trInstr.simps, erule nth-via-drop-append)

apply (drule drop-Suc-append)
apply (subgoal-tac  $P \vdash$  (None, sha, h, inih,
       $[[[the\ (sha\ (stackC, Sf))],$ 
       $vs, safeP, sigSafeMain, the\ (cdm\ (l, i)) + 1, aa, ba]]$ )
       $-jvm \rightarrow$ 
      (None, sha, h, inih,
       $[[[the\ (sha\ (stackC, topf)), the\ (sha\ (stackC, Sf))],$ 
       $vs, safeP, sigSafeMain, the\ (cdm\ (l, i)) + 1 + 1, aa, ba]]$ ))

apply (unfold exec-all-def)
apply (erule rtrancl-trans)
prefer 2
apply (rule r-into-rtrancl)
apply (clarify)
apply (unfold JVMExec.exec.simps)
apply (unfold Let-def)
apply (subgoal-tac (fst (snd (snd (snd (snd (the (method' (P', safeP) sigSafeMain)))))) ! Suc (the (cdm (l, i)))) =
      Getstatic topf stackC,simp)
apply (erule nth-via-drop-append)

apply (drule drop-Suc-append)
apply (subgoal-tac  $P \vdash$  (None, sha, h, inih,
       $[[[the\ (sha\ (stackC, topf)), the\ (sha\ (stackC, Sf))],$ 
       $vs, safeP, sigSafeMain, the\ (cdm\ (l, i)) + 1 + 1, aa, ba]]$ )
       $-jvm \rightarrow$ 
      (None, sha, h, inih,
       $[[[the\ (sha\ (stackC, topf)), the\ (sha\ (stackC, Sf)), the\ (sha$ 
      (stackC, topf)), the\ (sha\ (stackC, Sf))],
       $vs, safeP, sigSafeMain, the\ (cdm\ (l, i)) + 1 + 1 + 1, aa,$ 
      ba]]))
apply (unfold exec-all-def)
apply (erule rtrancl-trans)
prefer 2

```

**apply** (*rule r-into-rtrancl*)  
**apply** (*clarify*)  
**apply** (*unfold JVMEexec.exec.simps*)  
**apply** (*unfold Let-def*)  
**apply** (*subgoal-tac* (*fst* (*snd* (*snd* (*snd* (*snd* (*the* (*method'* (*P'*, *safeP*) *sigSafeMain*)))))) ! (*Suc* (*Suc* (*the* (*cdm* (*l*, *i*)))))) =  
*Dup2,simp*)  
**apply** (*erule nth-via-drop-append*)

**apply** (*drule drop-Suc-append*)  
**apply** (*subgoal-tac*  $P \vdash$  (*None,sha,h,inih*,  
[[*the* (*sha* (*stackC,topf*)),*the* (*sha* (*stackC, Sf*)),  
*the* (*sha* (*stackC,topf*)),*the* (*sha* (*stackC,Sf*))],  
*vs, safeP, sigSafeMain, the* (*cdm* (*l, i*)) + 1 + 1 + 1, *aa, ba*]])  
 $\rightarrow$   
(*None,sha,h,inih*,  
[[*the* (*sha* (*stackC,topf*)),*the* (*sha* (*stackC, Sf*)),  
*the* (*sha* (*stackC,topf*)),*the* (*sha* (*stackC, Sf*)),  
*the* (*sha* (*stackC,topf*)),*the* (*sha* (*stackC,Sf*))],  
*vs, safeP, sigSafeMain, the* (*cdm* (*l, i*)) + 1 + 1 + 1 + 1,  
*aa, ba*]))  
**apply** (*unfold exec-all-def*)  
**apply** (*erule rtrancl-trans*)  
**prefer** 2  
**apply** (*rule r-into-rtrancl*)  
**apply** (*clarify*)  
**apply** (*unfold JVMEexec.exec.simps*)  
**apply** (*unfold Let-def*)  
**apply** (*subgoal-tac* (*fst* (*snd* (*snd* (*snd* (*snd* (*the* (*method'* (*P'*, *safeP*) *sigSafeMain*)))))) ! (*Suc* (*Suc* (*Suc* (*the* (*cdm* (*l, i*)))))) =  
*Dup2,simp*)  
**apply** (*erule nth-via-drop-append*)

**apply** (*drule drop-Suc-append*)  
**apply** (*subgoal-tac*  $P \vdash$  (*None,sha,h,inih*,  
[[*the* (*sha* (*stackC,topf*)),*the* (*sha* (*stackC, Sf*)),  
*the* (*sha* (*stackC,topf*)),*the* (*sha* (*stackC, Sf*)),  
*the* (*sha* (*stackC,topf*)),*the* (*sha* (*stackC,Sf*))],  
*vs, safeP, sigSafeMain, the* (*cdm* (*l, i*)) + 1 + 1 + 1 + 1,  
*aa, ba*]])  
 $\rightarrow$   
(*None,sha,h,inih*,  
[[*the* (*sha* (*stackC,topf*)),*the* (*sha* (*stackC, Sf*)),  
*the* (*sha* (*stackC,topf*)),*the* (*sha* (*stackC, Sf*)),  
*the* (*sha* (*stackC,topf*)),*the* (*sha* (*stackC, Sf*))],

```

the (sha (stackC,topf)),the (sha (stackC, Sf)),
the (sha (stackC,topf)),the (sha (stackC,Sf))),
vs, safeP, sigSafeMain, the (cdm (l, i)) + 1 + 1 + 1 + 1 +
1, aa, ba]]))
apply (unfold exec-all-def)
apply (erule rtrancl-trans)
prefer 2
apply (rule r-into-rtrancl)
apply (clarify)
apply (unfold JVMExec.exec.simps)
apply (unfold Let-def)
apply (subgoal-tac (fst (snd (snd (snd (snd (the (method' (P', safeP) sigSafe-
Main))))))) !
(Suc (Suc (Suc (Suc (the (cdm (l, i))))))) =
Dup2,simp)
apply (erule nth-via-drop-append)

apply (drule drop-Suc-append)
apply (subgoal-tac P+ (None,sha,h,inih,
[[the (sha (stackC,topf)),the (sha (stackC, Sf)),
the (sha (stackC,topf)),the (sha (stackC, Sf)),
the (sha (stackC,topf)),the (sha (stackC, Sf)),
the (sha (stackC,topf)),the (sha (stackC,Sf))),
vs, safeP, sigSafeMain, the (cdm (l, i)) + 1 + 1 + 1 + 1 +
1, aa, ba]]))
-jvm→
(None,sha,h,inih,
[[the (S'a (nat n)),
the (sha (stackC,topf)),the (sha (stackC, Sf)),
the (sha (stackC,topf)),the (sha (stackC, Sf)),
the (sha (stackC,topf)),the (sha (stackC, Sf))),
vs, safeP, sigSafeMain, the (cdm (l, i)) + 1 + 1 + 1 + 1 +
1 + 1, aa, ba]]))
apply (unfold exec-all-def)
apply (erule rtrancl-trans)
prefer 2
apply (rule r-into-rtrancl)
apply (clarify)
apply (unfold JVMExec.exec.simps)
apply (unfold Let-def)
apply (subgoal-tac (fst (snd (snd (snd (snd (the (method' (P', safeP) sigSafe-
Main))))))) !
(Suc (Suc (Suc (Suc (Suc (the (cdm (l, i))))))) =
ArrLoad,simp)

apply (frule equivS-POPCONT, simp)
apply (simp add: raise-system-xcpt-def)
apply (erule nth-via-drop-append)

```

```

apply (drule drop-Suc-append)
apply (subgoal-tac  $P \vdash$  (None, sha, h, inih,
  [[(the (S'a (nat n)),
    the (sha (stackC, topf)), the (sha (stackC, Sf)),
    the (sha (stackC, topf)), the (sha (stackC, Sf)),
    the (sha (stackC, topf)), the (sha (stackC, Sf))],
    vs, safeP, sigSafeMain, the (cdm (l, i)) + 1 + 1 + 1 + 1 +
1 + 1, aa, ba]])
   $\text{-jvm}\rightarrow$ 
  (None, sha, h, inih,
  [[(the (sha (stackC, topf)), the (sha (stackC, Sf)),
    the (sha (stackC, topf)), the (sha (stackC, Sf)),
    the (sha (stackC, topf)), the (sha (stackC, Sf))],
    vs[Suc 0 := the (S'a (nat n))],
    safeP, sigSafeMain, the (cdm (l, i)) + 1 + 1 + 1 + 1 + 1
+ 1 + 1, aa, ba]])
apply (unfold exec-all-def)
apply (erule rtrancl-trans)
prefer 2
apply (rule r-into-rtrancl)
apply (clarify)
apply (unfold JVMExec.exec.simps)
apply (unfold Let-def)
apply (subgoal-tac (fst (snd (snd (snd (snd (the (method' (P', safeP) sigSafe-
Main)))))) !
  (Suc (Suc (Suc (Suc (Suc (Suc (the (cdm (l, i)))))))))) =
  Store 1, simp)
apply (erule nth-via-drop-append)

```

```

apply (drule drop-Suc-append)
apply (subgoal-tac  $P \vdash$  (None, sha, h, inih,
  [[(the (sha (stackC, topf)), the (sha (stackC, Sf)),
    the (sha (stackC, topf)), the (sha (stackC, Sf)),
    the (sha (stackC, topf)), the (sha (stackC, Sf))],
    vs[Suc 0 := the (S'a (nat n))],
    safeP, sigSafeMain, the (cdm (l, i)) + 1 + 1 + 1 + 1 + 1
+ 1 + 1, aa, ba]])
   $\text{-jvm}\rightarrow$ 
  (None, sha, h, inih,
  [[(Intg 1,
    the (sha (stackC, topf)), the (sha (stackC, Sf)),
    the (sha (stackC, topf)), the (sha (stackC, Sf)),
    the (sha (stackC, topf)), the (sha (stackC, Sf))],
    vs[Suc 0 := the (S'a (nat n))],
    safeP, sigSafeMain, the (cdm (l, i)) + 1 + 1 + 1 + 1 + 1

```

```

+ 1 + 1 + 1, aa, ba]))
apply (unfold exec-all-def)
apply (erule rtrancl-trans)
prefer 2
  apply (rule r-into-rtrancl)
  apply (clarify)
  apply (unfold JVMExec.exec.simps)
  apply (unfold Let-def)
  apply (subgoal-tac (fst (snd (snd (snd (snd (the (method' (P', safeP) sigSafe-
Main))))))) !
      (Suc (Suc (Suc (Suc (Suc (Suc (Suc (the (cdm (l, i)))))))))) =
      LitPush (Intg 1),simp)
apply (erule nth-via-drop-append)

apply (drule drop-Suc-append)
apply (subgoal-tac P+ (None, sha, h, inih,
  [[Intg 1,
   the (sha (stackC, topf)), the (sha (stackC, Sf)),
   the (sha (stackC, topf)), the (sha (stackC, Sf)),
   the (sha (stackC, topf)), the (sha (stackC, Sf)),
   vs[Suc 0 := the (S'a (nat n))],
   safeP, sigSafeMain, the (cdm (l, i)) + 1 + 1 + 1 + 1 + 1
+ 1 + 1 + 1, aa, ba]))
  -jvm→
  (None, sha, h, inih,
  [[Intg (the-Intg (the (sha (stackC, topf))) - 1),
   the (sha (stackC, Sf)), the (sha (stackC, topf)),
   the (sha (stackC, Sf)), the (sha (stackC, topf)),
   the (sha (stackC, Sf)),
   vs[Suc 0 := the (S'a (nat n))],
   safeP, sigSafeMain, the (cdm (l, i)) + 1 + 1 + 1 + 1 + 1
+ 1 + 1 + 1 + 1, aa, ba]))
apply (unfold exec-all-def)
apply (erule rtrancl-trans)
prefer 2
  apply (rule r-into-rtrancl)
  apply (clarify)
  apply (unfold JVMExec.exec.simps)
  apply (unfold Let-def)
  apply (subgoal-tac (fst (snd (snd (snd (snd (the (method' (P', safeP) sigSafe-
Main))))))) !
      (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (the (cdm (l, i))))))))))
=
      BinOp Subtract, simp)
apply (erule nth-via-drop-append)

```



```

apply (drule drop-Suc-append)
apply (subgoal-tac  $P \vdash$  (None,sha,h,inih,
  [[(Intg (the-Intg (the (sha (stackC, topf))) - 1),
    the (sha (stackC, Sf)),the (sha (stackC,topf)),
    the (sha (stackC, Sf)),the (sha (stackC,topf)),
    the (sha (stackC,Sf))),
  vs[Suc 0 := the (S'a (nat n))],
  safeP, sigSafeMain, the (cdm (l, i)) + 1 + 1 + 1 + 1 + 1
+ 1 + 1 + 1 + 1, aa, ba]])
  -jvm→
  (None,sha,h,inih,
  [[(the (S'a (nat (n - 1))),
    the (sha (stackC,topf)),the (sha (stackC, Sf)),
    the (sha (stackC,topf)),the (sha (stackC,Sf))),
  vs[Suc 0 := the (S'a (nat n))],
  safeP, sigSafeMain, the (cdm (l, i)) + 1 + 1 + 1 + 1 + 1
+ 1 + 1 + 1 + 1 + 1, aa, ba]]))
apply (unfold exec-all-def)
apply (erule rtrancl-trans)
prefer 2
apply (rule r-into-rtrancl)
apply (clarify)
apply (unfold JVMExec.exec.simps)
apply (unfold Let-def)
apply (subgoal-tac (fst (snd (snd (snd (snd (the (method' (P', safeP) sigSafe-
Main)))))) !
  (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (the (cdm (l,
i)))))))))))))) =
  ArrLoad,simp)

apply (frule equivS-POPCONT, simp)
apply (simp add: raise-system-xcpt-def)
apply (erule nth-via-drop-append)

```

```

apply (drule drop-Suc-append)
apply (subgoal-tac  $P \vdash$  (None,sha,h,inih,
  [[(the (S'a (nat (n - 1))),
    the (sha (stackC,topf)),the (sha (stackC, Sf)),
    the (sha (stackC,topf)),the (sha (stackC,Sf))),
  vs[Suc 0 := the (S'a (nat n))],
  safeP, sigSafeMain, the (cdm (l, i)) + 1 + 1 + 1 + 1 + 1
+ 1 + 1 + 1 + 1 + 1, aa, ba]])
  -jvm→
  (None,sha,h,inih,
  [[(the (sha (stackC,topf)),the (sha (stackC, Sf)),
    the (sha (stackC,topf)),the (sha (stackC,Sf))),
  vs[Suc 0 := the (S'a (nat n))],

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                2 := the (S'a (nat (n - 1))),
                safeP, sigSafeMain, the (cdm (l, i)) + 1 + 1 + 1 + 1 + 1
+ 1 + 1 + 1 + 1 + 1 + 1 + 1, aa, ba]))
apply (unfold exec-all-def)
apply (erule rtrancl-trans)
prefer 2
apply (rule r-into-rtrancl)
apply (clarify)
apply (unfold JVMExec.exec.simps)
apply (unfold Let-def)
apply (subgoal-tac (fst (snd (snd (snd (snd (the (method' (P', safeP) sigSafe-
Main))))))) !
(Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (the (cdm
(l, i)))))))))))))) =
Store 2, simp)
apply (erule nth-via-drop-append)

apply (drule drop-Suc-append)
apply (subgoal-tac P ⊢ (None, sha, h, inih,
[[[the (sha (stackC, topf)), the (sha (stackC, Sf)),
the (sha (stackC, topf)), the (sha (stackC, Sf))],
vs[Suc 0 := the (S'a (nat n)),
2 := the (S'a (nat (n - 1))),
safeP, sigSafeMain, the (cdm (l, i)) + 1 + 1 + 1 + 1 + 1
+ 1 + 1 + 1 + 1 + 1 + 1 + 1, aa, ba]])
-jvm→
(None, sha, h, inih,
[[[Intg 2,
the (sha (stackC, topf)), the (sha (stackC, Sf)),
the (sha (stackC, topf)), the (sha (stackC, Sf))],
vs[Suc 0 := the (S'a (nat n)),
2 := the (S'a (nat (n - 1))),
safeP, sigSafeMain, the (cdm (l, i)) + 1 + 1 + 1 + 1 + 1
+ 1 + 1 + 1 + 1 + 1 + 1 + 1, aa, ba]]))
apply (unfold exec-all-def)
apply (erule rtrancl-trans)
prefer 2
apply (rule r-into-rtrancl)
apply (clarify)
apply (unfold JVMExec.exec.simps)
apply (unfold Let-def)
apply (subgoal-tac (fst (snd (snd (snd (snd (the (method' (P', safeP) sigSafe-
Main))))))) !
(Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (the
(cdm (l, i)))))))))))))) =
LitPush (Intg 2), simp)
apply (erule nth-via-drop-append)

```

```

apply (drule drop-Suc-append)
apply (subgoal-tac P ⊢ (None, sha, h, inih,
  [[Intg 2,
    the (sha (stackC, topf)), the (sha (stackC, Sf)),
    the (sha (stackC, topf)), the (sha (stackC, Sf))]],
  vs[Suc 0 := the (S'a (nat n)),
    2 := the (S'a (nat (n - 1)))],
  safeP, sigSafeMain, the (cdm (l, i)) + 1 + 1 + 1 + 1 + 1
+ 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1, aa, ba]))
  -jvm→
  (None, sha, h, inih,
  [[Intg (the-Intg (the (sha (stackC, topf))) - 2),
    the (sha (stackC, Sf)), the (sha (stackC, topf)),
    the (sha (stackC, Sf))]],
  vs[Suc 0 := the (S'a (nat n)),
    2 := the (S'a (nat (n - 1)))],
  safeP, sigSafeMain, the (cdm (l, i)) + 1 + 1 + 1 + 1 + 1
+ 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1, aa, ba]))
apply (unfold exec-all-def)
apply (erule rtrancl-trans)
prefer 2
apply (rule r-into-rtrancl)
apply (clarify)
apply (unfold JVMExec.exec.simps)
apply (unfold Let-def)
apply (subgoal-tac (fst (snd (snd (snd (snd (the (method' (P', safeP) sigSafe-
Main)))))) !
  (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc
(the (cdm (l, i)))))))))))))) =
  BinOp Subtract, simp)
apply (erule nth-via-drop-append)

```

```

apply (drule drop-Suc-append)
apply (subgoal-tac P ⊢ (None, sha, h, inih,
  [[Intg (the-Intg (the (sha (stackC, topf))) - 2),
    the (sha (stackC, Sf)), the (sha (stackC, topf)),
    the (sha (stackC, Sf))]],
  vs[Suc 0 := the (S'a (nat n)),
    2 := the (S'a (nat (n - 1)))],
  safeP, sigSafeMain, the (cdm (l, i)) + 1 + 1 + 1 + 1 + 1
+ 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1, aa, ba]))
  -jvm→
  (None, sha, h, inih,
  [[the (S'a (nat (n - 2))),
    the (sha (stackC, topf)), the (sha (stackC, Sf))]],

```

```

vs[Suc 0 := the (S'a (nat n)),
   2 := the (S'a (nat (n - 1)))],
safeP, sigSafeMain, the (cdm (l, i) + 1 + 1 + 1 + 1 + 1
+ 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1, aa, ba)])
apply (unfold exec-all-def)
apply (erule rtrancl-trans)
prefer 2
apply (rule r-into-rtrancl)
apply (clarify)
apply (unfold JVMEExec.exec.simps)
apply (unfold Let-def)
apply (subgoal-tac (fst (snd (snd (snd (snd (the (method' (P', safeP) sigSafe-
Main))))))) !
(Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc
(Suc (the (cdm (l, i))))))))))))))))) =
ArrLoad,simp)

apply (frule equivS-POPCONT, simp)
apply (simp add: raise-system-xcpt-def)
apply (erule nth-via-drop-append)

apply (drule drop-Suc-append)
apply (subgoal-tac P+ (None, sha, h, inih,
  [([the (S'a (nat (n - 2))),
    the (sha (stackC, topf)), the (sha (stackC, Sf))],
vs[Suc 0 := the (S'a (nat n)),
   2 := the (S'a (nat (n - 1)))],
safeP, sigSafeMain, the (cdm (l, i) + 1 + 1 + 1 + 1 + 1
+ 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1, aa, ba)])
-jvm→
  (None, sha, h, inih,
  [([the (sha (stackC, topf)), the (sha (stackC, Sf))],
vs[Suc 0 := the (S'a (nat n)),
   2 := the (S'a (nat (n - 1))),
   3 := the (S'a (nat (n - 2)))],
safeP, sigSafeMain, the (cdm (l, i) + 1 + 1 + 1 + 1 + 1
+ 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1, aa, ba)]))
apply (unfold exec-all-def)
apply (erule rtrancl-trans)
prefer 2
apply (rule r-into-rtrancl)
apply (clarify)
apply (unfold JVMEExec.exec.simps)
apply (unfold Let-def)
apply (subgoal-tac (fst (snd (snd (snd (snd (the (method' (P', safeP) sigSafe-
Main))))))) !
(Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc
(Suc (Suc (the (cdm (l, i))))))))))))))))) =

```

*Store 3, simp*)  
**apply** (*erule nth-via-drop-append*)

**apply** (*drule drop-Suc-append*)  
**apply** (*subgoal-tac P* ⊢ (*None, sha, h, inih,*  
[[*the (sha (stackC, topf)), the (sha (stackC, Sf))*]],  
*vs*[*Suc 0 := the (S'a (nat n)),*  
*2 := the (S'a (nat (n - 1))),*  
*3 := the (S'a (nat (n - 2)))*]],  
*safeP, sigSafeMain, the (cdm (l, i)) + 1 + 1 + 1 + 1 + 1*  
+ *1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1, aa, ba*]))  
*-jvm* →  
(*None, sha, h, inih,*  
[[*Intg 2,*  
*the (sha (stackC, topf)), the (sha (stackC, Sf))*]],  
*vs*[*Suc 0 := the (S'a (nat n)),*  
*2 := the (S'a (nat (n - 1))),*  
*3 := the (S'a (nat (n - 2)))*]],  
*safeP, sigSafeMain, the (cdm (l, i)) + 1 + 1 + 1 + 1 + 1*  
+ *1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1, aa, ba*]))))  
**apply** (*unfold exec-all-def*)  
**apply** (*erule rtrancl-trans*)  
**prefer** 2  
**apply** (*rule r-into-rtrancl*)  
**apply** (*clarify*)  
**apply** (*unfold JVMExec.exec.simps*)  
**apply** (*unfold Let-def*)  
**apply** (*subgoal-tac (fst (snd (snd (snd (snd (the (method' (P', safeP) sigSafe-*  
*Main))))))) !*  
(*Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc*  
(*Suc (Suc (Suc (the (cdm (l, i)))))))))))))))))) =*  
*LitPush (Intg 2), simp*)  
**apply** (*erule nth-via-drop-append*)

**apply** (*drule drop-Suc-append*)  
**apply** (*subgoal-tac P* ⊢ (*None, sha, h, inih,*  
[[*Intg 2,*  
*the (sha (stackC, topf)), the (sha (stackC, Sf))*]],  
*vs*[*Suc 0 := the (S'a (nat n)),*  
*2 := the (S'a (nat (n - 1))),*  
*3 := the (S'a (nat (n - 2)))*]],  
*safeP, sigSafeMain, the (cdm (l, i)) + 1 + 1 + 1 + 1 + 1*  
+ *1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1, aa, ba*]))  
*-jvm* →  
(*None, sha, h, inih,*

```

      [[Intg (n - 2),
        the (sha (stackC,Sf))],
      vs[Suc 0 := the (S'a (nat n)),
          2 := the (S'a (nat (n - 1))),
          3 := the (S'a (nat (n - 2)))],
      safeP, sigSafeMain, the (cdm (l, i) + 1 + 1 + 1 + 1 + 1 +
1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1, aa, ba)])
apply (unfold exec-all-def)
apply (erule rtrancl-trans)
prefer 2
apply (rule r-into-rtrancl)
apply (clarify)
apply (unfold JVMEexec.exec.simps)
apply (unfold Let-def)
apply (subgoal-tac (fst (snd (snd (snd (snd (the (method' (P', safeP) sigSafe-
Main))))))) !
      (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc
(Suc (Suc (Suc (Suc (the (cdm (l, i)))))))))))))))))) =
      BinOp Subtract,simp)
apply (erule nth-via-drop-append)

apply (drule drop-Suc-append)
apply (subgoal-tac P+ (None,sha,h,inih,
      [[Intg (n - 2),
        the (sha (stackC,Sf))],
      vs[Suc 0 := the (S'a (nat n)),
          2 := the (S'a (nat (n - 1))),
          3 := the (S'a (nat (n - 2)))],
      safeP, sigSafeMain, the (cdm (l, i) + 1 + 1 + 1 + 1 + 1 +
1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1, aa, ba)])
      -jvm→
      (None,sha,h,inih,
      [[Intg (n - 2),
        Intg (n - 2),
        the (sha (stackC,Sf))],
      vs[Suc 0 := the (S'a (nat n)),
          2 := the (S'a (nat (n - 1))),
          3 := the (S'a (nat (n - 2)))],
      safeP, sigSafeMain,
      the (cdm (l, i) + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1, aa, ba)])
apply (unfold exec-all-def)
apply (erule rtrancl-trans)
prefer 2
apply (rule r-into-rtrancl)
apply (clarify)
apply (unfold JVMEexec.exec.simps)
apply (unfold Let-def)

```

**apply** (*subgoal-tac* (*fst* (*snd* (*snd* (*snd* (*snd* (*the* (*method'* (*P'*, *safeP*) *sigSafeMain*))))))) !  
(Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc  
(Suc (Suc (Suc (Suc (*the* (*cdm* (*l*, *i*)))))))))))))))) =  
*Dup,simp*)  
**apply** (*erule nth-via-drop-append*)

**apply** (*drule drop-Suc-append*)  
**apply** (*subgoal-tac*  $P \vdash$  (*None,sha,h,inih*,  
[[*Intg* (*n* - 2),  
*Intg* (*n* - 2),  
*the* (*sha* (*stackC,Sf*))],  
*vs*[*Suc* 0 := *the* (*S'a* (*nat* *n*)),  
2 := *the* (*S'a* (*nat* (*n* - 1))),  
3 := *the* (*S'a* (*nat* (*n* - 2)))]],  
*safeP, sigSafeMain*,  
*the* (*cdm* (*l*, *i*) + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1  
+ 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1, *aa, ba*]))  
-*jvm*→  
(*None, sha*((*stackC, topf*)  $\mapsto$  *Intg* (*n* - 2)),*h,inih*,  
[[*Intg* (*n* - 2),  
*the* (*sha* (*stackC,Sf*))],  
*vs*[*Suc* 0 := *the* (*S'a* (*nat* *n*)),  
2 := *the* (*S'a* (*nat* (*n* - 1))),  
3 := *the* (*S'a* (*nat* (*n* - 2)))]],  
*safeP, sigSafeMain*,  
*the* (*cdm* (*l*, *i*) + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1  
+ 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1, *aa, ba*]))))  
**apply** (*unfold exec-all-def*)  
**apply** (*erule rtrancl-trans*)  
**prefer** 2  
**apply** (*rule r-into-rtrancl*)  
**apply** (*clarify*)  
**apply** (*unfold JVMEExec.exec.simps*)  
**apply** (*unfold Let-def*)  
**apply** (*subgoal-tac* (*fst* (*snd* (*snd* (*snd* (*snd* (*the* (*method'* (*P'*, *safeP*) *sigSafeMain*))))))) !  
(Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc  
(Suc (Suc (Suc (Suc (*the* (*cdm* (*l*, *i*)))))))))))))))) =  
*Putstatic topf stackC,simp*)  
**apply** (*erule nth-via-drop-append*)

**apply** (*drule drop-Suc-append*)  
**apply** (*subgoal-tac*  $P \vdash$  (*None, sha*((*stackC, topf*)  $\mapsto$  *Intg* (*n* - 2)),*h,inih*,







```

+ 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1+1, aa, ba))))
apply (unfold exec-all-def)
apply (erule rtrancl-trans)
prefer 2
apply (rule r-into-rtrancl)
apply (clarify)
apply (unfold JVMExec.exec.simps)
apply (unfold Let-def)
apply (subgoal-tac (fst (snd (snd (snd (snd (the (method' (P', safeP) sigSafe-
Main)))))) !
      (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc
(Suc (Suc (Suc (Suc (Suc
      (Suc (Suc (Suc (Suc (the (cdm (l, i)))))))))))))))))) =
      Load 2,simp)
apply (erule nth-via-drop-append)

apply (drule drop-Suc-append)
apply (subgoal-tac  $P \vdash$  (None, sha((stackC, topf)  $\mapsto$  Intg (n - 2)),
      h(la  $\mapsto$  Arr ty m (S'a(nat (n - 2)  $\mapsto$  vs[Suc 0 := the
(S'a (nat n))] ! Suc 0))),inih,
      [([vs[Suc 0 := the (S'a (nat n)),
        2 := the (S'a (nat (n - 1)))]!2],
      vs[Suc 0 := the (S'a (nat n)),
        2 := the (S'a (nat (n - 1))),
        3 := the (S'a (nat (n - 2)))]),
      safeP, sigSafeMain,
      the (cdm (l, i) + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1
+ 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1+1, aa, ba))]
      -jvm→
      (None, sha((stackC, topf)  $\mapsto$  Intg (n - 2), (heapC, kOf)  $\mapsto$ 
      vs[Suc 0 := the (S'a (nat n)),
        2 := the
(S'a (nat (n - 1)))] ! 2),
      h(la  $\mapsto$  Arr ty m (S'a(nat (n - 2)  $\mapsto$  vs[Suc 0 := the
(S'a (nat n))] ! Suc 0))),inih,
      [([],
      vs[Suc 0 := the (S'a (nat n)),
        2 := the (S'a (nat (n - 1))),
        3 := the (S'a (nat (n - 2)))]),
      safeP, sigSafeMain,
      the (cdm (l, i) + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1
+ 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1+1 + 1 + 1, aa, ba))]
apply (unfold exec-all-def)
apply (erule rtrancl-trans)
prefer 2
apply (rule r-into-rtrancl)
apply (clarify)

```





Main)))))) !  
 (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc  
 (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc  
 (Suc (Suc (Suc (Suc (the (cdm (l, i)))))))))))))))))) =  
 Goto (trAddr pcc (the (cdm (l, i)) + incPop)), simp)  
**apply** (erule nth-via-drop-append)

**apply** (subgoal-tac drop pcc (extractBytecode P') =  
 [Tableswitch 1 (int (length ps)) ps] @ bytecode'  $\wedge$  ctm (nat (ps!i))  
 = Some i)  
**apply** (elim conjE)  
**prefer** 2 **apply** (rule Tableswitch-POPCONT)  
**apply** (subgoal-tac P $\vdash$  (None, sha((stackC, topf)  $\mapsto$  Intg (n - 2), (heapC, kOf)  
 $\mapsto$  vs[Suc 0 := the (S'a (nat n)),  
 2 := the  
 (S'a (nat (n - 1))) ! 2),  
 h(la  $\mapsto$  Arr ty m (S'a(nat (n - 2)  $\mapsto$  vs[Suc 0 := the  
 (S'a (nat n)) ! Suc 0)), inih,  
 [[(vs[Suc 0 := the (S'a (nat n)),  
 2 := the (S'a (nat (n - 1))),  
 3 := the (S'a (nat (n - 2)))] ! 3],  
 vs[Suc 0 := the (S'a (nat n)),  
 2 := the (S'a (nat (n - 1))),  
 3 := the (S'a (nat (n - 2)))]],  
 safeP, sigSafeMain,  
 nat (24 + int (the (cdm (l, i))) +  
 trAddr pcc (the (cdm (l, i)) + incPop)), aa, ba))]  
 -jvm $\rightarrow$   
 (None, sha((stackC, topf)  $\mapsto$  Intg (n - 2), (heapC, kOf)  $\mapsto$   
 vs[Suc 0 := the (S'a (nat n)),  
 2 := the  
 (S'a (nat (n - 1))) ! 2),  
 h(la  $\mapsto$  Arr ty m (S'a(nat (n - 2)  $\mapsto$  vs[Suc 0 := the  
 (S'a (nat n)) ! Suc 0)), inih,  
 [[[],  
 vs[Suc 0 := the (S'a (nat n)),  
 2 := the (S'a (nat (n - 1))),  
 3 := the (S'a (nat (n - 2)))]],  
 safeP, sigSafeMain,  
 if' (((the-Intg (vs[Suc 0 := the (S'a (nat n)),  
 2 := the (S'a (nat (n - 1))),  
 3 := the (S'a (nat (n - 2)))] ! 3) < 1)  $\vee$   
 (int (length ps) < the-Intg (vs[Suc 0 := the (S'a (nat n)),  
 2 := the (S'a (nat (n - 1))),  
 3 := the (S'a (nat (n - 2)))] ! 3))))  
 (nat (24 + int (the (cdm (l, i))) + trAddr pcc (the (cdm (l, i))  
 + incPop)) +  
 nat (ps ! nat (int (length ps))))

```

      (nat (24 + int (the (cdm (l, i))) + trAddr pcc (the (cdm (l, i))
+ incPop)) +
      nat (ps ! nat (the-Intg (vs[Suc 0 := the (S'a (nat n)),
      2 := the (S'a (nat (n - 1))),
      3 := the (S'a (nat (n - 2)))] ! 3)
- 1))),
      aa, ba]))
apply (unfold exec-all-def)
apply (erule rtrancl-trans)
prefer 2
apply (rule r-into-rtrancl)
apply (clarify)
apply (unfold extractBytecode-def)
apply (unfold JVMExec.exec.simps)
apply (unfold Let-def)
apply (subgoal-tac (fst (snd (snd (snd (snd (the (method' (P', safeP) sigSafe-
Main))))))) !
      (nat (24 + int (the (cdm (l, i))) + trAddr pcc (the (cdm (l, i))
+ incPop))) =
      Tableswitch 1 (int (length ps)) ps,simp)
apply (simp add: if'-def)
apply (subgoal-tac
  pcc = nat (24 + int (the (cdm (l, i))) + (int pcc - int (the (cdm (l, i)) +
incPop))))
prefer 2 apply (simp add: incPop-def)
apply (simp add: trAddr-def)
apply (erule nth-via-drop)
apply simp

apply (rule-tac x=hm in exI)
apply (rule-tac x=k in exI)
apply (rule-tac x=k1 in exI)
apply (rule-tac x=(l1,0) in exI)
apply (rule-tac x=Val v1 # S' in exI)

apply (rule-tac x=sha((stackC, topf) ↦ Intg (n - 2),
      (heapC, kOf) ↦ vs[Suc 0 := the (S'a (nat n)), 2 := the (S'a
(nat (n - 1)))] ! 2) in exI)
apply (rule-tac x=h(la ↦ Arr ty m (S'a(nat (n - 2) ↦ vs[Suc 0 := the (S'a
(nat n)] ! Suc 0))) in exI)
apply (rule-tac x=inih in exI)
apply (rule-tac x=vs[Suc 0 := the (S'a (nat n)),
      2 := the (S'a (nat (n - 1))),
      3 := the (S'a (nat (n - 2)))] in exI)
apply (rule-tac x=if' (((the-Intg (vs[Suc 0 := the (S'a (nat n)),
      2 := the (S'a (nat (n - 1))),

```

```

      3 := the (S'a (nat (n - 2))) ! 3) < 1) ∨
      (int (length ps) < the-Intg (vs[Suc 0 := the (S'a (nat n)),
      2 := the (S'a (nat (n - 1))),
      3 := the (S'a (nat (n - 2))) ! 3])))
+ incPop)) +
      nat (ps ! nat (int (length ps)))
+ incPop)) +
      nat (ps ! nat (the-Intg (vs[Suc 0 := the (S'a (nat n)),
      2 := the (S'a (nat (n - 1))),
      3 := the (S'a (nat (n - 2))) ! 3]
- 1))) in exI)
apply (rule-tac x=ref in exI)

apply (rule-tac x=nat (the-Intg (the ((sha((stackC, topf) ↦ Intg (n - 2),
      (heapC, k0f) ↦ vs[Suc 0 := the (S'a (nat n)),
      2 := the (S'a (nat (n - 1)))])
! 2)) (heapC, kf)))) in exI)

apply (rule-tac x=nat (the-Intg (the ((sha((stackC, topf) ↦ Intg (n - 2),
      (heapC, k0f) ↦ vs[Suc 0 := the (S'a (nat n)),
      2 := the (S'a (nat (n - 1)))])
! 2)) (heapC, k0f)))) in exI)

apply (rule-tac x=la in exI)
apply (rule-tac x=ty in exI)
apply (rule-tac x=m in exI)
apply (rule-tac x=S'a(nat (n - 2) ↦ vs[Suc 0 := the (S'a (nat n))] ! Suc 0) in
exI)

apply (rule-tac x= n - 2 in exI)
apply (rule-tac x=l' in exI)
apply (rule-tac x=regS in exI)
apply (rule-tac x=l'' in exI)
apply (rule-tac x=m' in exI)
apply (rule-tac x=m'' in exI)
apply (rule-tac x=d in exI)
apply (rule-tac x=g in exI)

apply (rule conjI, rule refl)
apply (rule conjI) apply (unfold trAddr-def)
      apply (rule refl)
apply (rule conjI, rule refl)
apply (rule conjI, rule refl)
apply (rule conjI) apply clarsimp
      apply (simp add: Sf-def add: topf-def)
      apply (simp add: stackC-def add: heapC-def)

```

```

apply (rule conjI, assumption)
apply (rule conjI) apply clarsimp
      apply (simp add: stackC-def add: heapC-def)
      apply (simp add: kf-def add: k0f-def)
apply (rule conjI) apply (erule activeCells)
apply (rule conjI) apply (erule activeCells-2,assumption)
      apply (erule activeCells-2,assumption)
apply (rule conjI, simp)
apply (rule conjI) apply (subgoal-tac length vs = 10)
      apply (simp add: stackC-def add: heapC-def)
      apply (rule length-vs)
apply (rule conjI) apply clarsimp
      apply (simp add: stackC-def add: heapC-def)
      apply (simp add: regionsf-def add: k0f-def)
apply (rule conjI) apply clarsimp
apply (rule conjI) apply clarsimp
      apply (simp add: dirCellC-def add: stackC-def)
      apply (simp add: heapC-def add: safeDirf-def)
apply (rule conjI) apply clarsimp
apply (rule conjI) apply clarsimp
apply (rule conjI) apply clarsimp
apply (rule conjI) apply (subgoal-tac length vs = 10)
      apply (rule equivH-POPCONT,assumption+)
      apply simp
      apply (rule length-vs)
apply (rule conjI) apply clarsimp
apply (rule conjI) apply clarsimp
      apply (subgoal-tac length vs = 10)
      apply (rule equivS2-POPCONT,simp,assumption+)
      apply (rule length-vs)
apply (rule conjI) apply (subgoal-tac length vs = 10)
      apply clarsimp
      apply (rule equivS-k1,assumption+)
      apply (rule length-vs)
by (rule equivpc-POPCONT)

end

theory dem-PRIMOP
imports ../JVMSAFE/JVMSAFE JVMExec SVM2JVM SVMSemantics CertifSVM2JVM
      dem-translation
begin

no-translations Norm s == (None,s)

no-translations ex-table-of m == snd (snd (snd m))

```



```

declare trInstr.simps [simp del]
declare equivS.simps [simp del]
declare extractBytecode-def [simp del]
declare initConsTable-def [simp del]
declare exec-instr.simps [simp del]

```

```

lemma equivS-PRIMOP:
  equivS (Val v1 # Val v2 # S) S' n ctm d g  $\implies$  n  $\geq$  1
apply (induct S)
apply (unfold equivS.simps, clarsimp)
by (case-tac a, simp-all)

```

```

lemma equivS2-PRIMOP:
  [[ 2 < length vs; equivS (Val (IntT i) # Val (IntT i') # S') S'a n ctm d g]]
   $\implies$  equivS (Val (execOp oper (IntT i) (IntT i')) # S')
    (S'a(nat (n - 1)  $\mapsto$  applyBinOp oper (the (S'a (nat n))))
      (vs[Suc 0 := the (S'a (nat (n - 1)))] ! Suc 0)) (n - 1) ctm d g)
apply (simp add: equivS.simps, clarsimp)
apply (rule conjI)
apply (case-tac oper, simp-all)
apply (drule equivS-ge-n)
apply (erule-tac x=1 in allE, simp)
by (subgoal-tac -1 + n = n - 1, simp, simp)

```

```

lemma activeCells-PRIMOP:
  [[ la  $\notin$  activeCells regS h (nat (the-Intg (the (sha (heapC, kf))))))]
   $\implies$  activeCells regS h (nat (the-Intg (the (sha (heapC, kf)))))) =
    activeCells regS (h(la  $\mapsto$  Arr ty m (S'a(nat (n - 1)  $\mapsto$ 
      applyBinOp oper (the (S'a (nat n))) (vs[Suc 0 := the (S'a (nat (n
- 1)))] ! Suc 0))))))
      (nat (the-Intg (the (sha (heapC, kf))))))
apply (frule l-not-in-cellReg)
apply (unfold activeCells-def, auto)
apply (unfold region-def, simp add: Let-def, elim conjE)
apply (rule-tac x=j in exI, simp)
apply (rule cellReg-step)
apply (rule cellReg-basic)
apply (erule-tac x=j in allE)
apply (rule cellsReg-monotone-1, assumption+)
apply (erule l-not-in-regs)
apply (simp add: Let-def, elim conjE)
apply (rule-tac x=j in exI, simp)
apply (rule cellReg-step)
apply (rule cellReg-basic)

```

**apply** (erule-tac x=j in allE)  
**apply** (erule cellsReg-monotone-2,assumption+)  
**by** (erule l-not-in-regs)

**lemma** domH-PRIMOP:

$\llbracket$  la  $\notin$  activeCells regS h (nat (the-Intg (the (sha (heapC, kf))));  
 $\forall l \in \text{dom } H. \exists l'. l' = \text{the } (g \ l) \wedge \text{equivC } (\text{the } (H \ l)) \ h \ l' \ (\text{the } (h \ l')) \ (\text{nat } (\text{the-Intg } (\text{the } (\text{sha } (\text{heapC}, \text{kf})))))) \ \text{com regS } d \ g \rrbracket$   
 $\implies \forall l \in \text{dom } H.$   
 $\exists l'. l' = \text{the } (g \ l) \wedge$   
 $\text{equivC } (\text{the } (H \ l))$   
 $(h(la \mapsto$   
 $\text{Arr ty m } (S'a(\text{nat } (n - 1) \mapsto \text{applyBinOp oper } (\text{the } (S'a (\text{nat } n))))$   
 $(\text{vs}[Suc \ 0 := \text{the } (S'a (\text{nat } (n - 1))]) ! \text{Suc } 0)))))$   
 $l' \ (\text{the } ((h(la \mapsto$   
 $\text{Arr ty m}$   
 $(S'a(\text{nat } (n - 1) \mapsto \text{applyBinOp oper } (\text{the } (S'a (\text{nat } n))))$   
 $(\text{vs}[Suc \ 0 := \text{the } (S'a (\text{nat } (n - 1))]) ! \text{Suc } 0)))))$   
 $l')$   
 $(\text{nat } (\text{the-Intg } (\text{the } (\text{sha } (\text{heapC}, \text{kf}))))))$   
 $\text{com regS } d \ g$

**apply** (frule l-not-in-cellReg)  
**apply** (rule ballI)  
**apply** (erule-tac x=l in ballE)  
**prefer** 2 **apply** simp  
**apply** (erule exE)  
**apply** (rule-tac x=l' in exI) **apply** (elim conjE)  
**apply** (rule conjI, assumption)  
**apply** clarsimp  
**apply** (rule conjI) **apply** (rule impI)  
**apply** clarsimp **apply** (simp add: region-def)  
**apply** clarsimp  
**apply** (rule-tac x=Obj in exI)  
**apply** (rule-tac x=flds in exI) **apply** simp  
**apply** (simp add: region-def) **apply** (elim conjE)  
**apply** (rule cellReg-step)  
**apply** (rule cellReg-basic)  
**apply** (erule-tac x=j in allE)  
**apply** (rule cellsReg-monotone-1,assumption+)  
**by** (erule l-not-in-regs)

**lemma** equivH-PRIMOP:

$\llbracket S = \text{Val } (\text{IntT } ia) \# \text{Val } (\text{IntT } i') \# S'; ((hm, k), k0, (l, i), S) = ((H, ka),$   
 $k0a, PC, Sa);$   
 $k' = \text{nat } (\text{the-Intg } (\text{the } (\text{sha } (\text{heapC}, \text{kf}))))); k0' = \text{nat } (\text{the-Intg } (\text{the } (\text{sha } (\text{heapC}, k0f))))); \text{sha } (\text{stackC}, Sf) = \text{Some } (\text{Addr } la);$

$distinct [la, l', l'']; activeCells regS h k' \cap \{la, l', l''\} = \{\}; h la = Some (Arr ty m S'a);$   
 $sha (stackC, topf) = Some (Intg n); sha (heapC, regionsf) = Some (Addr l');$   
 $h l' = Some (Arr ty m' regS);$   
 $sha (dirCellC, safeDirf) = Some (Addr l''); h l'' = Some (Arr ty m'' d); inj-on$   
 $g (dom H); equivH (H, ka) h k' com regS d g$   
 $\Downarrow$   
 $\implies equivH (hm, k)$   
 $(h(la \mapsto Arr ty m (S'a(nat (n - 1) \mapsto applyBinOp oper (the (S'a (nat n))))$   
 $(vs[Suc 0 := the (S'a (nat (n - 1))]) ! Suc 0))))))$   
 $(nat (the-Intg (the ((sha((stackC, topf) \mapsto Intg (n - 1))) (heapC, kf))))))$   
 $com regS d g$   
**apply** (*simp add: heapC-def add: stackC-def*)  
**apply** (*simp add: kf-def add: k0f-def*)  
**apply** (*fold heapC-def, fold kf-def*)  
**apply** (*unfold equivH.simps, elim conjE*)  
**apply** (*rule conjI, assumption*)  
**apply** (*rule conjI, frule activeCells-PRIMOP, simp*)  
**apply** (*rule conjI, assumption*)  
**by** (*rule domH-PRIMOP, assumption+*)

**declare** *exec-instr.simps* [*simp*]

**lemma** *execSVMInstr-PRIMOP* :

$\llbracket (P', cdm, ctm, com) = trSVM2JVM ((svms, ctmap), ini, ct, ah, ai, bc);$   
 $cdm, ctm, com \vdash ((hm, k), k0, (l, i), S) \triangleq S1';$   
 $(fst (the (map-of svms l)) ! i) = PRIMOP oper;$   
 $execSVMInst (PRIMOP oper) (map-of ct) (hm, k) k0 (l, i) S = Either.Right$   
 $S2;$   
 $drop (the (cdm (l, i))) (extractBytecode P') =$   
 $trInstr (the (cdm (l, i))) cdm' ctm' com pcc (PRIMOP oper) @ bytecode'$   
 $\llbracket \implies \exists v' sh' dh' ih' fms'.$   
 $P' \vdash S1' -jvm \rightarrow (v', sh', dh', ih', fms') \wedge$   
 $cdm, ctm, com \vdash S2 \triangleq (v', sh', dh', ih', fms')$

**apply** (*case-tac S1'*)  
**apply** (*rename-tac v tup*)  
**apply** (*case-tac tup*)  
**apply** (*rename-tac sh tup*)  
**apply** (*case-tac tup*)  
**apply** (*rename-tac dh tup*)  
**apply** (*case-tac tup*)  
**apply** (*rename-tac ih fms*)  
**apply** (*simp*)

**apply** (*subgoal-tac*  $\exists i i' S'. S = \text{Val } (\text{IntT } i) \# \text{Val } (\text{IntT } i') \# S'$ )  
**prefer** 2

**apply** (*case-tac*  $S, \text{simp-all}$ )  
**apply** (*insert RightNotUndefined*, *erule-tac*  $x = S2$  **in**  $\text{allE}, \text{force}$ )  
**apply** *clarsimp*  
**apply** (*case-tac*  $af, \text{simp-all}$ )  
**apply** (*case-tac*  $\text{Vala}, \text{simp-all}$ )  
**apply** (*erule-tac*  $x = S2$  **in**  $\text{allE}, \text{force}$ )  
**apply** (*case-tac*  $\text{list}, \text{simp}, \text{simp-all}$ )  
**apply** *clarsimp*  
**apply** (*case-tac*  $a, \text{simp-all}$ )  
**apply** (*case-tac*  $\text{Vala}, \text{simp-all}$ )  
**apply** (*insert RightNotUndefined*, *erule-tac*  $x = S2$  **in**  $\text{allE}, \text{force}$ )  
**apply** (*erule-tac*  $x = S2$  **in**  $\text{allE}, \text{force}$ )  
**apply** (*erule-tac*  $x = S2$  **in**  $\text{allE}, \text{force}$ )  
**apply** (*erule-tac*  $x = S2$  **in**  $\text{allE}, \text{force}$ )

**apply** (*unfold equivState-def*)  
**apply** (*elim exE, elim conjE*)

**apply** (*rule-tac*  $x = \text{None}$  **in**  $\text{exI}$ )  
**apply** (*rule-tac*  $x = \text{sha}((\text{stackC}, \text{topf}) \mapsto \text{Intg } (n - 1))$  **in**  $\text{exI}$ )  
**apply** (*rule-tac*  $x = h(\text{la} \mapsto \text{Arr } \text{ty } m (S'a(\text{nat } (n - 1) \mapsto \text{applyBinOp } \text{oper } (\text{the } (S'a (\text{nat } n))))$   
( $\text{vs}[\text{Suc } 0 := \text{the } (S'a (\text{nat } (n - 1))]$ )  
**! Suc 0))**) **in**  $\text{exI}$ )  
**apply** (*rule-tac*  $x = \text{inih}$  **in**  $\text{exI}$ )  
**apply** (*rule-tac*  $x = [(\ [], \text{vs}[\text{Suc } 0 := \text{the } (S'a (\text{nat } (n - 1))]), \text{safeP}, \text{sigSafeMain}, \text{the } (\text{cdm } (l, i)) +$   
 $1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1$   
 $+ 1 + 1 + 1, \text{ref}]$ ) **in**  $\text{exI}$ )

**apply** (*subgoal-tac*  $\exists \text{vs } \text{pc} . \text{fms} = (\ [], \text{vs}, \text{safeP}, \text{sigSafeMain}, \text{pc}, \text{ref}) \# (\ [])$ )  
**prefer** 2 **apply** *simp*  
**apply** (*erule exE*) +  
**apply** (*unfold exec-all-def*)  
**apply** (*rule conjI*)

**apply** (*subgoal-tac*  $\text{PC} = (l, i)$ ) **prefer** 2 **apply** *simp*  
**apply** *simp* **apply** (*elim conjE*) **apply** *clarsimp*  
**apply** (*subgoal-tac*  $P \uparrow \vdash (\text{None}, \text{sha}, h, \text{inih}, [(\ [], \text{vs}, \text{safeP}, \text{sigSafeMain}, \text{the } (\text{cdm } (l, i)), \text{ag}, \text{bd})]) - \text{jvm} \rightarrow$

```

      (None, sha, h, inih, [[the (sha (stackC, Sf))], vs, safeP, sigSafeMain,
the (cdm (l, i)) + 1, ag, bd]))
apply (unfold exec-all-def)
apply (erule rtrancl-trans)

```

```

prefer 2
apply (rule r-into-rtrancl)
apply (clarify)
apply (unfold JVMExec.exec.simps)
apply (unfold Let-def)
apply (unfold extractBytecode-def)
apply (subgoal-tac (fst (snd (snd (snd (snd (the (method' (P', safeP) sigSafe-
Main))))))) !
      the (cdm (l, i))) = Getstatic Sf stackC, simp)
apply (unfold trInstr.simps)
apply (erule nth-via-drop-append)

```

```

apply (drule drop-Suc-append)
apply (subgoal-tac P⊢ (None, sha, h, inih,
      [[the (sha (stackC, Sf))],
      vs, safeP, sigSafeMain, the (cdm (l, i)) + 1, ag, bd]])
      -jvm→
      (None, sha, h, inih,
      [[the (sha (stackC, topf)),
      the (sha (stackC, Sf))],
      vs, safeP, sigSafeMain, the (cdm (l, i)) + 1 + 1, ag, bd]))

```

```

apply (unfold exec-all-def)
apply (erule rtrancl-trans)
prefer 2
apply (rule r-into-rtrancl)
apply (clarify)
apply (unfold JVMExec.exec.simps)
apply (unfold Let-def)
apply (subgoal-tac (fst (snd (snd (snd (snd (the (method' (P', safeP) sigSafe-
Main))))))) ! Suc (the (cdm (l, i)))) =
      Getstatic topf stackC, simp)
apply (erule nth-via-drop-append)

```

```

apply (drule drop-Suc-append)
apply (subgoal-tac P⊢ (None, sha, h, inih,
      [[the (sha (stackC, topf)),
      the (sha (stackC, Sf))],
      vs, safeP, sigSafeMain, the (cdm (l, i)) + 1 + 1, ag, bd]])
      -jvm→
      (None, sha, h, inih,
      [[Intg 1,

```

```

      the (sha (stackC,topf)),
      the (sha (stackC,Sf))),
      vs, safeP, sigSafeMain, the (cdm (l, i) + 1 + 1 + 1, ag, bd)))
apply (unfold exec-all-def)
apply (erule rtrancl-trans)
prefer 2
  apply (rule r-into-rtrancl)
  apply (clarify)
  apply (unfold JVMExec.exec.simps)
  apply (unfold Let-def)
  apply (subgoal-tac (fst (snd (snd (snd (snd (the (method' (P', safeP) sigSafe-
Main)))))) ! (Suc (Suc (the (cdm (l, i)))))) =
      LitPush (Intg 1),simp)
apply (erule nth-via-drop-append)

```

```

apply (drule drop-Suc-append)
apply (subgoal-tac P† (None,sha,h,inih,
      [[Intg 1,
        the (sha (stackC,topf)),
        the (sha (stackC,Sf))],
      vs, safeP, sigSafeMain, the (cdm (l, i) + 1 + 1 + 1, ag, bd))]
      -jvm→
      (None,sha,h,inih,
      [[Intg (the-Intg (the (sha (stackC, topf))) - 1),
        the (sha (stackC,Sf))],
      vs, safeP, sigSafeMain, the (cdm (l, i) + 1 + 1 + 1 + 1, ag,
      bd]]))
apply (unfold exec-all-def)
apply (erule rtrancl-trans)
prefer 2
  apply (rule r-into-rtrancl)
  apply (clarify)
  apply (unfold JVMExec.exec.simps)
  apply (unfold Let-def)
  apply (subgoal-tac (fst (snd (snd (snd (snd (the (method' (P', safeP) sigSafe-
Main)))))) ! (Suc (Suc (Suc (the (cdm (l, i)))))) =
      BinOp Subtract,simp)
apply (erule nth-via-drop-append)

```

```

apply (drule drop-Suc-append)
apply (subgoal-tac P† (None,sha,h,inih,
      [[Intg (the-Intg (the (sha (stackC, topf))) - 1),
        the (sha (stackC,Sf))],
      vs, safeP, sigSafeMain, the (cdm (l, i) + 1 + 1 + 1 + 1,

```

```

ag, bd)])
      -jvm→
      (None,sha,h,inih,
      [[Intg (the-Intg (the (sha (stackC, topf))) - 1),
      the (sha (stackC,Sf)),
      Intg (the-Intg (the (sha (stackC, topf))) - 1),
      the (sha (stackC,Sf))],
      vs, safeP, sigSafeMain, the (cdm (l, i)) + 1 + 1 + 1 + 1 +
1, ag, bd]]))
apply (unfold exec-all-def)
apply (erule rtrancl-trans)
prefer 2
apply (rule r-into-rtrancl)
apply (clarify)
apply (unfold JVMExec.exec.simps)
apply (unfold Let-def)
apply (subgoal-tac (fst (snd (snd (snd (snd (the (method' (P', safeP) sigSafe-
Main))))))) !
      (Suc (Suc (Suc (Suc (the (cdm (l, i))))))) =
      Dup2,simp)
apply (erule nth-via-drop-append)

apply (drule drop-Suc-append)
apply (subgoal-tac P ⊢ (None,sha,h,inih,
      [[Intg (the-Intg (the (sha (stackC, topf))) - 1),
      the (sha (stackC,Sf)),
      Intg (the-Intg (the (sha (stackC, topf))) - 1),
      the (sha (stackC,Sf))],
      vs, safeP, sigSafeMain, the (cdm (l, i)) + 1 + 1 + 1 + 1 +
1, ag, bd]]))
      -jvm→
      (None,sha,h,inih,
      [[the (S'a (nat (n - 1))),
      Intg (the-Intg (the (sha (stackC, topf))) - 1),
      the (sha (stackC,Sf))],
      vs, safeP, sigSafeMain, the (cdm (l, i)) + 1 + 1 + 1 + 1 +
1 + 1, ag, bd]]))
apply (unfold exec-all-def)
apply (erule rtrancl-trans)
prefer 2
apply (rule r-into-rtrancl)
apply (clarify)
apply (unfold JVMExec.exec.simps)
apply (unfold Let-def)
apply (subgoal-tac (fst (snd (snd (snd (snd (the (method' (P', safeP) sigSafe-
Main))))))) !
      (Suc (Suc (Suc (Suc (Suc (the (cdm (l, i))))))) =
      ArrLoad,simp)

```

**apply** (*frule equivS-PRIMOP, simp*)  
**apply** (*simp add: raise-system-xcpt-def*)  
**apply** (*erule nth-via-drop-append*)

**apply** (*drule drop-Suc-append*)  
**apply** (*subgoal-tac P<sup>⊢</sup> (None, sha, h, inih,*  
 $[[[the (S'a (nat (n - 1))),$   
 $Intg (the-Intg (the (sha (stackC, topf))) - 1),$   
 $the (sha (stackC, Sf))],$   
 $vs, safeP, sigSafeMain, the (cdm (l, i)) + 1 + 1 + 1 + 1 +$   
 $1 + 1, ag, bd]]]$   
 $-jvm \rightarrow$   
 $(None, sha, h, inih,$   
 $[[[Intg (the-Intg (the (sha (stackC, topf))) - 1),$   
 $the (sha (stackC, Sf))],$   
 $vs[Suc 0 := the (S'a (nat (n - 1))],$   
 $safeP, sigSafeMain, the (cdm (l, i)) + 1 + 1 + 1 + 1 + 1$   
 $+ 1 + 1, ag, bd]]]$ )  
**apply** (*unfold exec-all-def*)  
**apply** (*erule rtrancl-trans*)  
**prefer** 2  
**apply** (*rule r-into-rtrancl*)  
**apply** (*clarify*)  
**apply** (*unfold JVMExec.exec.simps*)  
**apply** (*unfold Let-def*)  
**apply** (*subgoal-tac (fst (snd (snd (snd (snd (the (method' (P', safeP) sigSafe-*  
 $Main)))))) !$   
 $(Suc (Suc (Suc (Suc (Suc (Suc (the (cdm (l, i)))))))) =$   
 $Store 1, simp)$ )  
**apply** (*erule nth-via-drop-append*)

**apply** (*drule drop-Suc-append*)  
**apply** (*subgoal-tac P<sup>⊢</sup> (None, sha, h, inih,*  
 $[[[Intg (the-Intg (the (sha (stackC, topf))) - 1),$   
 $the (sha (stackC, Sf))],$   
 $vs[Suc 0 := the (S'a (nat (n - 1))],$   
 $safeP, sigSafeMain, the (cdm (l, i)) + 1 + 1 + 1 + 1 + 1$   
 $+ 1 + 1, ag, bd]]]$   
 $-jvm \rightarrow$   
 $(None, sha, h, inih,$   
 $[[[the (sha (stackC, Sf)),$   
 $Intg (the-Intg (the (sha (stackC, topf))) - 1),$   
 $the (sha (stackC, Sf))],$   
 $vs[Suc 0 := the (S'a (nat (n - 1))],$   
 $safeP, sigSafeMain, the (cdm (l, i)) + 1 + 1 + 1 + 1 + 1$



```

+ 1 + 1 + 1, ag, bd]))
apply (unfold exec-all-def)
apply (erule rtrancl-trans)
prefer 2
apply (rule r-into-rtrancl)
apply (clarify)
apply (unfold JVMExec.exec.simps)
apply (unfold Let-def)
apply (subgoal-tac (fst (snd (snd (snd (snd (the (method' (P', safeP) sigSafe-
Main))))))) !
      (Suc (Suc (Suc (Suc (Suc (Suc (Suc (the (cdm (l, i)))))))))) =
      Getstatic Sf stackC,simp)
apply (erule nth-via-drop-append)

apply (drule drop-Suc-append)
apply (subgoal-tac P† (None,sha,h,inih,
  [(the (sha (stackC,Sf)),
    Intg (the-Intg (the (sha (stackC, topf))) - 1),
    the (sha (stackC,Sf))],
    vs[Suc 0 := the (S'a (nat (n - 1))]),
    safeP, sigSafeMain, the (cdm (l, i)) + 1 + 1 + 1 + 1 + 1
+ 1 + 1 + 1, ag, bd]))
  -jvm→
  (None,sha,h,inih,
  [(the (sha (stackC,topf)),
    the (sha (stackC,Sf)),
    Intg (the-Intg (the (sha (stackC, topf))) - 1),
    the (sha (stackC,Sf))],
    vs[Suc 0 := the (S'a (nat (n - 1))]),
    safeP, sigSafeMain, the (cdm (l, i)) + 1 + 1 + 1 + 1 + 1
+ 1 + 1 + 1 + 1, ag, bd]))
apply (unfold exec-all-def)
apply (erule rtrancl-trans)
prefer 2
apply (rule r-into-rtrancl)
apply (clarify)
apply (unfold JVMExec.exec.simps)
apply (unfold Let-def)
apply (subgoal-tac (fst (snd (snd (snd (snd (the (method' (P', safeP) sigSafe-
Main))))))) !
      (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (the (cdm (l, i))))))))))
=
      Getstatic topf stackC,simp)
apply (erule nth-via-drop-append)

```

```

apply (drule drop-Suc-append)
apply (subgoal-tac P⊢ (None,sha,h,inih,
  [[(the (sha (stackC,topf))),
    the (sha (stackC,Sf)),
    Intg (the-Intg (the (sha (stackC, topf)))− 1),
    the (sha (stackC,Sf))]],
  vs[Suc 0 := the (S'a (nat (n − 1)))],
  safeP, sigSafeMain, the (cdm (l, i)) + 1 + 1 + 1 + 1 + 1
+ 1 + 1 + 1 + 1, ag, bd]))
  −jvm→
  (None,sha,h,inih,
  [[(the (S'a (nat n))),
    Intg (the-Intg (the (sha (stackC, topf)))− 1),
    the (sha (stackC,Sf))]],
  vs[Suc 0 := the (S'a (nat (n − 1)))],
  safeP, sigSafeMain, the (cdm (l, i)) + 1 + 1 + 1 + 1 + 1
+ 1 + 1 + 1 + 1 + 1, ag, bd]))))
apply (unfold exec-all-def)
apply (erule rtrancl-trans)
prefer 2
apply (rule r-into-rtrancl)
apply (clarify)
apply (unfold JVMExec.exec.simps)
apply (unfold Let-def)
apply (subgoal-tac (fst (snd (snd (snd (snd (the (method' (P', safeP) sigSafe-
Main)))))) !
  (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (the (cdm (l,
i)))))))))) =
  ArrLoad,simp)

  apply (frule equivS-PRIMOP, simp)
  apply (simp add: raise-system-xcpt-def)
apply (erule nth-via-drop-append)

```

```

apply (drule drop-Suc-append)
apply (subgoal-tac P⊢ (None,sha,h,inih,
  [[(the (S'a (nat n))),
    Intg (the-Intg (the (sha (stackC, topf)))− 1),
    the (sha (stackC,Sf))]],
  vs[Suc 0 := the (S'a (nat (n − 1)))],
  safeP, sigSafeMain, the (cdm (l, i)) + 1 + 1 + 1 + 1 + 1
+ 1 + 1 + 1 + 1 + 1, ag, bd]))
  −jvm→
  (None,sha,h,inih,
  [[(vs[Suc 0 := the (S'a (nat (n − 1)))] ! Suc 0,
  the (S'a (nat n)),
  Intg (the-Intg (the (sha (stackC, topf)))− 1),
  the (sha (stackC,Sf))]],

```

```

      vs[Suc 0 := the (S'a (nat (n - 1))),
        safeP, sigSafeMain, the (cdm (l, i)) + 1 + 1 + 1 + 1 + 1
+ 1 + 1 + 1 + 1 + 1 + 1, ag, bd)])
apply (unfold exec-all-def)
apply (erule rtrancl-trans)
prefer 2
apply (rule r-into-rtrancl)
apply (clarify)
apply (unfold JVMExec.exec.simps)
apply (unfold Let-def)
apply (subgoal-tac (fst (snd (snd (snd (snd (the (method' (P', safeP) sigSafe-
Main))))))) !
(Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (the (cdm
(l, i)))))))))))))) =
      Load 1, simp)
apply (erule nth-via-drop-append)

apply (drule drop-Suc-append)
apply (subgoal-tac P⊢ (None, sha, h, inih,
  [[(vs[Suc 0 := the (S'a (nat (n - 1)))] ! Suc 0,
    the (S'a (nat n)),
    Intg (the-Intg (the (sha (stackC, topf))) - 1),
    the (sha (stackC, Sf))),
    vs[Suc 0 := the (S'a (nat (n - 1)))]],
    safeP, sigSafeMain, the (cdm (l, i)) + 1 + 1 + 1 + 1 + 1
+ 1 + 1 + 1 + 1 + 1 + 1, ag, bd)])
  -jvm→
  (None, sha, h, inih,
    [[(applyBinOp oper (the (S'a (nat n))) (vs[Suc 0 := the (S'a
(nat (n - 1)))] ! Suc 0),
      Intg (the-Intg (the (sha (stackC, topf))) - 1),
      the (sha (stackC, Sf))),
      vs[Suc 0 := the (S'a (nat (n - 1)))]],
      safeP, sigSafeMain, the (cdm (l, i)) + 1 + 1 + 1 + 1 + 1
+ 1 + 1 + 1 + 1 + 1 + 1, ag, bd)])
apply (unfold exec-all-def)
apply (erule rtrancl-trans)
prefer 2
apply (rule r-into-rtrancl)
apply (clarify)
apply (unfold JVMExec.exec.simps)
apply (unfold Let-def)
apply (subgoal-tac (fst (snd (snd (snd (snd (the (method' (P', safeP) sigSafe-
Main))))))) !
(Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (the
(cdm (l, i)))))))))))))) =
      BinOp oper, simp)

```

**apply** (*erule nth-via-drop-append*)

**apply** (*drule drop-Suc-append*)

**apply** (*subgoal-tac*  $P \vdash$  (*None, sha, h, inih,*  
 $[[[$ *applyBinOp oper* (*the* ( $S'a$  (*nat*  $n$ )))] (*vs*[ $Suc\ 0 := the$  ( $S'a$   
(*nat* ( $n - 1$ )))] !  $Suc\ 0$ ),  
*Intg* (*the-Intg* (*the* (*sha* (*stackC, topf*))) - 1),  
*the* (*sha* (*stackC, Sf*))),  
*vs*[ $Suc\ 0 := the$  ( $S'a$  (*nat* ( $n - 1$ )))]],  
*safeP, sigSafeMain, the* (*cdm* ( $l, i$ )) + 1 + 1 + 1 + 1 + 1  
+ 1 + 1 + 1 + 1 + 1 + 1 + 1, *ag, bd*]))  
 $-jvm \rightarrow$   
(*None, sha,*  
*h*( $la \mapsto Arr\ ty\ m$  ( $S'a$ (*nat* ( $n - 1$ ))  $\mapsto$  *applyBinOp oper* (*the*  
( $S'a$  (*nat*  $n$ )))] (*vs*[ $Suc\ 0 := the$  ( $S'a$  (*nat* ( $n - 1$ )))] !  
 $Suc\ 0$ ))),  
*inih,*  
 $[[[$  $\square$ ],  
*vs*[ $Suc\ 0 := the$  ( $S'a$  (*nat* ( $n - 1$ )))]],  
*safeP, sigSafeMain, the* (*cdm* ( $l, i$ )) + 1 + 1 + 1 + 1 + 1  
+ 1 + 1 + 1 + 1 + 1 + 1 + 1, *ag, bd*]]))

**apply** (*unfold exec-all-def*)

**apply** (*erule rtrancl-trans*)

**prefer** 2

**apply** (*rule r-into-rtrancl*)

**apply** (*clarify*)

**apply** (*unfold JVMExec.exec.simps*)

**apply** (*unfold Let-def*)

**apply** (*subgoal-tac* (*fst* (*snd* (*snd* (*snd* (*snd* (*the* (*method'* ( $P'$ , *safeP*) *sigSafe-*  
*Main*)))))) !

(*Suc* (*Suc* (*Suc* (*Suc* (*Suc* (*Suc* (*Suc* (*Suc* (*Suc* (*Suc* (*Suc* (*Suc* (*Suc*  
(*the* (*cdm* ( $l, i$ )))))))))))))) =  
*ArrStore, simp*)

**apply** (*simp add: raise-system-xcpt-def*)

**apply** (*frule equivS-PRIMOP, simp*)

**apply** (*erule nth-via-drop-append*)

**apply** (*drule drop-Suc-append*)

**apply** (*subgoal-tac*  $P \vdash$  (*None, sha,*  
*h*( $la \mapsto Arr\ ty\ m$  ( $S'a$ (*nat* ( $n - 1$ ))  $\mapsto$  *applyBinOp oper* (*the*  
( $S'a$  (*nat*  $n$ )))] (*vs*[ $Suc\ 0 := the$  ( $S'a$  (*nat* ( $n - 1$ )))] !  
 $Suc\ 0$ ))),  
*inih,*  
 $[[[$  $\square$ ],  
*vs*[ $Suc\ 0 := the$  ( $S'a$  (*nat* ( $n - 1$ )))]],

```

      safeP, sigSafeMain, the (cdm (l, i)) + 1 + 1 + 1 + 1 + 1
+ 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1, ag, bd)])
    -jvm→
      (None, sha,
        h(la ↦ Arr ty m (S'a(nat (n - 1)) ↦ applyBinOp oper (the
(S'a (nat n))) (vs[Suc 0 := the (S'a (nat (n - 1))]) !
          Suc 0))))),
        inih,
        [([the (sha (stackC, topf))],
          vs[Suc 0 := the (S'a (nat (n - 1))]),
          safeP, sigSafeMain, the (cdm (l, i)) + 1 + 1 + 1 + 1 + 1
+ 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1, ag, bd)])])
apply (unfold exec-all-def)
apply (erule rtrancl-trans)
prefer 2
apply (rule r-into-rtrancl)
apply (clarify)
apply (unfold JVMExec.exec.simps)
apply (unfold Let-def)
apply (subgoal-tac (fst (snd (snd (snd (snd (the (method' (P', safeP) sigSafe-
Main))))))) !
      (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc
(Suc (the (cdm (l, i)))))))))))))))))) =
        Getstatic topf stackC, simp)
apply (erule nth-via-drop-append)

apply (drule drop-Suc-append)
apply (subgoal-tac P ⊢ (None, sha,
      h(la ↦ Arr ty m (S'a(nat (n - 1)) ↦ applyBinOp oper (the
(S'a (nat n))) (vs[Suc 0 := the (S'a (nat (n - 1))]) !
        Suc 0))))),
      inih,
      [([the (sha (stackC, topf))],
        vs[Suc 0 := the (S'a (nat (n - 1))]),
        safeP, sigSafeMain, the (cdm (l, i)) + 1 + 1 + 1 + 1 + 1
+ 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1, ag, bd)])])
    -jvm→
      (None, sha,
        h(la ↦ Arr ty m (S'a(nat (n - 1)) ↦ applyBinOp oper (the
(S'a (nat n))) (vs[Suc 0 := the (S'a (nat (n - 1))]) !
          Suc 0))))),
        inih,
        [([Intg 1,
          the (sha (stackC, topf))],
          vs[Suc 0 := the (S'a (nat (n - 1))]),
          safeP, sigSafeMain, the (cdm (l, i)) + 1 + 1 + 1 + 1 + 1
+ 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1, ag, bd)])])
apply (unfold exec-all-def)

```

```

apply (erule rtrancl-trans)
prefer 2
apply (rule r-into-rtrancl)
apply (clarify)
apply (unfold JVMExec.exec.simps)
apply (unfold Let-def)
apply (subgoal-tac (fst (snd (snd (snd (snd (the (method' (P', safeP) sigSafe-
Main))))))) !
      (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc
(Suc (Suc (the (cdm (l, i)))))))))))))))))) =
      LitPush (Intg 1),simp)
apply (erule nth-via-drop-append)

apply (drule drop-Suc-append)
apply (subgoal-tac P+ (None, sha,
      h(la  $\mapsto$  Arr ty m (S'a(nat (n - 1)  $\mapsto$  applyBinOp oper (the
(S'a (nat n))) (vs[Suc 0 := the (S'a (nat (n - 1))]) !
      Suc 0))),
      inih,
      [([Intg 1,
      the (sha (stackC, topf))],
      vs[Suc 0 := the (S'a (nat (n - 1))],
      safeP, sigSafeMain, the (cdm (l, i)) + 1 + 1 + 1 + 1 + 1
+ 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1, ag, bd)])
      -jvm $\rightarrow$ 
      (None, sha,
      h(la  $\mapsto$  Arr ty m (S'a(nat (n - 1)  $\mapsto$  applyBinOp oper (the
(S'a (nat n))) (vs[Suc 0 := the (S'a (nat (n - 1))]) !
      Suc 0))),
      inih,
      [([Intg (n - 1)],
      vs[Suc 0 := the (S'a (nat (n - 1))],
      safeP, sigSafeMain,
      the (cdm (l, i)) + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 +
1 + 1 + 1 + 1 + 1 + 1 + 1, ag, bd)]))
apply (unfold exec-all-def)
apply (erule rtrancl-trans)
prefer 2
apply (rule r-into-rtrancl)
apply (clarify)
apply (unfold JVMExec.exec.simps)
apply (unfold Let-def)
apply (subgoal-tac (fst (snd (snd (snd (snd (the (method' (P', safeP) sigSafe-
Main))))))) !
      (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc
(Suc (Suc (Suc (the (cdm (l, i)))))))))))))))))) =
      BinOp Subtract, simp)
apply (erule nth-via-drop-append)

```

```

apply (drule drop-Suc-append)
apply (subgoal-tac P ⊢ (None, sha,
      h(la ⊢ Arr ty m (S'a(nat (n - 1) ⊢ applyBinOp oper (the
(S'a (nat n))) (vs[Suc 0 := the (S'a (nat (n - 1)))]) !
      Suc 0))),
      inih,
      [([Intg (n - 1)],
      vs[Suc 0 := the (S'a (nat (n - 1)))],
      safeP, sigSafeMain,
      the (cdm (l, i) + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 +
1 + 1 + 1 + 1 + 1 + 1 + 1, ag, bd))]
      -jvm→
      (None, sha((stackC, topf) ⊢ Intg (n - 1))),
      h(la ⊢ Arr ty m (S'a(nat (n - 1) ⊢ applyBinOp oper (the
(S'a (nat n))) (vs[Suc 0 := the (S'a (nat (n - 1)))]) !
      Suc 0))),
      inih,
      [([],
      vs[Suc 0 := the (S'a (nat (n - 1)))],
      safeP, sigSafeMain,
      the (cdm (l, i) + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 +
1 + 1 + 1 + 1 + 1 + 1 + 1, ag, bd]))])
apply (unfold exec-all-def)
apply (erule rtrancl-trans)
prefer 2
apply (rule r-into-rtrancl)
apply (clarify)
apply (unfold JVMEexec.exec.simps)
apply (unfold Let-def)
apply (subgoal-tac (fst (snd (snd (snd (snd (the (method' (P', safeP) sigSafe-
Main)))))) !
      (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc
(Suc (Suc (Suc (Suc (the (cdm (l, i)))))))))))))))))) =
      Putstatic topf stackC, simp)
apply (erule nth-via-drop-append)
apply simp

apply (frule nonJumping-Suc-pc)
apply (erule-tac sym [where t=PRIMOP oper])
apply (simp add: nonJumping.simps)
apply (simp add: trInstr.simps)

apply (rule-tac x=hm in exI)
apply (rule-tac x=k in exI)
apply (rule-tac x=k0 in exI)

```

```

apply (rule-tac x=(l,Suc i) in exI)
apply (rule-tac x=Val (execOp oper (IntT ia) (IntT i')) # S' in exI)

apply (rule-tac x=sha((stackC, topf) ↦ Intg (n - 1)) in exI)
apply (rule-tac x=h(la ↦ Arr ty m (S'a(nat (n - 1) ↦ applyBinOp oper (the
(S'a (nat n))))
(vs[Suc 0 := the (S'a (nat (n - 1)))]
! Suc 0)))) in exI)
apply (rule-tac x=inih in exI)
apply (rule-tac x=vs[Suc 0 := the (S'a (nat (n - 1)))] in exI)
apply (rule-tac x=the (cdm (l, i)) + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 +
1 + 1 + 1 + 1 + 1 + 1 in exI)
apply (rule-tac x=ref in exI)

apply (rule-tac x=nat (the-Intg (the ((sha((stackC, topf) ↦ Intg (n - 1)))
(heapC, kf)))) in exI)
apply (rule-tac x=nat (the-Intg (the ((sha((stackC, topf) ↦ Intg (n - 1)))
(heapC, k0f)))) in exI)

apply (rule-tac x=la in exI)
apply (rule-tac x=ty in exI)
apply (rule-tac x=m in exI)
apply (rule-tac x=(S'a(nat (n - 1) ↦ applyBinOp oper (the (S'a (nat n)))
(vs[Suc 0 := the (S'a (nat (n - 1)))] ! Suc 0)))) in exI)

apply (rule-tac x=n - 1 in exI)
apply (rule-tac x=l' in exI)
apply (rule-tac x=regS in exI)
apply (rule-tac x=l'' in exI)
apply (rule-tac x=m' in exI)
apply (rule-tac x=m'' in exI)
apply (rule-tac x=d in exI)
apply (rule-tac x=g in exI)

apply (rule conjI, clarsimp)
apply (rule conjI, rule refl)
apply (rule conjI, rule refl)
apply (rule conjI, rule refl)
apply (rule conjI) apply clarsimp
apply (rule conjI, assumption)
apply (rule conjI) apply clarsimp
apply (simp add: heapC-def add: stackC-def)
apply (rule conjI) apply (rule activeCells,assumption)
apply (rule conjI) apply (rule activeCells-2,assumption+)
apply (rule activeCells-2,assumption+)
apply (rule conjI) apply clarsimp
apply (rule conjI) apply clarsimp
apply (rule conjI) apply clarsimp
apply (rule conjI) apply clarsimp

```



```

apply (rule conjI) apply clarsimp
apply (rule conjI) apply clarsimp
apply (rule conjI) apply clarsimp
apply (rule conjI) apply clarsimp
apply (rule conjI) apply (rule equivH-PRIMOP, assumption+)
apply (rule conjI) apply simp
apply (rule conjI) apply clarsimp
      apply (subgoal-tac length vs = 10)
      apply (rule equivS2-PRIMOP,simp, assumption+)
      apply (rule length-vs)
apply (rule conjI) apply simp
      apply (simp add: heapC-def add: stackC-def)
by simp

```

**end**

**theory** dem-SLIDE

**imports** ../JVMSAFE/JVMExec ../JVMSAFE/State SVM2JVM SVMSemantics  
 CertifSVM2JVM

*dem-translation*

**begin**

**no-translations** Norm s == (None,s)

**no-translations** ex-table-of m == snd (snd (snd m))

**declare** trInstr.simps [simp del]

**declare** equivS.simps [simp del]

**declare** extractBytecode-def [simp del]

**declare** initConsTable-def [simp del]

**constdefs**

*slideArray* :: [entries- ,int,nat,nat] ⇒ entries-

*slideArray* S ntop m n ≡ (  
 let bs = map ( % i . the (S (nat(ntop)-m+i+1))) [0.. $m$ ]  
 in S ( [nat(ntop)-m-n+1.. $nat(ntop)-n+1$ ] [↦] bs))

**axioms** *slide-method*:

[[ s = (None, shp, hp, ihp,  
 ([Intg (int n),Intg (int m)],loc,safeP, sigSafeMain,pc,z1,z2)#[]);

shp (stackC,Sf) = Some (Addr l);

shp (StackC,topf) = Some (Intg ntop);

hp l = Some (Arr ty ma S');

Invoke-static stackC slide [PrimT Integer,PrimT Integer] =

fst(snd(snd(snd(snd(the(method' (P',safeP) sigSafeMain)))))) !

pc

] ⇒

$$P' \mid - s - jvm \rightarrow (None, shp ((stackC, topf) \mapsto Intg (ntop-int\ n)), \\ hp(l \mapsto Arr\ ty\ ma\ (slideArray\ S'\ ntop\ m\ n)), \\ ihp, ([], loc, safeP, sigSafeMain, pc+1, z1, z2) \# [])$$

**declare** *equivH.simps* [*simp del*]

**lemma** *activeCells-SLIDE*:

$$\llbracket la \notin activeCells\ regS\ h\ (nat\ (the-Intg\ (the\ (sh\ (heapC,\ kf)))))) \rrbracket \\ \implies activeCells\ regS\ h\ (nat\ (the-Intg\ (the\ (sh\ (heapC,\ kf)))))) = \\ activeCells\ regS\ (h(la \mapsto Arr\ ty\ ma\ (slideArray\ S'\ na\ m\ n))) \\ (nat\ (the-Intg\ (the\ (sh\ (heapC,\ kf))))))$$

**apply** (*frule l-not-in-cellReg*)  
**apply** (*unfold activeCells-def, auto*)  
**apply** (*unfold region-def, simp add: Let-def, elim conjE*)  
**apply** (*rule-tac x=j in exI, simp*)  
**apply** (*rule cellReg-step*)  
**apply** (*rule cellReg-basic*)  
**apply** (*erule-tac x=j in allE*)  
**apply** (*rule cellsReg-monotone-1, assumption+*)  
**apply** (*erule l-not-in-regs*)  
**apply** (*simp add: Let-def, elim conjE*)  
**apply** (*rule-tac x=j in exI, simp*)  
**apply** (*rule cellReg-step*)  
**apply** (*rule cellReg-basic*)  
**apply** (*erule-tac x=j in allE*)  
**apply** (*erule cellsReg-monotone-2, assumption+*)  
**by** (*erule l-not-in-regs*)

**lemma** *domH-SLIDE*:

$$\llbracket la \notin activeCells\ regS\ h\ (nat\ (the-Intg\ (the\ (sha\ (heapC,\ kf)))))) \rrbracket; \\ \forall l \in dom\ H. \exists l'. l' = the\ (g\ l) \wedge equivC\ (the\ (H\ l))\ h\ l'\ (the\ (h\ l'))\ (nat \\ (the-Intg\ (the\ (sha\ (heapC,\ kf))))))\ com\ regS\ d\ g \rrbracket \\ \implies \forall l \in dom\ H. \\ \exists l'. l' = the\ (g\ l) \wedge \\ equivC\ (the\ (H\ l))\ (h(la \mapsto Arr\ ty\ ma\ (slideArray\ S'\ na\ m\ n)))\ l' \\ (the\ ((h(la \mapsto Arr\ ty\ ma\ (slideArray\ S'\ na\ m\ n)))\ l')) \\ (nat\ (the-Intg\ (the\ (sha\ (heapC,\ kf))))))\ com\ regS\ d\ g$$

**apply** (*frule l-not-in-cellReg*)  
**apply** (*rule ballI*)  
**apply** (*erule-tac x=l in ballE*)  
**prefer** 2 **apply** *simp*  
**apply** (*erule exE*)  
**apply** (*rule-tac x=l' in exI*) **apply** (*elim conjE*)  
**apply** (*rule conjI, assumption*)

**apply** *clarsimp*  
**apply** (*rule conjI*, *rule impI*)  
**apply** (*simp add: region-def*)  
**apply** (*rule impI*)  
**apply** (*rule-tac x=Obj in exI*)  
**apply** (*rule-tac x=flds in exI*) **apply** *simp*  
**apply** (*simp add: region-def*)  
**apply** *clarsimp*  
**apply** (*rule cellReg-step*)  
**apply** (*rule cellReg-basic*)  
**apply** (*erule-tac x=j in allE*)  
**apply** (*rule cellsReg-monotone-1,assumption+*)  
**by** (*erule l-not-in-regs*)

**lemma** *equivH-SLIDE*:

$\llbracket ((hm, k), k0, (l, i), S) = ((H, ka), k0a, PC, Sa);$   
 $k' = \text{nat } (\text{the-Intg } (\text{the } (\text{sha } (\text{heapC}, kf))));$   
 $k0' = \text{nat } (\text{the-Intg } (\text{the } (\text{sha } (\text{heapC}, k0f))));$   
 $\text{sha } (\text{stackC}, Sf) = \text{Some } (\text{Addr } la);$   
 $\text{distinct } [la, l', l'']; h \text{ la} = \text{Some } (\text{Arr ty ma } S');$   
 $\text{sha } (\text{stackC}, \text{topf}) = \text{Some } (\text{Intg } na);$   
 $\text{sha } (\text{heapC}, \text{regionsf}) = \text{Some } (\text{Addr } l');$   
 $h \text{ l}' = \text{Some } (\text{Arr ty m}' \text{ regS});$   
 $\text{sha } (\text{dirCellC}, \text{safeDirf}) = \text{Some } (\text{Addr } l'');$   
 $h \text{ l}'' = \text{Some } (\text{Arr ty m}'' \text{ d}); \text{inj-on } g \text{ (dom } H);$   
 $\text{equivH } (H, ka) \text{ h } k' \text{ com regS } d \text{ g}; k0a = k0';$   
 $\text{activeCells regS } h \text{ k}' \cap \{la, l', l''\} = \{\}$   
 $\implies \text{equivH } (hm, k) \text{ (h(la} \mapsto \text{Arr ty ma (slideArray } S' \text{ na m n))}$   
 $\text{ (nat (the-Intg (the ((sha((stackC, topf)} \mapsto \text{Intg (na - int n)))) (heapC,$   
 $kf)))) \text{ com regS } d \text{ g}$   
**apply** (*simp add: heapC-def add: stackC-def*)  
**apply** (*fold heapC-def*)  
**apply** (*unfold equivH.simps, elim conjE*)  
**apply** (*rule conjI, clarsimp*)  
**apply** (*rule conjI, frule activeCells-SLIDE, simp*)  
**apply** (*rule conjI, assumption*)  
**by** (*rule domH-SLIDE, assumption+*)

**axioms** *equivS-concat*:

$\llbracket \text{equivS } (S1@S2) \text{ S}' \text{ ntop ctm } d \text{ g};$   
 $n = \text{length } S1;$   
 $\forall i < n . \forall z . S!i \neq \text{Cont } z;$   
 $\forall i < n . S' (\text{nat ntop} - i) = S'' (\text{nat ntop} - i);$   
 $\text{equivS } S3 \text{ S}'' (\text{ntop} - \text{int } n) \text{ ctm } d \text{ g}$   
 $\rrbracket \implies \text{equivS } (S1@S3) \text{ S}'' \text{ ntop ctm } d \text{ g}$

**axioms** *equivS-concat2*:

$\llbracket \text{equivS } (S1@S2) \text{ S}' \text{ ntop ctm } d \text{ g};$

$n = \text{length } S1;$   
 $\forall i < n . \forall z . S!i \neq \text{Cont } z$   
 $\] \Longrightarrow \text{equivS } S2 \ S' \ (\text{ntop} - \text{int } n) \ \text{ctm } d \ g$

**axioms** *equivS-same-subarray*:

$\text{equivS} \ (\text{drop} \ (m + n) \ S)$   
 $\quad (S'([\text{nat } \text{ntop} - m - n + 1 .. < \text{nat } \text{ntop} - n + 1] \ [\mapsto]$   
 $\quad \text{map} \ (\lambda i. \ \text{the} \ (S' \ (\text{nat } \text{ntop} - m + i + 1))) \ [0 .. < m]))$   
 $\quad (\text{ntop} - \text{int } n - \text{int } m) \ \text{ctm } d \ g$   
 $\Longrightarrow \text{equivS} \ (\text{drop} \ (\text{Suc } m + n) \ S)$   
 $\quad (S'([\text{nat } \text{ntop} - \text{Suc } m - n + 1 .. < \text{nat } \text{ntop} - n + 1] \ [\mapsto]$   
 $\quad \text{map} \ (\lambda i. \ \text{the} \ (S' \ (\text{nat } \text{ntop} - \text{Suc } m + i + 1))) \ [0 .. < \text{Suc } m]))$   
 $\quad (\text{ntop} - \text{int } n - \text{int } m - 1) \ \text{ctm } d \ g$

**axioms** *equivS-nth*:

$\[ \text{equivS} \ S \ S' \ \text{ntop} \ \text{ctm } d \ g;$   
 $\quad n \leq \text{length } S;$   
 $\quad \forall i < n . \forall z . S!i \neq \text{Cont } z;$   
 $\quad i < n;$   
 $\quad S!i = \text{Val } v \] \Longrightarrow \text{equivV } v \ (\text{the} \ (S' \ (\text{nat } \text{ntop} - i))) \ d \ g$

**axioms** *equivS-nth-reg*:

$\[ \text{equivS} \ S \ S' \ \text{ntop} \ \text{ctm } d \ g;$   
 $\quad n \leq \text{length } S;$   
 $\quad \forall i < n . \forall z . S!i \neq \text{Cont } z;$   
 $\quad i < n;$   
 $\quad S!i = \text{Reg } \text{reg} \] \Longrightarrow \text{reg} = \text{nat} \ (\text{the-Intg} \ (\text{the} \ (S' \ (\text{nat } \text{ntop} - i))))$

**lemma** *equivS-SLIDE* [rule-format]:

$\text{equivS} \ S \ S' \ \text{ntop} \ \text{ctm } d \ g \longrightarrow$   
 $\text{length } S \geq m + n \longrightarrow$   
 $(\forall i < m+n . \forall z . S!i \neq \text{Cont } z) \longrightarrow$   
 $\text{equivS} \ (\text{take } m \ S \ @ \ \text{drop} \ (m + n) \ S)$   
 $\quad (\text{slideArray } S' \ \text{ntop} \ m \ n) \ (\text{ntop} - \text{int } n) \ \text{ctm } d \ g$

**apply** (*unfold slideArray-def*)

**apply** (*unfold Let-def*)

**apply** (*induct m*)

**apply** *simp*

**apply** (*induct-tac n*)

**apply** *simp*

**apply** (*rule impI*) $+$

**apply** (*drule mp, simp*) $+$

**apply** (*subgoal-tac drop n S = S!n # drop (Suc n) S*)

**apply** *clarsimp*

**apply** (*thin-tac equivS ?x ?y ?z ?u ?v ?w*)

```

apply (case-tac S!n)
apply (simp add: equivS.simps)
apply (elim conjE)
apply (subgoal-tac ntop - int n - 1=ntop - (1 + int n),simp,simp)
apply (simp add: equivS.simps)
apply (elim conjE)
apply (subgoal-tac ntop - int n - 1=ntop - (1 + int n),simp,simp)
apply (erule-tac x=n in allE)
apply (drule mp,simp,force)
apply (rule drop-nth3,simp)

```

```

apply (rule impI)+
apply (drule mp,simp)+
apply (subgoal-tac take m S @ (S!m # drop (Suc m + n) S) =
      take (Suc m) S @ drop (Suc m + n) S)
apply (erule-tac s=take m S @ (S!m # drop (Suc m + n) S) in subst)
prefer 2
apply (subgoal-tac (take m S @ [S ! m]) @ drop (Suc m + n) S =
      take m S @ S ! m # drop (Suc m + n) S)
prefer 2 apply (rule sym,rule concat1)
apply (erule subst)
apply (rule concat2)
apply (rule sym,rule take-append3,simp)

```

```

apply (subgoal-tac ntop = int (length S) - 1)
prefer 2 apply (erule equivS-length)
apply (rule equivS-concat)
apply (simp del: upt-Suc)
apply (subgoal-tac m=length (take m S))
prefer 2 apply (rule take-length,simp)
apply simp
apply (rule allI,rule impI,erule-tac x=i in allE,drule mp,simp,simp)
defer

```

```

apply (case-tac S!m)

```

```

apply (simp only: equivS.simps)
apply (rule conjI,simp)
apply (rule conjI) defer
apply (subgoal-tac equivS (drop (m + n) S)
      (S'([nat ntop - m - n + 1..prefer 2 apply (rule equivS-concat2) apply simp
apply (rule take-length,simp)

```

```

apply (rule allI,rule impI,erule-tac x=i in allE,drule mp,simp,simp)
apply clarify
apply (erule equivS-same-subarray)

```

```

apply (simp only: equivS.simps)
apply (rule conjI,simp)
apply (rule conjI) defer
apply (subgoal-tac equivS (drop (m + n) S)
  (S'([nat ntop - m - n + 1..<nat ntop - n + 1] [↦])
  map (λi. the (S' (nat ntop - m + i + 1))) [0..<m]))
  ((ntop - int n) - int m) ctm d g)
prefer 2 apply clarify apply (erule equivS-concat2)
apply (rule take-length,simp)
apply (rule allI,rule impI,erule-tac x=i in allE,drule mp,simp,simp)
apply clarify
apply (erule equivS-same-subarray)

```

```

apply (erule-tac x=m in allE)
apply (drule mp,simp)
apply (erule-tac x=x in allE)
apply force

```

```

apply (rule allI,rule impI)
apply (subgoal-tac (Suc (nat ntop - (m + n))) + ((m-i) - 1) = nat (ntop - int
n) - i)
apply (erule-tac s=Suc (nat ntop - (m + n)) + ((m-i) - 1) in subst)
apply (subgoal-tac (S'([Suc (nat ntop - (m + n))..<Suc (nat ntop - n)] [↦])
  map (λi. the (S' (Suc (nat ntop - m + i)))) [0..<m]))
  (Suc (nat ntop - (m + n)) + ((m - i) - 1)) =
  Some ((map (λi. the (S' (Suc (nat ntop - m + i)))) [0..<m])!((m -
i) - 1)))
prefer 2 apply (rule map-upds-nth)
apply simp
apply clarify
apply (rule-tac s=Some (map (λi. the (S' (Suc (nat (int (length S) - 1) - m +
i)))) [0..<m]
  ! (m - i - 1)) in subst)
prefer 2 apply simp
apply (thin-tac ?x=?y)
apply (subgoal-tac (nat (int (length S) - 1) - Suc m - n + 1) + (m-i) =
  Suc (nat (int (length S) - 1) - (m + n)) + (m - i - 1))
apply (erule-tac s=nat (int (length S) - 1) - Suc m - n + 1 + (m - i) in
subst)
apply (subgoal-tac (S'([nat (int (length S) - 1) - Suc m - n + 1..<nat (int
(length S) - 1) - n + 1]

```

```

      [↦] map (λi. the (S' (nat (int (length S) - 1) - Suc m + i + 1)))
[0..<Suc m]))
      (nat (int (length S) - 1) - Suc m - n + 1 + (m - i))=
      Some ((map (λi. the (S' (nat (int (length S) - 1) - Suc m + i +
1)))) [0..<Suc m])
      ! (m - i)))
prefer 2 apply (rule map-upds-nth)
apply (erule additions1,simp)
apply (rule-tac s=Some (map (λi. the (S' (nat (int (length S) - 1) - Suc m +
i + 1)))) [0..<Suc m] !
      (m - i)) and
      t=(S'([nat (int (length S) - 1) - Suc m - n + 1..<nat (int (length S) -
1) - n + 1] [↦]
      map (λi. the (S' (nat (int (length S) - 1) - Suc m + i + 1))) [0..<Suc
m]))
      (nat (int (length S) - 1) - Suc m - n + 1 + (m - i)) in subst)
apply (rule sym,assumption)
apply (thin-tac ?x=?y)
apply (simp only: nth-map-upt)
apply (subgoal-tac Suc (nat (int (length S) - 1) - m + (0 + (m - i - 1)))=
      nat (int (length S) - 1) - Suc m + (0 + (m - i)) + 1)
apply (erule-tac s=Suc (nat (int (length S) - 1) - m + (0 + (m - i - 1))) in
subst,rule refl)

apply (simp (no-asm))
apply (erule additions2,simp)
apply (simp (no-asm))
apply (erule additions3,simp)
apply (simp (no-asm))

apply (subgoal-tac (nat (int (length S) - 1) - Suc m - n + 1) + 0
      = nat (int (length S) - 1 - int n - int m))
apply (erule-tac s=(nat (int (length S) - 1) - Suc m - n + 1) + 0 in subst)
apply (subgoal-tac
      (S'([nat (int (length S) - 1) - Suc m - n + 1..<nat (int (length S) - 1)
- n + 1] [↦]
      map (λi. the (S' (nat (int (length S) - 1) - Suc m + i + 1))) [0..<Suc
m]))
      (nat (int (length S) - 1) - Suc m - n + 1 + 0)=
      Some (map (λi. the (S' (nat (int (length S) - 1) - Suc m + i + 1)))
[0..<Suc m] ! 0))
prefer 2 apply (rule map-upds-nth) apply simp
defer
apply (rule-tac t=(S'([nat (int (length S) - 1) - Suc m - n + 1..<nat (int
(length S) - 1) - n + 1] [↦]
      map (λi. the (S' (nat (int (length S) - 1) - Suc m + i + 1))) [0..<Suc
m]))

```

```

      (nat (int (length S) - 1) - Suc m - n + 1 + 0) and
      s=Some (map (λi. the (S' (nat (int (length S) - 1) - Suc m + i + 1)))
[0..<Suc m] ! 0)
    in subst) apply (rule sym,assumption)
  apply (thin-tac ?x=Some ?y)
  apply (subgoal-tac
    map (λi. the (S' (nat (int (length S) - 1) - Suc m + i + 1))) [0..<Suc m]
    ! 0
    =(λi. the (S' (nat (int (length S) - 1) - Suc m + i + 1))) (0+0))
  prefer 2 apply (rule nth-map-upt, simp (no-asm))
  apply (rule-tac s=the (S' (nat (int (length S) - 1) - Suc m + (0 + 0) + 1))
    and t=map (λi. the (S' (nat (int (length S) - 1) - Suc m + i + 1)))
[0..<Suc m] ! 0)
    in subst,rule sym,assumption)
  apply (rule-tac s=nat ntop - m and t=nat (int (length S) - 1) - Suc m + (0
+ 0) + 1 in subst)
  apply (simp (no-asm)) apply clarify
  apply (erule additions4)
  apply (simp (no-asm),clarify)
  apply (erule-tac n=Suc m + n in equivS-nth)
  apply (assumption,rule allI,rule impI,force,simp,assumption)
  apply (simp (no-asm))
  apply (erule additions5)
  prefer 2 apply (rule additions6,simp)

  apply (subgoal-tac (nat (int (length S) - 1) - Suc m - n + 1) + 0
    = nat (int (length S) - 1 - int n - int m))
  apply (erule-tac s=(nat (int (length S) - 1) - Suc m - n + 1) + 0 in subst)
  apply (subgoal-tac
    (S'([nat (int (length S) - 1) - Suc m - n + 1..<nat (int (length S) - 1)
- n + 1] [↦]
    map (λi. the (S' (nat (int (length S) - 1) - Suc m + i + 1))) [0..<Suc
m]))
    (nat (int (length S) - 1) - Suc m - n + 1 + 0)=
    Some (map (λi. the (S' (nat (int (length S) - 1) - Suc m + i + 1)))
[0..<Suc m] ! 0))
  prefer 2 apply (rule map-upds-nth) apply simp
  defer
  apply (rule-tac t=(S'([nat (int (length S) - 1) - Suc m - n + 1..<nat (int
(length S) - 1) - n + 1] [↦]
    map (λi. the (S' (nat (int (length S) - 1) - Suc m + i + 1))) [0..<Suc
m]))
    (nat (int (length S) - 1) - Suc m - n + 1 + 0) and
    s=Some (map (λi. the (S' (nat (int (length S) - 1) - Suc m + i + 1)))
[0..<Suc m] ! 0)
    in subst) apply (rule sym,assumption)
  apply (thin-tac ?x=Some ?y)
  apply (subgoal-tac

```



$\text{map } (\lambda i. \text{the } (S' (\text{nat } (\text{length } S) - 1) - \text{Suc } m + i + 1))) [0..<\text{Suc } m]$   
 $! 0$   
 $= (\lambda i. \text{the } (S' (\text{nat } (\text{length } S) - 1) - \text{Suc } m + i + 1))) (0+0)$   
**prefer 2 apply** (rule nth-map-upt, simp (no-asm))  
**apply** (rule-tac s=the (S' (nat (length S) - 1) - Suc m + (0 + 0) + 1))  
**and** t=map ( $\lambda i. \text{the } (S' (\text{nat } (\text{length } S) - 1) - \text{Suc } m + i + 1)))$   
 $[0..<\text{Suc } m] ! 0$   
**in** subst,rule sym,assumption)  
**apply** (rule-tac s=nat ntop - m **and** t=nat (int (length S) - 1) - Suc m + (0  
+ 0) + 1 **in** subst)  
**apply** (simp (no-asm)) **apply** clarify  
**apply** (erule additions4)  
**apply** (simp (no-asm),clarify)  
**apply** (erule-tac n=Suc m + n **in** equivS-nth-reg)  
**apply** (assumption,rule allI,rule impI,force,simp,assumption)  
**apply** (simp (no-asm))  
**apply** (erule additions5)  
**apply** (rule additions6,simp)  
**done**

**axioms** good-stack-SLIDE:

$\llbracket (P', \text{cdm}, \text{ctm}, \text{com}) = \text{trSVM2JVM } ((\text{svms}, \text{ctmap}), \text{ini}, \text{ct}, \text{ah}, \text{ai}, \text{bc});$   
 $\text{cdm}, \text{ctm}, \text{com} \vdash ((\text{hm}, \text{k}), \text{k0}, (\text{l}, \text{i}), \text{S}) \triangleq \text{S1}';$   
 $(\text{fst } (\text{the } (\text{map-of } \text{svms } \text{l}))) ! \text{i} = \text{SLIDE } m \text{ n}$   
 $\rrbracket \implies \text{length } S \geq m + n \wedge (\forall i < m+n . \forall z . S!i \neq \text{Cont } z)$

**lemma** execSVMInstr-SLIDE :

$\llbracket (P', \text{cdm}, \text{ctm}, \text{com}) = \text{trSVM2JVM } ((\text{svms}, \text{ctmap}), \text{ini}, \text{ct}, \text{ah}, \text{ai}, \text{bc});$   
 $\text{cdm}, \text{ctm}, \text{com} \vdash ((\text{hm}, \text{k}), \text{k0}, (\text{l}, \text{i}), \text{S}) \triangleq \text{S1}';$   
 $(\text{fst } (\text{the } (\text{map-of } \text{svms } \text{l}))) ! \text{i} = \text{SLIDE } m \text{ n};$   
 $\text{execSVMInst } (\text{SLIDE } m \text{ n}) (\text{map-of } \text{ct}) (\text{hm}, \text{k}) \text{k0 } (\text{l}, \text{i}) \text{S} = \text{Either.Right}$   
 $\text{S2};$   
 $\text{drop } (\text{the } (\text{cdm } (\text{l}, \text{i}))) (\text{extractBytecode } P') =$   
 $\text{trInstr } (\text{the } (\text{cdm } (\text{l}, \text{i}))) \text{cdm}' \text{ctm}' \text{com } \text{pcc } (\text{SLIDE } m \text{ n}) @ \text{bytecode}'$   
 $\rrbracket \implies \exists v' \text{sh}' \text{dh}' \text{ih}' \text{fms}' .$   
 $P' \vdash \text{S1}' - \text{jvm} \rightarrow (v', \text{sh}', \text{dh}', \text{ih}', \text{fms}') \wedge$   
 $\text{cdm}, \text{ctm}, \text{com} \vdash \text{S2} \triangleq (v', \text{sh}', \text{dh}', \text{ih}', \text{fms}') \wedge$   
 $\text{cdm}, \text{ctm}, \text{com} \vdash \text{S2} \triangleq (v', \text{sh}', \text{dh}', \text{ih}', \text{fms}')$

**apply** (subgoal-tac length S  $\geq m + n \wedge (\forall i < m+n . \forall z . S!i \neq \text{Cont } z)$ )  
**prefer 2 apply** (rule good-stack-SLIDE,assumption+)  
**apply** (elim conjE)

**apply** (case-tac S1')  
**apply** (rename-tac v tup)  
**apply** (case-tac tup)

```

apply (rename-tac sh tup)
apply (case-tac tup)
apply (rename-tac dh tup)
apply (case-tac tup)
apply (rename-tac ih fms)
apply (simp)

```

```

apply (unfold equivState-def)
apply (elim exE,elim conjE)

```

```

apply (rule-tac x=None in exI)
apply (rule-tac x=sha((stackC, topf) ↦ Intg (na - int(n))) in exI)
apply (rule-tac x=h(la ↦ Arr ty ma (slideArray S' na m n)) in exI)
apply (rule-tac x=ih in exI)
apply (rule-tac x=[[[], vs, safeP, sigSafeMain, the (cdm (l, i)) + 1 + 1 + 1,
ref]] in exI)
apply (rule conjI)
apply (unfold extractBytecode-def)
apply (clarify)

```

```

apply (subgoal-tac P⊢ (None,sha,h,inih,[[[], vs, safeP, sigSafeMain, the (cdm
(l,i), ag,bd]])
  -jvm→ (None,sha,h,inih,[[[Intg (int m)], vs, safeP, sigSafeMain, the (cdm (l,i)
+ 1, ag,bd]]))
apply (unfold exec-all-def )
apply (erule rtrancl-trans)

```

```

prefer 2
apply (rule r-into-rtrancl)
apply (clarify)
apply (unfold JVMEexec.exec.simps)
apply (unfold Let-def)
apply (subgoal-tac (fst (snd (snd (snd (snd (the (method' (P', safeP) sigSafe-
Main))))))) !
  the (cdm (l, i))) = LitPush (Intg (int m)),simp)

```

```

apply (unfold trInstr.simps)
apply (erule nth-via-drop-append)

```

```

apply (drule drop-Suc-append)
apply (subgoal-tac P⊢ (None,sha,h,inih,[[[Intg (int m)],
vs, safeP, sigSafeMain, the (cdm (l,i))+1, ag,bd]])
  -jvm→ (None,sha,h,inih,[[[Intg (int n), Intg (int m)],

```

```

vs, safeP, sigSafeMain, (the (cdm (l,i) + 1)+1, ag,bd)))
apply (unfold exec-all-def )
apply (erule rtrancl-trans)
prefer 2
apply (rule r-into-rtrancl)
apply (clarify)
apply (unfold JVMExec.exec.simps)
apply (unfold Let-def)
apply (subgoal-tac (fst (snd (snd (snd (snd (the (method' (P', safeP) sigSafe-
Main))))))) !
Suc (the (cdm (l, i)))) = LitPush (Intg (int n))) apply (simp)
apply (erule nth-via-drop-append)

apply (drule drop-Suc-append)
apply (fold exec-all-def)
apply (rule slide-method)
apply (simp,assumption,simp,simp)

apply (rule sym)
apply (rule-tac ys=[] in nth-via-drop-append)
apply simp

apply (frule nonJumping-Suc-pc)
apply (erule-tac sym [where t=SLIDE m n])
apply (simp add: nonJumping.simps)
apply (simp add: trInstr.simps)

apply (rule-tac x=hm in exI)
apply (rule-tac x=k in exI)
apply (rule-tac x=k0 in exI)
apply (rule-tac x=(l,Suc i) in exI)
apply (rule-tac x=take m S @ drop (m + n) S in exI)

apply (rule-tac x=sha((stackC, topf)  $\mapsto$  Intg (na - int n)) in exI)
apply (rule-tac x=h(la  $\mapsto$  Arr ty ma (slideArray S' na m n)) in exI)
apply (rule-tac x=ih in exI)
apply (rule-tac x=vs in exI)
apply (rule-tac x=the (cdm (l, i) + 1 + 1 + 1) in exI)
apply (rule-tac x=ref in exI)

apply (rule-tac x=nat (the-Intg (the ((sha((stackC, topf)  $\mapsto$  Intg (na - int n)))
(heapC, kf)))) in exI)
apply (rule-tac x=nat (the-Intg (the ((sha((stackC, topf)  $\mapsto$  Intg (na - int n)))
(heapC, k0f)))) in exI)

apply (rule-tac x=la in exI)

```

```

apply (rule-tac x=ty in exI)
apply (rule-tac x=ma in exI)
apply (rule-tac x=(slideArray S' na m n) in exI)

apply (rule-tac x=(na - int n) in exI)
apply (rule-tac x=l' in exI)
apply (rule-tac x=regS in exI)
apply (rule-tac x=l'' in exI)
apply (rule-tac x=m' in exI)
apply (rule-tac x=m'' in exI)
apply (rule-tac x=d in exI)
apply (rule-tac x=g in exI)

apply (rule conjI, clarsimp)
apply (rule conjI, rule refl)
apply (rule conjI, rule refl)
apply (rule conjI, rule refl)
apply (rule conjI) apply clarsimp
apply (rule conjI, assumption)
apply (rule conjI) apply clarsimp
      apply (simp add: heapC-def add: stackC-def)
apply (rule conjI) apply (rule activeCells,assumption)
apply (rule conjI) apply (rule activeCells-2,assumption+)
      apply (rule activeCells-2,assumption+)
apply (rule conjI) apply clarsimp
apply (rule conjI) apply clarsimp
apply (rule conjI) apply clarsimp
apply (rule conjI) apply clarsimp
apply (rule conjI) apply clarsimp
apply (rule conjI) apply clarsimp
apply (rule conjI) apply clarsimp
apply (rule conjI) apply clarsimp
apply (rule conjI) apply (rule equivH-SLIDE, assumption+)
apply (rule conjI) apply simp
apply (rule conjI) apply (rule equivS-SLIDE)
      apply (simp,assumption,erule-tac x=ia in alle,drule mp,assumption,case-tac
z,force)
apply (rule conjI) apply simp
      apply (simp add: heapC-def add: stackC-def)
by simp

end

theory dem-MATCHN
imports ../JVMSAFE/JVMEExec SVM2JVM SVMSemantics CertifSVM2JVM
begin

```

**no-translations**  $Norm\ s == (None, s)$

**no-translations**  $ex\text{-}table\text{-}of\ m == snd\ (snd\ (snd\ m))$

**declare**  $trInstr.simps\ [simp\ del]$

**declare**  $equivS.simps\ [simp\ del]$

**axioms**  $equivS\text{-}MATCHN$ :

$n \geq int\ off$

**declare**  $extractBytecode\text{-}def\ [simp\ del]$

**declare**  $initConsTable\text{-}def\ [simp\ del]$

**constdefs**  $if'\ :: [bool, 'a, 'a] => 'a$

$if'\ c\ a\ b \equiv if\ c\ then\ a\ else\ b$

**axioms**  $pc\text{-}MATCHN$ :

$if'\ (the\ Intg\ (Intg\ (the\ Intg\ (the\ (S'\ (nat\ (n - int\ off)))) - int\ v)) < 0 \vee$   
 $int\ (m + 1) < the\ Intg\ (Intg\ (the\ Intg\ (the\ (S'\ (nat\ (n - int\ off))))$   
 $- int\ v))$   
 $(Suc\ (Suc\ (Suc\ (Suc\ (Suc\ (Suc\ (the\ (cdm\ (l, i)) +$   
 $nat\ (map\ (\lambda n. trAddr\ n\ (the\ (cdm\ (l, i)) + incMatchN))\ (map\ (\lambda p. the$   
 $(cdm'\ (p, 0))\ ps)\ !$   
 $nat\ (int\ (2 + m))))))))))$   
 $(Suc\ (Suc\ (Suc\ (Suc\ (Suc\ (Suc\ (the\ (cdm\ (l, i)) +$   
 $nat\ (map\ (\lambda n. trAddr\ n\ (the\ (cdm\ (l, i)) + incMatchN))\ (map\ (\lambda p. the$   
 $(cdm'\ (p, 0))\ ps)\ !$   
 $nat\ (the\ Intg\ (Intg\ (the\ Intg\ (the\ (S'\ (nat\ (n - int\ off)))) - int\ v)) -$   
 $0)))))))))) =$   
 $the\ (cdm\ (case\ S\ !+ off\ of\ Val\ (IntT\ i) \Rightarrow let\ r = nat\ i - v; p = if\ 0 \leq r \wedge r \leq$   
 $m - 1\ then\ ps\ !\ r\ else\ ps\ !\ m\ in\ (p, 0)$   
 $| Val\ (BoolT\ b) \Rightarrow let\ p = if\ \neg b\ then\ ps\ !\ 0\ else\ ps\ !\ 1\ in\ (p, 0)))$

**declare**  $if'\text{-}def\ [simp\ del]$

**lemma**  $execSVMInstr\text{-}MATCHN$  :

$\llbracket (P', cdm, ctm, com) = trSVM2JVM\ ((svms, ctmap), ini, ct, ah, ai, bc);$   
 $cdm, ctm, com \vdash ((hm, k), k0, (l, i), S) \hat{=} S1';$   
 $(fst\ (the\ (map\text{-}of\ svms\ l))\ !\ i) = MATCHN\ off\ v\ m\ ps;$   
 $execSVMInst\ (MATCHN\ off\ v\ m\ ps)\ (map\text{-}of\ ct)\ (hm, k)\ k0\ (l, i)\ S =$   
 $Either.Right\ S2;$   
 $drop\ (the\ (cdm\ (l, i)))\ (extractBytecode\ P') =$   
 $trInstr\ (the\ (cdm\ (l, i)))\ cdm'\ ctm'\ com\ pcc\ (MATCHN\ off\ v\ m\ ps)\ @$   
 $bytecode';$

$$\begin{aligned}
& cdm' \subseteq_m cdm \\
\mathbb{J} & \Longrightarrow \exists v' sh' dh' ih' fms' . \\
P' \vdash & S1' -jvm \rightarrow (v', sh', dh', ih', fms') \wedge \\
& cdm, ctm, com \vdash S2 \triangleq (v', sh', dh', ih', fms')
\end{aligned}$$

```

apply (case-tac S1')
apply (rename-tac v1 tup)
apply (case-tac tup)
apply (rename-tac sh tup)
apply (case-tac tup)
apply (rename-tac dh tup)
apply (case-tac tup)
apply (rename-tac ih fms)
apply (simp)

```

```

apply (unfold equivState-def)
apply (elim exE, elim conjE)

```

```

apply (rule-tac x=None in exI)
apply (rule-tac x=sha in exI)
apply (rule-tac x=h in exI)
apply (rule-tac x=inh in exI)
apply (rule-tac x=[[], vs,
    safeP, sigSafeMain,
    if' ((the-Intg (Intg (the-Intg (the (S' (nat (n - int off)))))) - int
v)) < 0) ∨
    (the-Intg (Intg (the-Intg (the (S' (nat (n - int off)))))) - int
v)) > (int (m + 1))))
    (Suc (Suc (Suc (Suc (Suc (Suc (Suc (the (cdm (l, i)) +
    nat (map (λn. trAddr n (the (cdm (l, i)) +
incMatchN)) (map (λp. the (cdm' (p, 0))) ps) !
    nat (int (2 + m))))))))))
    (Suc (Suc (Suc (Suc (Suc (Suc (Suc (the (cdm (l, i)) +
    nat (map (λn. trAddr n (the (cdm (l, i)) +
incMatchN)) (map (λp. the (cdm' (p, 0))) ps) !
    (nat (the-Intg (Intg (the-Intg (the (S' (nat (n -
int off)))))) - int v)) - 0))))))))),
    ref]) in exI)

```

```

apply (subgoal-tac ∃ vs pc . fms=([], vs, safeP, sigSafeMain, pc, ref)#[])
prefer 2 apply simp
apply (erule exE)+

```

```

apply (rule conjI)

```

```

apply (unfold exec-all-def)
apply (subgoal-tac  $P \vdash (None, sha, h, inih, [([], vs, safeP, sigSafeMain, the (cdm (l, i)), ag, bd)]) -jvm \rightarrow$ 
      (None, sha, h, inih, [([the (sha (stackC, Sf))], vs, safeP, sigSafeMain,
the (cdm (l, i)) + 1, ag, bd)]))
apply (unfold exec-all-def)
apply (erule rtrancl-trans)

```

```

prefer 2
apply (rule r-into-rtrancl)
apply (clarify)
apply (unfold JVMExec.exec.simps)
apply (unfold Let-def)
apply (unfold extractBytecode-def)
apply (subgoal-tac (fst (snd (snd (snd (snd (the (method' (P', safeP) sigSafe-
Main))))))) !
      the (cdm (l, i))) = Getstatic Sf stackC, simp)
apply (unfold trInstr.simps, unfold Let-def)
apply (erule nth-via-drop-append)

```

```

apply (drule drop-Suc-append)
apply (subgoal-tac  $P \vdash (None, sha, h, inih,$ 
      [([the (sha (stackC, Sf))],
      vs, safeP, sigSafeMain, the (cdm (l, i)) + 1, ag, bd)])
      -jvm \rightarrow
      (None, sha, h, inih,
      [([the (sha (stackC, topf)), the (sha (stackC, Sf))],
      vs, safeP, sigSafeMain, the (cdm (l, i)) + 1 + 1, ag, bd)]))
apply (unfold exec-all-def)
apply (erule rtrancl-trans)
prefer 2
apply (rule r-into-rtrancl)
apply (clarify)
apply (unfold JVMExec.exec.simps)
apply (unfold Let-def)
apply (subgoal-tac (fst (snd (snd (snd (snd (the (method' (P', safeP) sigSafe-
Main))))))) ! Suc (the (cdm (l, i)))) =
      Getstatic topf stackC, simp)
apply (erule nth-via-drop-append)

```

```

apply (drule drop-Suc-append)
apply (subgoal-tac  $P \vdash (None, sha, h, inih,$ 

```

```

[[[the (sha (stackC,topf)),the (sha (stackC, Sf))],
vs, safeP, sigSafeMain, the (cdm (l, i) + 1 + 1, ag, bd)]]
-jvm→
(None,sha,h,inih,
[[[Intg (int off),
the (sha (stackC,topf)),the (sha (stackC, Sf))],
vs, safeP, sigSafeMain, the (cdm (l, i) + 1 + 1 + 1, ag,
bd]]))
apply (unfold exec-all-def)
apply (erule rtrancl-trans)
prefer 2
apply (rule r-into-rtrancl)
apply (clarify)
apply (unfold JVMExec.exec.simps)
apply (unfold Let-def)
apply (subgoal-tac (fst (snd (snd (snd (snd (the (method' (P', safeP) sigSafe-
Main)))))) ! (Suc (Suc (the (cdm (l, i)))))) =
LitPush (Intg (int off)),simp)
apply (erule nth-via-drop-append)

```

```

apply (drule drop-Suc-append)
apply (subgoal-tac P+ (None,sha,h,inih,
[[[Intg (int off),
the (sha (stackC,topf)),the (sha (stackC, Sf))],
vs, safeP, sigSafeMain, the (cdm (l, i) + 1 + 1 + 1, ag, bd)]]
-jvm→
(None,sha,h,inih,
[[[Intg ( n - int off),
the (sha (stackC, Sf))],
vs, safeP, sigSafeMain, the (cdm (l, i) + 1 + 1 + 1 + 1,
ag, bd]]))
apply (unfold exec-all-def)
apply (erule rtrancl-trans)
prefer 2
apply (rule r-into-rtrancl)
apply (clarify)
apply (unfold JVMExec.exec.simps)
apply (unfold Let-def)
apply (subgoal-tac (fst (snd (snd (snd (snd (the (method' (P', safeP) sigSafe-
Main)))))) ! (Suc (Suc (Suc (the (cdm (l, i)))))) =
BinOp Subtract,simp)
apply (erule nth-via-drop-append)

```

```

apply (drule drop-Suc-append)
apply (subgoal-tac P+ (None,sha,h,inih,

```



```

      [[([Intg (n - int off),
         the (sha (stackC, Sf))],
         vs, safeP, sigSafeMain, the (cdm (l, i)) + 1 + 1 + 1 + 1,
ag, bd))]
      -jvm→
      (None,sha,h,inih,
      [[the (S' (nat (n - int off)))]],
      vs, safeP, sigSafeMain, the (cdm (l, i)) + 1 + 1 + 1 + 1 +
1, ag, bd]]))
apply (unfold exec-all-def)
apply (erule rtrancl-trans)
prefer 2
apply (rule r-into-rtrancl)
apply (clarify)
apply (unfold JVMExec.exec.simps)
apply (unfold Let-def)
apply (subgoal-tac (fst (snd (snd (snd (snd (the (method' (P', safeP) sigSafe-
Main))))))) !
      (Suc (Suc (Suc (Suc (the (cdm (l, i))))))) =
      ArrLoad)

apply (case-tac S !+ off, simp)
apply (case-tac Vala)
apply (simp, insert RightNotUndefined, erule-tac x = S2 in allE, force)
apply simp apply (rename-tac n1)
apply (subgoal-tac n >= int off) prefer 2 apply (rule equivS-MATCHN,
simp add: raise-system-xcpt-def)
apply simp apply (subgoal-tac n >= int off) prefer 2 apply (rule equivS-MATCHN,
simp add: raise-system-xcpt-def)
apply (simp, insert RightNotUndefined, erule-tac x = S2 in allE, force)
apply (simp, insert RightNotUndefined, erule-tac x = S2 in allE, force)
apply (erule nth-via-drop-append)

apply (drule drop-Suc-append)
apply (subgoal-tac P ⊢ (None,sha,h,inih,
      [[the (S' (nat (n - int off)))]],
      vs, safeP, sigSafeMain, the (cdm (l, i)) + 1 + 1 + 1 + 1 +
1, ag, bd))]
      -jvm→
      (None,sha,h,inih,
      [[([Intg (int v),
         the (S' (nat (n - int off)))]],
      vs, safeP, sigSafeMain, the (cdm (l, i)) + 1 + 1 + 1 + 1 +
1 + 1, ag, bd]]))
apply (unfold exec-all-def)
apply (erule rtrancl-trans)
prefer 2
apply (rule r-into-rtrancl)

```

**apply** (*clarify*)  
**apply** (*unfold JVMExec.exec.simps*)  
**apply** (*unfold Let-def*)  
**apply** (*subgoal-tac* (*fst* (*snd* (*snd* (*snd* (*snd* (*the* (*method'* (*P'*, *safeP*) *sigSafeMain*))))))))) !  

$$(Suc (Suc (Suc (Suc (Suc (the (cdm (l, i)))))))))) =$$

$$LitPush (Intg (int v)), simp)$$
**apply** (*erule nth-via-drop-append*)

**apply** (*drule drop-Suc-append*)  
**apply** (*subgoal-tac*  $P \vdash$  (*None, sha, h, inih*,  
 $[[[Intg (int v),$   
 $the (S' (nat (n - int off))]]],$   
 $vs, safeP, sigSafeMain, the (cdm (l, i)) + 1 + 1 + 1 + 1 +$   
 $1 + 1, ag, bd]]$ )  
 $-jvm \rightarrow$   
 $(None, sha, h, inih,$   
 $[[[Intg (the-Intg (the (S' (nat (n - int off)))) - int v)],$   
 $vs, safeP, sigSafeMain, the (cdm (l, i)) + 1 + 1 + 1 + 1 +$   
 $1 + 1 + 1, ag, bd]]])$ )  
**apply** (*unfold exec-all-def*)  
**apply** (*erule rtrancl-trans*)  
**prefer** 2  
**apply** (*rule r-into-rtrancl*)  
**apply** (*clarify*)  
**apply** (*unfold JVMExec.exec.simps*)  
**apply** (*unfold Let-def*)  
**apply** (*subgoal-tac* (*fst* (*snd* (*snd* (*snd* (*snd* (*the* (*method'* (*P'*, *safeP*) *sigSafeMain*))))))))) !  

$$(Suc (Suc (Suc (Suc (Suc (Suc (the (cdm (l, i)))))))))) =$$

$$BinOp Subtract, simp)$$
**apply** (*erule nth-via-drop-append*)

**apply** (*drule drop-Suc-append*)  
**apply** (*subgoal-tac*  $P \vdash$  (*None, sha, h, inih*,  
 $[[[Intg (the-Intg (the (S' (nat (n - int off)))) - int v)],$   
 $vs, safeP, sigSafeMain, the (cdm (l, i)) + 1 + 1 + 1 + 1 +$   
 $1 + 1 + 1, ag, bd]]$ )  
 $-jvm \rightarrow$   
 $(None, sha, h, inih,$   
 $[[[[],$   
 $vs, safeP, sigSafeMain,$   
 $if' ((the-Intg (Intg (the-Intg (the (S' (nat (n - int off)))) - int$   
 $v)) < 0) \vee$   
 $(the-Intg (Intg (the-Intg (the (S' (nat (n - int off)))) - int$   
 $v)) > (int (m + 1)))]])$ )

```

      (Suc (Suc (Suc (Suc (Suc (Suc (Suc (the (cdm (l, i)) +
        nat (map (λn. trAddr n (the (cdm (l, i)) +
incMatchN)) (map (λp. the (cdm' (p, 0))) ps) !
        nat (int (2 + m ))))))))))))
      (Suc (Suc (Suc (Suc (Suc (Suc (Suc (the (cdm (l, i)) +
        nat (map (λn. trAddr n (the (cdm (l, i)) +
incMatchN)) (map (λp. the (cdm' (p, 0))) ps) !
        (nat (the-Intg (Intg (the-Intg (the (S' (nat (n -
int off)))) - int v)) - 0)))))))))),
      ag, bd]))
apply (unfold exec-all-def)
apply (erule rtrancl-trans)
prefer 2
apply (rule r-into-rtrancl)
apply (clarify)
apply (unfold JVMExec.exec.simps)
apply (unfold Let-def)
apply (subgoal-tac (fst (snd (snd (snd (snd (the (method' (P', safeP) sigSafe-
Main)))))) !
      (Suc (Suc (Suc (Suc (Suc (Suc (Suc (the (cdm (l, i)))))))))) =
      Tableswitch 0 (int (m + 1)) (map (λn. trAddr n (the (cdm (l,
i)) + incMatchN))
      (map (λp. the (cdm' (p, 0))) ps)),simp)

apply (simp add: if'-def)
apply (erule nth-via-drop-append)
apply simp

apply (rule-tac x=hm in exI)
apply (rule-tac x=k in exI)
apply (rule-tac x=k0 in exI)
apply (rule-tac x=(case S !+ off of
  Val (IntT i) ⇒ let r = nat i - v;
    p = if 0 ≤ r ∧ r ≤ m - 1 then ps ! r else ps ! m
    in (p,0)
  | Val (BoolT b) ⇒ let p = if ¬ b then ps ! 0 else ps ! 1
    in (p,0)) in exI)
apply (rule-tac x=S in exI)

apply (rule-tac x=sha in exI)
apply (rule-tac x=h in exI)
apply (rule-tac x=inih in exI)
apply (rule-tac x=vs in exI)
apply (rule-tac x=if' ((the-Intg (Intg (the-Intg (the (S' (nat (n - int off)))) -
int v)) < 0) ∨
  (the-Intg (Intg (the-Intg (the (S' (nat (n - int off)))) - int
v)) > (int (m + 1))))
  (Suc (Suc (Suc (Suc (Suc (Suc (the (cdm (l, i)) +
    nat (map (λn. trAddr n (the (cdm (l, i)) +

```

```

incMatchN)) (map (λp. the (cdm' (p, 0)) ps) !
              nat (int (2 + m )))))))))))
      (Suc (Suc (Suc (Suc (Suc (Suc (Suc (the (cdm (l, i)) +
              nat (map (λn. trAddr n (the (cdm (l, i)) +
incMatchN)) (map (λp. the (cdm' (p, 0)) ps) !
              (nat (the-Intg (Intg (the-Intg (the (S' (nat (n -
int off)))) - int v)) - 0))))))))) in exI)
apply (rule-tac x=ref in exI)

apply (rule-tac x=nat (the-Intg (the (sha (heapC, kf)))) in exI)

apply (rule-tac x= nat (the-Intg (the (sha (heapC, kOf)))) in exI)

apply (rule-tac x=la in exI)
apply (rule-tac x=ty in exI)
apply (rule-tac x=ma in exI)
apply (rule-tac x=S' in exI)

apply (rule-tac x=n in exI)
apply (rule-tac x=l' in exI)
apply (rule-tac x=regS in exI)
apply (rule-tac x=l'' in exI)
apply (rule-tac x=m' in exI)
apply (rule-tac x=m'' in exI)
apply (rule-tac x=d in exI)
apply (rule-tac x=g in exI)

apply (rule conjI)
apply simp
apply (case-tac S!+off)
  apply simp
  apply (case-tac Vala)
  apply simp
  apply (insert RightNotUndefined)
  apply (erule-tac x= S2 in allE, force)
  apply clarsimp
apply clarsimp
apply clarsimp
apply (insert RightNotUndefined)
apply (erule-tac x= S2 in allE, force)
apply clarsimp
apply (insert RightNotUndefined)
apply (erule-tac x= S2 in allE, force)
apply (rule conjI, rule refl)
apply (rule conjI, rule refl)
apply (rule conjI, rule refl)
apply (rule conjI, assumption)
apply (rule conjI, assumption)
apply (rule conjI, simp)

```

```

apply (rule conjI, assumption)
apply (rule conjI, assumption)
apply (rule conjI, assumption)
apply (rule conjI, assumption)
apply (rule conjI, assumption)
apply (rule conjI, assumption)
apply (rule conjI, simp)
apply (rule conjI, simp)
apply (rule conjI, simp)
apply (rule conjI, simp)
apply (rule conjI, simp)
apply (rule conjI, clarsimp)
by (rule pc-MATCHN)

```

**end**

```

theory dem-BUILDENV
imports ../JVMSAFE/JVMExec SVM2JVM SVMSemantics CertifSVM2JVM
          dem-translation

```

**begin**

```

no-translations Norm s == (None,s)

```

```

no-translations ex-table-of m == snd (snd (snd m))

```

```

declare trInstr.simps [simp del]
declare equivS.simps [simp del]
declare extractBytecode-def [simp del]
declare initConsTable-def [simp del]

```

```

declare exec-instr.simps [simp]

```

```

fun Item2val :: entries-  $\Rightarrow$  sheap  $\Rightarrow$  Item  $\Rightarrow$  val

```

**where**

```

  Item2val S sh (ItemConst v) = (if (isBool v = True)
    then Bool (the-BoolT v)
    else Intg (the-IntT v))
| Item2val S sh (ItemVar l)   = the (S (nat (the-Intg (the (sh (stackC, topf)))
  - int l)))
| Item2val S sh (ItemRegSelf) = the (sh (heapC, kf))

```

**axioms** *maxPush-BUILDENV*:

$\text{fst } (\text{the } (\text{map-of svms } l)) ! i = \text{BUILDENV items} \implies n + \text{int } (\text{length items}) < \text{int } m$

**axioms** *pushAux-instrs*:

$P \vdash$   
 (None, sha, h, inih,  
 [([],  
 vs[Suc 0 := Intg (n + int (length items))],  
 safeP, sigSafeMain, the (cdm (l, i)) + 1 + 1 + 1 + 1, ag, bd])  
 $\text{-jvm} \rightarrow$   
 (None,  
 sha,  
 (let vitems = map (Item2val S' sha) items;  
 idxs = decreasing (nat (the-Intg (the (sha (stackC, topf))) + int (length items)));  
 S'' = S' ++ map-of (zip idxs vitems)  
 in h(the-Addr (the (sha (stackC, Sf)))  $\mapsto$  Arr ty m S''),  
 inih,  
 [([],  
 vs[Suc 0 := Intg (n + int (length items))],  
 safeP, sigSafeMain, the (cdm (l, i)) + 1 + 1 + 1 + 1 +  
 length (concat (map pushAux (zip items [0.. $\text{length items}$ ]))), ag, bd])

**axioms** *equivS-BUILDENV*:

$\text{equivS } S S' n \text{ ctm } d g$   
 $\implies \text{equivS } (\text{map } (\text{item2Stack } k S) \text{ items } @ S)$   
 (S' ++ map-of (zip (decreasing (nat (n + int (length items)))) (map  
 (Item2val S' sha) items))  
 (n + int (length items)) ctm d g

**declare** *equivH.simps* [simp del]

**lemma** *activeCells-BUILDENV*:

$\llbracket \text{the-Addr } (\text{the } (\text{sha } (\text{stackC}, \text{Sf}))) \notin \text{activeCells } \text{regS } h \text{ (nat } (\text{the-Intg } (\text{the } (\text{sha } (\text{heapC}, \text{kf})))))) \rrbracket$   
 $\implies \text{activeCells } \text{regS } h \text{ (nat } (\text{the-Intg } (\text{the } (\text{sha } (\text{heapC}, \text{kf})))))) =$   
 $\text{activeCells } \text{regS}$

```

      (h(the-Addr (the (sha (stackC, Sf)))  $\mapsto$ 
      Arr ty m
      (S' ++ map-of (zip (decreasing (nat (the-Intg (the (sha (stackC, topf)))
+ int (length items))))
      (map (Item2val S' sha) items)))))) (nat (the-Intg (the (sha
(heapC, kf))))))
apply (frule l-not-in-cellReg)
apply (unfold activeCells-def,auto)
apply (unfold region-def, simp add: Let-def, elim conjE)
apply (rule-tac x=j in exI,simp)
apply (rule cellReg-step)
apply (rule cellReg-basic)
apply (erule-tac x=j in allE)
apply (rule cellsReg-monotone-1,assumption+)
apply (erule l-not-in-regs)
apply (simp add: Let-def, elim conjE)
apply (rule-tac x=j in exI, simp)
apply (rule cellReg-step)
apply (rule cellReg-basic)
apply (erule-tac x=j in allE)
apply (erule cellsReg-monotone-2,assumption+)
by (erule l-not-in-regs)

```

**lemma** domH-BUILDENV:

```

[[ la  $\notin$  activeCells regS h (nat (the-Intg (the (sha (heapC, kf)))));
   $\forall l \in \text{dom } H. \exists l'. l' = \text{the } (g \ l) \wedge \text{equivC } (\text{the } (H \ l)) \ h \ l' \ (\text{the } (h \ l')) \ (\text{nat}$ 
  (the-Intg (the (sha (heapC, kf)))))) com regS d g ]]
 $\implies \forall l \in \text{dom } H.$ 
   $\exists l'. l' = \text{the } (g \ l) \wedge$ 
  equivC (the (H l))
  (h(la  $\mapsto$ 
  Arr ty m
  (S' ++ map-of (zip (decreasing (nat (the-Intg (the (sha (stackC,
topf))) + int (length items))))
  (map (Item2val S' sha) items))))))
  l' (the ((h(la  $\mapsto$ 
  Arr ty m
  (S' ++ map-of (zip (decreasing (nat (the-Intg (the (sha
(stackC, topf))) + int (length items))))
  (map (Item2val S' sha) items))))))
  l'))
  (nat (the-Intg (the (sha (heapC, kf))))))
  com regS d g
apply (frule l-not-in-cellReg)
apply (rule ballI)
apply (erule-tac x=l in ballE)

```

```

prefer 2 apply simp
apply (erule exE)
apply (rule-tac x=l' in exI) apply (elim conjE)
  apply (rule conjI, assumption)
apply clarsimp
apply (rule conjI) apply (rule impI)
apply clarsimp apply (simp add: region-def)
apply clarsimp
apply (rule-tac x=Obj in exI)
apply (rule-tac x=flds in exI) apply simp
apply (simp add: region-def) apply (elim conjE)
apply (rule cellReg-step)
  apply (rule cellReg-basic)
  apply (erule-tac x=j in allE)
  apply (rule cellsReg-monotone-1, assumption+)
by (erule l-not-in-regs)

```

**lemma** *equivH-BUILDENV*:

```

[[ k' = nat (the-Intg (the (sha (heapC, kf))));
  (hm, k), k0, (l, i), S) = ((H, ka), k0a, PC, Sa);
  equivH (H, ka) h k' com regS d g;
  sha (stackC, Sf) = Some (Addr la);
  activeCells regS h k' ∩ {la, l', l''} = {}]]
⇒ equivH (hm, k)
  (h(the-Addr (the (sha (stackC, Sf)))) ↦
    Arr ty m
    (S' ++ map-of (zip (decreasing (nat (the-Intg (the (sha (stackC, topf))))
      + int (length items))))
      (map (Item2val S' sha) items))))
  (nat (the-Intg (the ((sha((stackC, topf) ↦ vs[Suc 0 := Intg (n + int
    (length items))) ! Suc 0)) (heapC, kf)))) com regS d g
  apply (simp add: heapC-def add: stackC-def)
apply (simp add: kf-def add: k0f-def)
apply (fold heapC-def, fold kf-def)
apply (unfold equivH.simps, elim conjE)
apply (rule conjI, assumption)
apply (fold stackC-def)
apply (rule conjI)
apply (subgoal-tac
  the-Addr (the (sha (stackC, Sf))) ∉ activeCells regS h (nat (the-Intg (the (sha
    (heapC, kf))))))
  apply (frule activeCells-BUILDENV, simp)
apply simp
apply (rule conjI, assumption)
by (rule domH-BUILDENV, simp, assumption)

```

**lemma** *execSVMInstr-BUILDENV*:



$\llbracket (P', cdm, ctm, com) = trSVM2JVM ((svms, ctm), ini, ct, ah, ai, bc);$   
 $cdm, ctm, com \vdash ((hm, k), k0, (l, i), S) \triangleq S1';$   
 $(fst (the (map-of svms l)) ! i) = BUILDENV items;$   
 $execSVMInst (BUILDENV items) (map-of ct) (hm, k) k0 (l, i) S = E-$   
*ther.Right S2;*  
 $drop (the (cdm (l, i))) (extractBytecode P') =$   
 $trInstr (the (cdm (l, i))) cdm' ctm' com pcc (BUILDENV items) @ bytecode'$   
 $\rrbracket \implies \exists v' sh' dh' ih' fms' .$   
 $P' \vdash S1' -jvm \rightarrow (v', sh', dh', ih', fms') \wedge$   
 $cdm, ctm, com \vdash S2 \triangleq (v', sh', dh', ih', fms')$

**apply** (case-tac S1')  
**apply** (rename-tac v tup)  
**apply** (case-tac tup)  
**apply** (rename-tac sh tup)  
**apply** (case-tac tup)  
**apply** (rename-tac dh tup)  
**apply** (case-tac tup)  
**apply** (rename-tac ih fms)  
**apply** (simp)

**apply** (unfold equivState-def)  
**apply** (elim exE, elim conjE)

**apply** (rule-tac x=None in exI)  
**apply** (rule-tac x=sha((stackC, topf)  $\mapsto$  vs[Suc 0 := Intg (n + int (length items))]  
! Suc 0) in exI)  
**apply** (rule-tac x=(let vitems = map (Item2val S' sha) items;  
idxs = decreasing (nat (the-Intg (the (sha (stackC, topf))) +  
int (length items)));  
S'' = S' ++ map-of (zip idxs vitems)  
in h(the-Addr (the (sha (stackC, Sf)))  $\mapsto$  Arr ty m S'')) in exI)  
**apply** (rule-tac x=inih in exI)  
**apply** (rule-tac x=[[],  
vs[Suc 0 := Intg (n + int (length items))],  
safeP, sigSafeMain,  
the (cdm (l, i)) + 1 + 1 + 1 + 1 + 1 + 1 +  
length (concat (map pushAux (zip items [0..<length items])))],  
ref]) in exI)

**apply** (subgoal-tac  $\exists$  vs pc . fms=([], vs, safeP, sigSafeMain, pc, ref)#[])  
**prefer** 2 **apply** simp  
**apply** (erule exE)+

```

apply (rule conjI)

apply (unfold exec-all-def)
apply (subgoal-tac PC = (l,i)) prefer 2 apply simp
apply simp apply (elim conjE) apply clarsimp
apply (subgoal-tac P⊢ (None,sha,h,inih,[[[], vs, safeP, sigSafeMain, the (cdm
(l,i)), ag,bd]])
      -jvm→
      (None,sha,h,inih,[[the (sha (stackC,topf))], vs, safeP,
sigSafeMain, the (cdm (l,i)) + 1, ag,bd]])
apply (unfold exec-all-def)
apply (erule rtrancl-trans)

prefer 2
apply (rule r-into-rtrancl)
apply (clarify)
apply (unfold JVMExec.exec.simps)
apply (unfold Let-def)
apply (unfold extractBytecode-def)
apply (unfold trInstr.simps)
apply (unfold Let-def)
apply (subgoal-tac (fst (snd (snd (snd (snd (the (method' (P', safeP) sigSafe-
Main))))))) !
      the (cdm (l, i))) = Getstatic topf stackC,simp)
apply (erule nth-via-drop-append-2)

apply (erule drop-Suc-append-2)
apply (subgoal-tac P⊢ (None, sha, h, inih,
      [[the (sha (stackC,topf))],
      vs, safeP, sigSafeMain, the (cdm (l, i)) + 1, ag, bd]])
      -jvm→
      (None,sha,h,inih,
      [[Intg (int (length items)),
      the (sha (stackC,topf))],
      vs, safeP, sigSafeMain, the (cdm (l, i)) + 1 + 1, ag, bd]])
apply (unfold exec-all-def)
apply (erule rtrancl-trans)
prefer 2
apply (rule r-into-rtrancl)
apply (clarify)
apply (unfold JVMExec.exec.simps)
apply (unfold Let-def)
apply (subgoal-tac (fst (snd (snd (snd (snd (the (method' (P', safeP) sigSafe-
Main))))))) ! Suc (the (cdm (l, i))) =
      LitPush (Intg (int (length items))),simp)
apply (erule nth-via-drop-append)

```

```

apply (drule drop-Suc-append)
apply (subgoal-tac  $P \vdash$  (None,sha,h,inih,
  [[Intg (int (length items)),
   the (sha (stackC,topf)]]),
  vs, safeP, sigSafeMain, the (cdm (l, i) + 1 + 1, ag, bd)]
   $\rightarrow$ 
  (None,sha,h,inih,
  [[Intg (n + int (length items)]]),
  vs, safeP, sigSafeMain, the (cdm (l, i) + 1 + 1 + 1, ag, bd)))
apply (unfold exec-all-def)
apply (erule rtrancl-trans)
prefer 2
apply (rule r-into-rtrancl)
apply (clarify)
apply (unfold JVMExec.exec.simps)
apply (unfold Let-def)
apply (subgoal-tac (fst (snd (snd (snd (snd (the (method' (P', safeP) sigSafeMain)))))) ! (Suc (Suc (the (cdm (l, i)))))) =
  BinOp Add,simp)
apply (erule nth-via-drop-append)

```

```

apply (drule drop-Suc-append)
apply (subgoal-tac  $P \vdash$  (None,sha,h,inih,
  [[Intg (n + int (length items)]]),
  vs, safeP, sigSafeMain, the (cdm (l, i) + 1 + 1 + 1, ag, bd)]
   $\rightarrow$ 
  (None,sha,h,inih,
  [[ $\square$ ],
  vs[Suc 0 := Intg (n + int (length items)]]),
  safeP, sigSafeMain, the (cdm (l, i) + 1 + 1 + 1 + 1, ag, bd]))
apply (unfold exec-all-def)
apply (erule rtrancl-trans)
prefer 2
apply (rule r-into-rtrancl)
apply (clarify)
apply (unfold JVMExec.exec.simps)
apply (unfold Let-def)
apply (subgoal-tac (fst (snd (snd (snd (snd (the (method' (P', safeP) sigSafeMain)))))) ! (Suc (Suc (Suc (the (cdm (l, i)))))) =
  Store 1,simp)
apply (erule nth-via-drop-append)

```

```

apply (drule drop-Suc-append)
apply (subgoal-tac  $P \vdash$  (None,sha,h,inih,

```

```

    [([],
      vs[Suc 0 := Intg (n + int (length items))],
      safeP, sigSafeMain, the (cdm (l, i)) + 1 + 1 + 1 + 1, ag,
      bd))]
  -jvm→
    (None, sha,
     (let vitems = map (Item2val S' sha) items;
        idxs = decreasing (nat (the-Intg (the (sha (stackC,
topf))) + int (length items))));
      S'' = S' ++ map-of (zip idxs vitems)
      in h(the-Addr (the (sha (stackC, Sf))) ↦ Arr ty m S'')),
     inih,
    [([],
      vs[Suc 0 := Intg (n + int (length items))],
      safeP, sigSafeMain, the (cdm (l, i)) + 1 + 1 + 1 + 1 +
        length (concat (map pushAux (zip items [0..<length
items]))), ag, bd))]
prefer 2 apply (rule pushAux-instrs)
apply (unfold exec-all-def)
apply (erule rtrancl-trans)

apply (drule drop-append-length)
apply (subgoal-tac P ⊢ (None, sha,
  (let vitems = map (Item2val S' sha) items;
     idxs = decreasing (nat (the-Intg (the (sha (stackC,
topf))) + int (length items))));
    S'' = S' ++ map-of (zip idxs vitems)
    in h(the-Addr (the (sha (stackC, Sf))) ↦ Arr ty m S'')),
     inih,
  [([],
    vs[Suc 0 := Intg (n + int (length items))],
    safeP, sigSafeMain, the (cdm (l, i)) + 1 + 1 + 1 + 1 +
      length (concat (map pushAux (zip items [0..<length
items]))), ag, bd))]
  -jvm→
    (None, sha,
     (let vitems = map (Item2val S' sha) items;
        idxs = decreasing (nat (the-Intg (the (sha (stackC,
topf))) + int (length items))));
      S'' = S' ++ map-of (zip idxs vitems)
      in h(the-Addr (the (sha (stackC, Sf))) ↦ Arr ty m S'')),
     inih,
    [([vs[Suc 0 := Intg (n + int (length items))] ! Suc 0],
      vs[Suc 0 := Intg (n + int (length items))],
      safeP, sigSafeMain, the (cdm (l, i)) + 1 + 1 + 1 + 1 + 1 +
        length (concat (map pushAux (zip items [0..<length
items]))), ag, bd))]
apply (unfold exec-all-def)

```

```

apply (erule rtrancl-trans)
prefer 2
apply (rule r-into-rtrancl)
apply (clarify)
apply (unfold JVMExec.exec.simps)
apply (unfold Let-def)
apply (subgoal-tac (fst (snd (snd (snd (snd (the (method' (P', safeP) sigSafe-
Main))))))) !
      (((Suc (Suc (Suc (Suc (the (cdm (l, i))))))) +
length (concat (map pushAux (zip items [0..apply (erule nth-via-drop-append)

apply (drule drop-Suc-append)
apply (subgoal-tac P+ (None,sha,
      (let vitems = map (Item2val S' sha) items;
          idxs = decreasing (nat (the-Intg (the (sha (stackC,
topf)))) + int (length items)));
          S'' = S' ++ map-of (zip idxs vitems)
          in h(the-Addr (the (sha (stackC, Sf)))  $\mapsto$  Arr ty m S''),
          inih,
          [[vs[Suc 0 := Intg (n + int (length items))] ! Suc 0],
          vs[Suc 0 := Intg (n + int (length items))],
          safeP, sigSafeMain, the (cdm (l, i)) + 1 + 1 + 1 + 1 +
length (concat (map pushAux (zip items [0..\rightarrow
      (None,
        sha((stackC, topf)  $\mapsto$  vs[Suc 0 := Intg (n + int (length
items))] ! Suc 0),
      (let vitems = map (Item2val S' sha) items;
          idxs = decreasing (nat (the-Intg (the (sha (stackC,
topf)))) + int (length items)));
          S'' = S' ++ map-of (zip idxs vitems)
          in h(the-Addr (the (sha (stackC, Sf)))  $\mapsto$  Arr ty m S''),
          inih,
          [([],
          vs[Suc 0 := Intg (n + int (length items))],
          safeP, sigSafeMain, the (cdm (l, i)) + 1 + 1 + 1 + 1 + 1
+ 1 +
length (concat (map pushAux (zip items [0..apply (unfold exec-all-def)
prefer 2
apply (rule r-into-rtrancl)
apply (clarify)
apply (unfold JVMExec.exec.simps)

```

```

apply (unfold Let-def)
apply (subgoal-tac (fst (snd (snd (snd (snd (the (method' (P', safeP) sigSafe-
Main))))))) !
      (((Suc (Suc (Suc (Suc (Suc (the (cdm (l, i)))))))))) +
      length (concat (map pushAux (zip items [0..<length items]))) =
      Putstatic topf stackC, simp)
apply (rule nth-via-drop-append, force)
apply simp

apply (frule nonJumping-Suc-pc)
apply (erule-tac sym [where t=BUILDENV items])
apply (simp add: nonJumping.simps)
apply (simp add: trInstr.simps)

apply (rule-tac x=hm in exI)
apply (rule-tac x=k in exI)
apply (rule-tac x=k0 in exI)
apply (rule-tac x=(l, Suc i) in exI)
apply (rule-tac x=map (item2Stack k S) items @ S in exI)

apply (rule-tac x=sha((stackC, topf)  $\mapsto$  vs[Suc 0 := Intg (n + int (length items))]
! Suc 0) in exI)
apply (rule-tac x=h(the-Addr (the (sha (stackC, Sf)))  $\mapsto$  Arr ty m
      (S' ++ map-of (zip (decreasing (nat (the-Intg (the (sha (stackC,
topf)))) + int (length items))))
      (map (Item2val S' sha) items))) in exI)
apply (rule-tac x=inih in exI)
apply (rule-tac x=vs[Suc 0 := Intg (n + int (length items))] in exI)
apply (rule-tac x=the (cdm (l, i)) + 1 + 1 + 1 + 1 + 1 + 1 + length (concat
(map pushAux (zip items [0..<length items]))) in exI)
apply (rule-tac x=ref in exI)

apply (rule-tac x=nat (the-Intg (the ((sha((stackC, topf)  $\mapsto$ 
      vs[Suc 0 := Intg (n + int (length items))] ! Suc 0)) (heapC,
kf)))) in exI)

apply (rule-tac x= nat (the-Intg (the ((sha((stackC, topf)  $\mapsto$ 
      vs[Suc 0 := Intg (n + int (length items))] ! Suc 0)) (heapC,
k0f)))) in exI)

apply (rule-tac x=la in exI)
apply (rule-tac x=ty in exI)
apply (rule-tac x=m in exI)
apply (rule-tac x=S' ++ map-of (zip (decreasing (nat (n + int (length items))))
(map (Item2val S' sha) items)) in exI)

apply (rule-tac x=(n + int (length items)) in exI)
apply (rule-tac x=l' in exI)

```

```

apply (rule-tac x=regS in exI)
apply (rule-tac x=l'' in exI)
apply (rule-tac x=m' in exI)
apply (rule-tac x=m'' in exI)
apply (rule-tac x=d in exI)
apply (rule-tac x=g in exI)

apply (rule conjI, erule sym)
apply (rule conjI, rule refl)
apply (rule conjI, rule refl)
apply (rule conjI, rule refl)
apply (rule conjI) apply clarsimp
apply (rule conjI, assumption)
apply (rule conjI) apply clarsimp
      apply (simp add: heapC-def add: stackC-def)
apply (rule conjI) apply (erule activeCells)
apply (rule conjI) apply (erule activeCells-2,assumption)
      apply (erule activeCells-2,assumption)
apply (rule conjI) apply clarsimp
apply (rule conjI) apply clarsimp
      apply (subgoal-tac length vs = 10)
      apply clarsimp
      apply (rule length-vs)
apply (rule conjI) apply clarsimp
apply (rule conjI) apply clarsimp
apply (rule conjI) apply clarsimp
apply (rule conjI) apply clarsimp
apply (rule conjI) apply clarsimp
apply (rule conjI) apply (rule equivH-BUILDENV,assumption+)
apply (rule conjI) apply clarsimp
      apply (rule maxPush-BUILDENV,assumption+)
apply (rule conjI) apply clarsimp
      apply (erule equivS-BUILDENV)
apply (rule conjI) apply (subgoal-tac length vs = 10)
      apply (simp add: heapC-def add: stackC-def)
      apply (rule length-vs)

by simp

end

theory dem-CALL
imports ../JVMSAFE/JVMExec ../JVMSAFE/State SVM2JVM SVMSemantics
CertifSVM2JVM
      dem-translation

begin

```

```

no-translations Norm s == (None,s)
no-translations ex-table-of m == snd (snd (snd m))

declare trInstr.simps [simp del]
declare equivS.simps [simp del]
declare extractBytecode-def [simp del]
declare initConsTable-def [simp del]

axioms pushregion-method:
  [| s = (None, shp, hp, ihp,
    ([],vs, safeP, sigSafeMain,pc,z1,z2)#[]);
   shp (stackC,topf) = Some (Intg k);
   shp (stackC,Sf) = Some (Addr l);
   shp (heapC,regionsf) = Some (Addr l'); hp l' = Some (Arr ty m regS);
   shp (dirCellC,safeDirf) = Some (Addr l'');
   Invoke-static heapC pushRegion [] = fst(snd(snd(snd(snd(the(method' (P',safeP)
sigSafeMain)))))) ! pc
  |] ==>
  P' |- s -jvm-> (None, shp ((heapC,kf)↦ Intg (k+1)),
    hp(l'↦ Arr ty ma (regS (nat k + 1 ↦ Addr lnew)))
    (lnew ↦ Obj cellC (empty ((izqf,cellC)↦ Addr lnew)
    ((derf,cellC)↦ Addr lnew))),
    ihp, ([],loc, safeP, sigSafeMain,pc+1,z1,z2)#[]) ^
  lnew ∉ {l,l',l''} ∧
  lnew ∉ activeCells regS h (nat (the-Intg (the (sha (heapC, kf))))))

declare equivH.simps [simp del]

declare equivH.simps [simp del]

lemma cellsReg-monotone-3 [rule-format]:
  x ∈ cellsReg (h(l'↦ A, l''↦ B)) l
  → x ≠ l
  → l ≠ l'
  → l ≠ l''
  → x = nextCell h l
apply (rule impI)
apply (erule cellsReg.induct,simp)
by (simp add: nextCell-def)

lemma cellsReg-monotone-4 [rule-format]:
  x ∈ cellsReg h l

```



$\longrightarrow x \neq l$   
 $\longrightarrow l \neq l'$   
 $\longrightarrow l \neq l''$   
 $\longrightarrow x = \text{nextCell } (h(l' \mapsto A, l'' \mapsto B)) \ l$   
**apply** (rule impI)  
**apply** (erule cellsReg.induct,simp)  
**by** (simp add: nextCell-def)

**lemma** activeCells-CALL:

$\llbracket l' \notin \text{activeCells regS dh } (\text{nat } (\text{the-Intg } (\text{the } (\text{sha } (\text{heapC}, \text{kf})))));$   
 $\quad \text{lnew} \notin \text{activeCells regS dh } (\text{nat } (\text{the-Intg } (\text{the } (\text{sha } (\text{heapC}, \text{kf}))))))$   
 $\rrbracket$   
 $\implies \text{activeCells regS dh } (\text{nat } (\text{the-Intg } (\text{the } (\text{sha } (\text{heapC}, \text{kf})))))) =$   
 $\quad \text{activeCells } (\text{regS } (\text{Suc } (\text{nat } (\text{the-Intg } (\text{the } (\text{sha } (\text{heapC}, \text{kf})))))) \mapsto \text{Addr lnew}))$   
 $\quad (\text{dh}(l' \mapsto \text{Arr ty } m' (\text{regS}(\text{Suc } (\text{nat } (\text{the-Intg } (\text{the } (\text{sha } (\text{heapC}, \text{kf})))))) \mapsto$   
 $\text{Addr lnew})), \text{lnew} \mapsto$   
 $\quad \text{Obj cellC } [(izqf, \text{cellC}) \mapsto \text{Addr lnew}, (\text{derf}, \text{cellC}) \mapsto \text{Addr lnew}]))$   
 $\quad (\text{nat } (\text{the-Intg } (\text{the } (\text{sha } (\text{heapC}, \text{kf})))) + 1))$   
**apply** (frule l-not-in-cellReg)  
**apply** (frule-tac l=lnew in l-not-in-cellReg)  
**apply** (rule equalityI)  
**apply** (rule subsetI)  
**apply** (simp add: activeCells-def,clarsimp)  
**apply** (unfold region-def, simp add: Let-def, elim conjE)  
**apply** (rule-tac x=j in exI,simp)  
**apply** (rule conjI, simp)  
**apply** (rule cellReg-step)  
**apply** (rule cellReg-basic)  
**apply** (erule-tac x=j in allE)  
**apply** (rule cellsReg-monotone-1)  
**apply** (rule cellReg-step)  
**apply** (rule cellReg-basic)  
**apply** (rule cellsReg-monotone-1,assumption+)  
**apply** (rule l-not-in-regs)  
**apply** (erule-tac x=j in allE,simp,simp)  
**apply** (rule l-not-in-regs)  
**apply** (erule-tac x=j in allE)+  
**apply** simp  
**apply** (rule subsetI)  
**apply** (unfold activeCells-def)  
**apply** (unfold region-def, simp add: Let-def)  
**apply** (erule exE)  
**apply** (case-tac  
 $j < \text{Suc } (\text{nat } (\text{the-Intg } (\text{the } (\text{sha } (\text{heapC}, \text{kf}))))))$ ,simp)  
**apply** (rule-tac x=j in exI,simp)  
**apply** (elim conjE)  
**apply** (rule cellReg-step)

```

apply (rule cellReg-basic)
apply (rule cellsReg-monotone-3,simp,simp)
  apply (rule l-not-in-regs)
  apply (erule-tac x=j in allE)+ apply simp
  apply (rule l-not-in-regs)
  apply (erule-tac x=j in allE)+ apply simp
apply (elim conjE)
apply (subgoal-tac
  j = nat (the-Intg (the (sha (heapC, kf))) + 1),simp)
prefer 2 apply simp
apply (subgoal-tac nat (the-Intg (the (sha (heapC, kf))) + 1) = Suc (nat (the-Intg
(the (sha (heapC, kf))))))
prefer 2 apply arith
apply simp
apply (subgoal-tac x=lnew,simp)
apply (erule cellsReg.induct)
apply simp
by (simp add: nextCell-def)

```

**lemma** domH-CALL:

```

[[ l' ∉ activeCells regS h (nat (the-Intg (the (sha (heapC, kf)))));
  lnew ∉ activeCells regS h (nat (the-Intg (the (sha (heapC, kf)))));
  lnew ∉ {la,l',l''};
  ∀ l ∈ dom H. ∃ l'. l' = the (g l) ∧ equivC (the (H l)) h l' (the (h l')) (nat
(the-Intg (the (sha (heapC, kf)))))) com regS d g ]]
⇒ ∀ l ∈ dom H.
  ∃ l'a. l'a = the (g l) ∧
    equivC (the (H l))
      (h(l' ↦ Arr ty m' (regS(Suc (nat (the-Intg (the (sha (heapC, kf)))))))
↦ Addr lnew)), lnew ↦
      Obj cellC [(izqf, cellC) ↦ Addr lnew, (derf, cellC) ↦ Addr lnew]))
    l'a (the ((h(l' ↦ Arr ty m' (regS(Suc (nat (the-Intg (the (sha (heapC,
kf)))))) ↦ Addr lnew)), lnew ↦
      Obj cellC [(izqf, cellC) ↦ Addr lnew, (derf, cellC) ↦ Addr
lnew]))
      l'a))
  (nat (the-Intg (the (sha (heapC, kf))) + 1))
    com (regS(Suc (nat (the-Intg (the (sha (heapC, kf)))))) ↦ Addr
lnew)) d g
apply (frule l-not-in-cellReg)
apply (frule-tac l=lnew in l-not-in-cellReg)
apply (rule ballI)
apply (erule-tac x=l in ballE)
prefer 2 apply simp

```

```

apply (erule exE)
apply (rule-tac x=l'a in exI) apply (elim conjE)
  apply (rule conjI, assumption)
apply clarsimp
apply (rule conjI, rule impI)
apply (simp add: activeCells-def)
apply (simp add: activeCells-def)
apply (erule-tac x=j in allE) apply simp
apply (subgoal-tac the (g l) ≠ lnew,simp)
  prefer 2 apply blast
apply (rule impI)
apply (rule conjI, simp)
apply (rule-tac x=Obja in exI)
apply (rule-tac x=flds in exI) apply simp
apply (simp add: region-def)
apply clarsimp
apply (rule cellReg-step)
  apply (rule cellReg-basic)
  apply (rule cellsReg-monotone-4) apply assumption+
  apply (rule l-not-in-regs)
  apply (erule-tac x=j in allE) apply simp
apply (rule l-not-in-regs)
apply (erule-tac x=j in allE)
apply (erule-tac x=j in allE)
apply (erule-tac x=j in allE)
by simp

```

**axioms** kf-ge-0:  
 $the-Intg (the (sha (heapC, kf))) \geq 0$

**lemma** equivH-CALL:

```

[[((hm, k), k0, (l, i), S) = ((H, ka), k0a, PC, Sa);
 (v, sh, dh, ih, fms) = (None, sha, h, inih, [(), vs, safeP, sigSafeMain, pc,
 ref])]];
k' = nat (the-Intg (the (sha (heapC, kf))));
k0' = nat (the-Intg (the (sha (heapC, k0f))));
sha (stackC, Sf) = Some (Addr la);
lnew ∉ activeCells regS h (nat (the-Intg (the (sha (heapC, kf))));
distinct [la, l', l']; activeCells regS h k' ∩ {la, l', l'} = {}; h la = Some (Arr
ty m S');
sha (stackC, topf) = Some (Intg n); sha (heapC, regionsf) = Some (Addr l');
h l' = Some (Arr ty m' regS);
lnew ∉ {la, l', l'};
inj-on g (dom H); equivH (H, ka) h k' com regS d g]]
⇒ equivH (hm, Suc k)
(dh(l' ↦ Arr ty m' (regS(k' + 1 ↦ Addr lnew)), lnew ↦ Obj cellC [(izqf,
cellC) ↦ Addr lnew, (derf, cellC) ↦ Addr lnew]))

```

```

      (nat (the-Intg (the ((sh((heapC, kf) ↦ Intg (the-Intg (the (sh (heapC,
kf))) + 1))) (heapC, kf)))) com
      (regS(Suc (nat (the-Intg (the (sha (heapC, kf)))))) ↦ Addr lnew)) d g
apply (simp add: heapC-def add: stackC-def)
apply (fold heapC-def)
apply (unfold equivH.simps, elim conjE)
apply (rule conjI, clarsimp)
apply (subgoal-tac the-Intg (the (sha (heapC, kf))) >= 0, simp)
  apply (rule kf-ge-0)
apply (rule conjI, frule-tac l'=l' in activeCells-CALL, assumption+, simp)
apply (rule conjI, assumption)
apply (subgoal-tac lnew ∉ {la, l', l''})
apply (rule domH-CALL, assumption+)
by clarsimp

```

**axioms** *activeCells-3*:

$$\llbracket l' \notin \text{activeCells } \text{regS } h \ k; l \neq l'; l \neq l'' \rrbracket \\ \implies l' \notin \text{activeCells } \text{regS } (h(l \mapsto A, l'' \mapsto B)) \ k$$

**axioms** *activeCells-3-1*:

$$\llbracket l' \notin \text{activeCells } \text{regS } h \ k \rrbracket \\ \implies l' \notin \text{activeCells } \text{regS } (h(l' \mapsto A, l'' \mapsto B)) \ k$$

**axioms** *activeCells-4*:

$$la \notin \text{activeCells } \text{regS } h \ k \\ \implies la \notin \text{activeCells } \text{regS}' \ h \ k'$$

**axioms** *kf-Intg*:

$$\text{sha } (heapC, kf) = \text{Some } (\text{Intg } j)$$

**axioms** *p-in-dom-cdm'*:

$$(p, 0) \in \text{dom } \text{cdm}'$$

**axioms** *lnew-notin-activeCells*:

$$lnew \notin \text{activeCells } \text{regS } h \ (\text{nat } (\text{the-Intg } (\text{the } (\text{sha } (heapC, kf))))))$$

**lemma** *execSVMInstr-CALL* :

$$\llbracket (P', \text{cdm}, \text{ctm}, \text{com}) = \text{trSVM2JVM } ((\text{svms}, \text{ctmap}), \text{ini}, \text{ct}, \text{ah}, \text{ai}, \text{bc}); \\ \text{cdm}, \text{ctm}, \text{com} \vdash ((\text{hm}, k), k0, (l, i), S) \hat{=} S1'; \\ (\text{fst } (\text{the } (\text{map-of } \text{svms } l)) ! i) = \text{CALL } p; \\ \text{execSVMInst } (\text{CALL } p) (\text{map-of } \text{ct}) (\text{hm}, k) k0 (l, i) S = \text{Either.Right } S2; \\ \text{drop } (\text{the } (\text{cdm } (l, i))) (\text{extractBytecode } P') = \\ \text{trInstr } (\text{the } (\text{cdm } (l, i))) \text{cdm}' \text{ctm}' \text{com } \text{pcc } (\text{CALL } p) @ \text{bytecode}'; \\ \text{cdm}' \subseteq_m \text{cdm} \rrbracket$$

$$\begin{aligned} \llbracket &\implies \exists v' sh' dh' ih' fms' . \\ P' \vdash S1' -jvm \rightarrow &(v', sh', dh', ih', fms') \wedge \\ cdm, ctm, com \vdash S2 &\triangleq (v', sh', dh', ih', fms') \end{aligned}$$

```

apply (case-tac S1')
apply (rename-tac v tup)
apply (case-tac tup)
apply (rename-tac sh tup)
apply (case-tac tup)
apply (rename-tac dh tup)
apply (case-tac tup)
apply (rename-tac ih fms)
apply (simp)

```

```

apply (unfold equivState-def)
apply (elim exE, elim conjE)

```

```

apply (rule-tac x=None in exI)
apply (rule-tac x=sh((heapC, kf)  $\mapsto$  Intg (the-Intg (the (sh (heapC, kf))+1))
in exI)
apply (rule-tac x=dh(l'  $\mapsto$  Arr ty m' (regS (k'+1  $\mapsto$  Addr lnew)))
(lnew  $\mapsto$  Obj cellC (empty ((izqf, cellC)  $\mapsto$  Addr lnew)
((derf, cellC)  $\mapsto$  Addr lnew))) in exI)
apply (rule-tac x=ih in exI)
apply (rule-tac x=[[], vs, safeP, sigSafeMain, the (cdm (p, 0)), ref]) in exI)
apply (rule conjI)
apply (unfold extractBytecode-def)

```

```

apply (subgoal-tac P^ (v, sh, dh, ih, fms) -jvm  $\rightarrow$ 
(None, sh ((heapC, kf)  $\mapsto$  Intg (the-Intg (the (sh (heapC, kf))+1)),
dh (l'  $\mapsto$  Arr ty m' (regS (nat (the-Intg (the (sh (heapC, kf))))+1  $\mapsto$ 
Addr lnew)))
(lnew  $\mapsto$  Obj cellC (empty ((izqf, cellC)  $\mapsto$  Addr lnew)
((derf, cellC)  $\mapsto$  Addr lnew))),
ih, ([[], vs, safeP, sigSafeMain, pc+1, ref])#[])
apply (unfold exec-all-def )
apply (erule rtrancl-trans)

```

**prefer** 2

```

apply (subgoal-tac (fst (snd (snd (snd (snd (the (method' (P', safeP) sigSafe-
Main))))))) !
the (cdm (l, i))) = Invoke-static heapC pushRegion [])

```

```

apply (fold exec-all-def)
apply clarify
apply (subgoal-tac
P' |-
  (None, sha, h, inih, [[[], vs, safeP, sigSafeMain, the (cdm (l, i)), ag, bd]])
-jvm→
  (None, sha ((heapC, kf) ↦ Intg (n+1)),
    h(l' ↦ Arr ty m' (regS (nat n + 1 ↦ Addr lnew)))
    (lnew ↦ Obj cellC (empty ((izqf, cellC) ↦ Addr lnew)
      ((derf, cellC) ↦ Addr lnew))),
    inih, [[[], vs, safeP, sigSafeMain, the (cdm (l, i)) + 1, ag, bd]]) ∧
    lnew ∉ {la, l', l''} ∧
    lnew ∉ activeCells regS h (nat (the-Intg (the (sha (heapC, kf))))))
prefer 2 apply (rule pushregion-method, simp, assumption+, simp)
apply simp apply (subgoal-tac sha (heapC, kf) = Some (Intg n))
prefer 2 apply (rule kf-Intg)
apply simp

apply (unfold trInstr.simps)
apply (unfold Let-def)
apply (erule nth-via-drop-append)

apply (drule drop-Suc-append)
apply (unfold exec-all-def)
apply (rule r-into-rtrancl)
apply (clarify)
apply (unfold JVMEExec.exec.simps)
apply (unfold Let-def)
apply (subgoal-tac (fst (snd (snd (snd (snd (the (method' (P', safeP) sigSafe-
Main))))))) !
  Suc (the (cdm (l, i)))) = Goto (trAddr (the (cdm (p, 0))) (the (cdm (l, i)) +
incCall))) apply (simp)
apply (simp add: trAddr-def incCall-def)
apply (rule-tac ys=[] in nth-via-drop-append)
apply simp
apply (rule conjI)
apply (subgoal-tac cdm' (p, 0) = cdm (p, 0), simp)
apply (subgoal-tac (p, 0) ∈ dom cdm')
apply (simp add: map-le-def)
apply (rule p-in-dom-cdm')
apply simp

apply (rule-tac x=hm in exI)
apply (rule-tac x=Suc k in exI)
apply (rule-tac x=k0 in exI)

```

**apply** (*rule-tac*  $x=(p,0)$  **in** *exI*)  
**apply** (*rule-tac*  $x=S$  **in** *exI*)

**apply** (*rule-tac*  $x=sh((heapC, kf) \mapsto Intg (the-Intg (the (sh (heapC, kf))) + 1))$   
**in** *exI*)  
**apply** (*rule-tac*  $x=dh(l' \mapsto Arr\ ty\ m' (regS(k' + 1 \mapsto Addr\ lnew)),$   
 $lnew \mapsto Obj\ cellC [(izgf, cellC) \mapsto Addr\ lnew,$   
 $(derf, cellC) \mapsto Addr\ lnew])$  **in** *exI*)

**apply** (*rule-tac*  $x=ih$  **in** *exI*)  
**apply** (*rule-tac*  $x=vs$  **in** *exI*)  
**apply** (*rule-tac*  $x=the (cdm (p, 0))$  **in** *exI*)  
**apply** (*rule-tac*  $x=ref$  **in** *exI*)

**apply** (*rule-tac*  $x=nat (the-Intg (the ((sh((heapC, kf) \mapsto$   
 $Intg (the-Intg (the (sh (heapC, kf))) + 1))) (heapC, kf))))$  **in** *exI*)  
**apply** (*rule-tac*  $x=nat (the-Intg (the ((sh((heapC, kf) \mapsto$   
 $Intg (the-Intg (the (sh (heapC, kf))) + 1))) (heapC, kOf))))$  **in**  
*exI*)

**apply** (*rule-tac*  $x=la$  **in** *exI*)  
**apply** (*rule-tac*  $x=ty$  **in** *exI*)  
**apply** (*rule-tac*  $x=m$  **in** *exI*)  
**apply** (*rule-tac*  $x=S'$  **in** *exI*)

**apply** (*rule-tac*  $x=n$  **in** *exI*)  
**apply** (*rule-tac*  $x=l'$  **in** *exI*)  
**apply** (*rule-tac*  $x=regS(Suc (nat (the-Intg (the (sha (heapC, kf)))))) \mapsto Addr$   
 $lnew)$  **in** *exI*)  
**apply** (*rule-tac*  $x=l''$  **in** *exI*)  
**apply** (*rule-tac*  $x=m'$  **in** *exI*)  
**apply** (*rule-tac*  $x=m''$  **in** *exI*)  
**apply** (*rule-tac*  $x=d$  **in** *exI*)  
**apply** (*rule-tac*  $x=g$  **in** *exI*)

**apply** (*subgoal-tac*  
 $la \neq lnew \wedge l' \neq lnew \wedge l'' \neq lnew)$   
**prefer** 2  
**apply** (*case-tac* *ref,simp*)  
**apply** (*frule-tac*  $lnew=lnew$   
**in** *pushregion-method, assumption+*)  
**apply** (*clarsimp, rule sym*)  
**apply** (*erule nth-via-drop*)  
**apply** *clarsimp*

**apply** (*rule conjI, clarsimp*)  
**apply** (*rule conjI, rule refl*)  
**apply** (*rule conjI, rule refl*)  
**apply** (*rule conjI, rule refl*)

```

apply (rule conjI) apply (simp add: heapC-def add: stackC-def)
apply (rule conjI, assumption)
apply (rule conjI) apply clarsimp
      apply (simp add: heapC-def add: stackC-def)
apply (rule conjI) apply (rule activeCells-3)
      apply (rule activeCells-4)
      apply simp apply force
      apply (assumption+)
apply (rule conjI) apply (rule activeCells-3-1)
      apply (rule activeCells-4)
      apply simp
      apply (rule activeCells-3)
      apply (rule activeCells-4)
      apply simp apply force
      apply (assumption+)
apply (rule conjI) apply simp
apply (rule conjI) apply (simp add: heapC-def add: stackC-def)
apply (rule conjI) apply (simp add: regionsf-def add: kf-def)
apply (rule conjI) apply simp
apply (rule conjI) apply (simp add: safeDirf-def add: kf-def)
apply (rule conjI) apply clarsimp
apply (rule conjI) apply clarsimp
apply (rule conjI) apply simp
apply (rule conjI) apply (subgoal-tac
  lnew  $\notin$  activeCells regS h (nat (the-Intg (the (sha (heapC, kf))))))
  apply (elim conjE)
  apply (rule equivH-CALL) apply assumption+
  apply simp apply assumption+
  apply (rule lnew-notin-activeCells)
apply (rule conjI) apply simp
apply (rule conjI) apply simp
apply (rule conjI) apply simp
      apply (simp add: k0f-def add: kf-def)
by simp

end

```

```

theory dem-DECREGION
imports ../JVMSAFE/JVMEExec SVM2JVM SVMSemantics CertifSVM2JVM
      dem-translation
begin

no-translations Norm s == (None,s)
no-translations ex-table-of m == snd (snd (snd m))

declare trInstr.simps [simp del]
declare equivS.simps [simp del]
declare extractBytecode-def [simp del]

```



**declare** *initConsTable-def* [simp del]

**axioms** *decregion-method*:

$$\begin{aligned} & \llbracket s = (None, shp, hp, ihp, (\[], vs, safeP, sigSafeMain, pc, z) \# \[]) \rrbracket; \\ & shp (heapC, kf) = Some (Intg k); \\ & shp (heapC, k0f) = Some (Intg k0); \\ & shp (heapC, regionsf) = Some (Addr l); \\ & hp l = Some (Arr ty m regS); \\ & Invoke-static heapC decregion \[] = fst(snd(snd(snd(snd(the(method' (P', safeP) sigSafeMain)))))) ! pc \\ & \llbracket \implies \\ & P' \vdash s -jvm \rightarrow (None, shp ((heapC, kf) \mapsto Intg k0), \\ & \quad hp(l \mapsto Arr ty ma regS'), \\ & \quad ihp, (\[], loc, safeP, sigSafeMain, pc+1, z) \# \[]) \\ & \wedge (\forall j \leq nat k0 . region regS' hp j = region regS hp j) \end{aligned}$$

**axioms** *kf-Some-i*:

$$sha (heapC, kf) = Some (Intg jj)$$

**axioms** *k0f-Some-i*:

$$sha (heapC, k0f) = Some (Intg ii)$$

**axioms** *activeCells-3*:

$$\begin{aligned} & \llbracket l' \notin activeCells regS h k; l \neq l'; l \neq l'' \rrbracket \\ & \implies l' \notin activeCells regS (h(l \mapsto A, l'' \mapsto B)) k \end{aligned}$$

**axioms** *activeCells-3-1*:

$$\begin{aligned} & \llbracket l' \notin activeCells regS h k \rrbracket \\ & \implies l' \notin activeCells regS (h(l' \mapsto A, l'' \mapsto B)) k \end{aligned}$$

**axioms** *activeCells-5*:

$$\begin{aligned} & la \notin activeCells regS h k \\ & \implies la \notin activeCells regS' h' k' \end{aligned}$$

**axioms** *equivS-DECREGION*:

$$\begin{aligned} & equivS S S' n ctm d g \\ & \implies equivS S S' n ctm d (g \mid' \{l \in dom hm. fst (the (hm l)) \leq nat (the-Intg (the (sha (heapC, k0f))))\}) \end{aligned}$$

**axioms** *inj-on-DECREGION*:

$$\begin{aligned} & \llbracket ((hm, k), k0, (l, i), S) = ((H, ka), k0a, PC, Sa); inj-on g (dom H) \rrbracket \\ & \implies inj-on (g \mid' \{l \in dom hm. fst (the (hm l)) \leq nat (the-Intg (the (sha (heapC, k0f))))\}) \\ & \quad (dom (hm \mid' \{p \in dom hm. fst (the (hm p)) \leq k0\})) \end{aligned}$$

**axioms** *equivH-DECREGION*:

$$\begin{aligned} & \llbracket ((hm, k), k0, (l, i), S) = ((H, ka), k0a, PC, Sa); equivH (H, ka) h k' com regS d g \rrbracket \\ & \implies equivH (hm \mid' \{p \in dom hm. fst (the (hm p)) \leq k0\}, k0) (dh(l' \mapsto Arr ty \end{aligned}$$

$ma\ regS')$   
 $(nat\ (the-Intg\ (the\ ((sh((heapC,\ kf)\ \mapsto\ the\ (sh\ (heapC,\ kOf))))\ (heapC,\ kf))))\ com\ regS'\ d$   
 $(g\ |\ ' \{l \in dom\ hm.\ fst\ (the\ (hm\ l)) \leq\ nat\ (the-Intg\ (the\ (sha\ (heapC,\ kOf))))\})$

**lemma** *execSVMInstr-DECREGION* :

$\llbracket (P',\ cdm,\ ctm,\ com) = trSVM2JVM\ ((svms,\ ctmap),\ ini,\ ct,\ ah,\ ai,\ bc);$   
 $cdm,\ ctm,\ com \vdash ((hm,\ k),\ k0,\ (l,\ i),\ S) \hat{=} S1';$   
 $(fst\ (the\ (map-of\ svms\ l))\ !\ i) = DECREGION;$   
 $execSVMInst\ DECREGION\ (map-of\ ct)\ (hm,\ k)\ k0\ (l,\ i)\ S = Either.Right$   
 $S2;$   
 $drop\ (the\ (cdm\ (l,\ i)))\ (extractBytecode\ P') =$   
 $trInstr\ (the\ (cdm\ (l,\ i)))\ cdm'\ ctm'\ com\ pcc\ DECREGION\ @\ bytecode'$   
 $\rrbracket \implies \exists\ v'\ sh'\ dh'\ ih'\ fms' .$   
 $P' \vdash S1' -jvm \rightarrow (v',sh',dh',ih',fms') \wedge$   
 $cdm,\ ctm,\ com \vdash S2 \hat{=} (v',sh',dh',ih',fms')$

**apply** (*case-tac*  $S1'$ )  
**apply** (*rename-tac*  $v\ tup$ )  
**apply** (*case-tac*  $tup$ )  
**apply** (*rename-tac*  $sh\ tup$ )  
**apply** (*case-tac*  $tup$ )  
**apply** (*rename-tac*  $dh\ tup$ )  
**apply** (*case-tac*  $tup$ )  
**apply** (*rename-tac*  $ih\ fms$ )  
**apply** (*simp*)

**apply** (*unfold equivState-def*)  
**apply** (*elim exE,elim conjE*)

**apply** (*rule-tac*  $x=None$  **in**  $exI$ )  
**apply** (*rule-tac*  $x=sh((heapC,kf)\ \mapsto\ the\ (sh\ (heapC,\ kOf)))$  **in**  $exI$ )  
**apply** (*rule-tac*  $x=dh(l' \mapsto Arr\ ty\ ma\ regS')$  **in**  $exI$ )  
**apply** (*rule-tac*  $x=ih$  **in**  $exI$ )  
**apply** (*rule-tac*  $x=[([],\ vs,\ safeP,\ sigSafeMain,\ pc+1,\ ref)]$  **in**  $exI$ )

**apply** (*rule conjI*)

**apply** (*subgoal-tac*  $sha\ (heapC,\ kOf) = Some\ (Intg\ ii)$ )  
**apply** (*subgoal-tac*  $sha\ (heapC,\ kf) = Some\ (Intg\ jj)$ )

**apply** (*subgoal-tac*  $P' \mid - (v, sh, dh, ih, fms) -jvm \rightarrow$   
 $(None, sh ((heapC, kf) \mapsto Intg (the-Intg (the (sha (heapC, k0f)))))$ ,  
 $dh (l' \mapsto Arr ty ma regS')$ ,  
 $ih, ([], vs, safeP, sigSafeMain, pc+1, ref) \# []$ )  
 $\wedge (\forall j \leq nat (the-Intg (the (sha (heapC, k0f)))) . region regS' dh j = region$   
 $regS dh j)$ )  
**apply** (*elim conjE*) **apply**(*simp*)

**apply** (*rule decregion-method*)  
**apply** (*simp, simp, simp, simp, simp*)

**apply** (*unfold extractBytecode-def*)  
**apply** (*unfold trInstr.simps*)  
**apply** *clarify*  
**apply** (*rule sym*)  
**apply** (*rule-tac*  $ys=[]$  **in** *nth-via-drop-append*)  
**apply** *simp*  
**apply** (*rule kf-Some-i*)  
**apply** (*rule k0f-Some-i*)

**apply** (*frule nonJumping-Suc-pc*)  
**apply** (*erule-tac sym* [**where**  $t=DECREGION$ ])  
**apply** (*simp add: nonJumping.simps*)  
**apply** (*simp add: trInstr.simps*)

**apply** (*rule-tac*  $x=hm \mid \{p \in dom hm. fst (the (hm p)) \leq k0\}$  **in** *exI*)  
**apply** (*rule-tac*  $x=k0$  **in** *exI*)  
**apply** (*rule-tac*  $x=k0$  **in** *exI*)  
**apply** (*rule-tac*  $x=(l, Suc i)$  **in** *exI*)  
**apply** (*rule-tac*  $x=S$  **in** *exI*)

**apply** (*rule-tac*  $x=sh((heapC, kf) \mapsto the (sh (heapC, k0f)))$  **in** *exI*)  
**apply** (*rule-tac*  $x=dh(l' \mapsto Arr ty ma regS')$  **in** *exI*)  
**apply** (*rule-tac*  $x=ih$  **in** *exI*)  
**apply** (*rule-tac*  $x=vs$  **in** *exI*)  
**apply** (*rule-tac*  $x=pc + 1$  **in** *exI*)  
**apply** (*rule-tac*  $x=ref$  **in** *exI*)

**apply** (*rule-tac*  $x=nat (the-Intg (the ((sh((heapC, kf) \mapsto the (sh (heapC, k0f))))$   
 $(heapC, kf))))$  **in** *exI*)  
**apply** (*rule-tac*  $x=nat (the-Intg (the ((sh((heapC, kf) \mapsto the (sh (heapC, k0f))))$   
 $(heapC, k0f))))$  **in** *exI*)

**apply** (*rule-tac*  $x=la$  **in** *exI*)  
**apply** (*rule-tac*  $x=ty$  **in** *exI*)  
**apply** (*rule-tac*  $x=m$  **in** *exI*)  
**apply** (*rule-tac*  $x=S'$  **in** *exI*)

```

apply (rule-tac x=n in exI)
apply (rule-tac x=l' in exI)
apply (rule-tac x=regS' in exI)
apply (rule-tac x=l'' in exI)
apply (rule-tac x=ma in exI)
apply (rule-tac x=m'' in exI)
apply (rule-tac x=d in exI)
apply (rule-tac x=g |' {l ∈ dom hm. fst (the (hm l)) ≤ nat (the-Intg (the (sha
(heapC, k0f))))}) in exI)

```

```

apply (rule conjI, clarsimp)
apply (rule conjI, rule refl)
apply (rule conjI, rule refl)
apply (rule conjI, rule refl)
apply (rule conjI) apply (simp add: stackC-def add: heapC-def)
apply (rule conjI, assumption)
apply (rule conjI) apply clarsimp
apply (rule conjI) apply (rule activeCells-5,assumption)
apply (rule conjI) apply (rule activeCells-5,assumption)
      apply (rule activeCells-5,assumption)
apply (rule conjI) apply clarsimp
apply (rule conjI) apply (simp add: stackC-def add: heapC-def)
apply (rule conjI) apply (simp add: regionsf-def add: kf-def)
apply (rule conjI) apply clarsimp
apply (rule conjI) apply (simp add: dirCellC-def add: heapC-def)
apply (rule conjI) apply clarsimp
apply (rule conjI) apply clarsimp
apply (rule conjI) apply (simp add: inj-on-def)
apply (rule conjI) apply (rule equivH-DECREGION, assumption+)
apply (rule conjI) apply simp
apply (rule conjI) apply (clarsimp, rule equivS-DECREGION, assumption)
apply (rule conjI) apply clarsimp
by simp

```

**end**

```

theory dem-BUILDCLS
imports ../JVMSAFE/JVMExec SVM2JVM SVMSemantics CertifSVM2JVM
      dem-translation

```

**begin**

```

no-translations Norm s == (None,s)
no-translations ex-table-of m == snd (snd (snd m))

```

```

declare trInstr.simps [simp del]
declare equivS.simps [simp del]
declare extractBytecode-def [simp del]
declare initConsTable-def [simp del]
declare exec-instr.simps [simp]

```

**lemma** *drop-Suc-2*:

```

  drop n xs = (y # ys) ==> drop (Suc n) xs = ys
apply (induct xs arbitrary: n, simp)
apply (simp add: drop-Cons nth-Cons split:nat.splits)
done

```

**axioms** *reserveCell-method*:

```

[[ s = (None, shp, hp, ihp, [([], loc, safeP, sigSafeMain, pc, ref)]);
  k' = nat (the-Intg (the (sha (heapC, kf))));
  shp (stackC, Sf) = Some (Addr l);
  shp (heapC, regionsf) = Some (Addr l');
  shp (dirCellC, safeDirf) = Some (Addr l''); hp l'' = Some (Arr ty m'' d);
  fst(snd(snd(snd(snd(the(method' (P', safeP) sigSafeMain)))))) ! pc =
  Invoke-static dirCellC reserveCell []
]] ==>
P' |- s -jvm -> (None, shp, hp', ihp, [( [Intg ii], loc, safeP, sigSafeMain, pc+1, ref)])
^
  d (nat ii) = Some (Addr lobj) ^
  hp' = hp (lobj ↦ obj) ^
  lobj ∉ activeCells regS h k' ^
  lobj ∉ {l, l', l''} ^
  the-obj obj = (cellC, empty) ^
  ii < int m'' ^ 0 < ii

```

**axioms** *insertCell-method-jvm*:

```

[[ s = (None, shp, hp(lobj ↦ obj), ihp, [( [Intg iloc, Intg (int j)], vs, safeP, sigSafe-
Main, pc, ref)]);
  shp (heapC, regionsf) = Some (Addr l'); hp l' = Some (Arr ty m' regS);
  shp (dirCellC, safeDirf) = Some (Addr l''); hp l'' = Some (Arr ty m'' d);
  d (nat iloc) = Some (Addr lobj);
  fst(snd(snd(snd(snd(the(method' (P', safeP) sigSafeMain)))))) ! pc =
  Invoke-static heapC insertCell [PrimT Integer, PrimT Integer]
]] ==>
P' |- s -jvm -> (None, shp, hp(lobj ↦ obj, l' ↦ Arr ty m' regS'), ihp, [( [], vs,
safeP, sigSafeMain, pc+1, ref)])

```

**axioms** *insertCell-method-equivState*:

```

[[ sh (stackC, Sf) = Some (Addr l); dh l = Some (Arr ty m S');
  sh (stackC, topf) = Some (Intg n);
  sh (heapC, regionsf) = Some (Addr l'); dh l' = Some (Arr ty m' regS);

```

$$\begin{aligned}
& sh \ (dirCellC, safeDirf) = Some \ (Addr \ l''); \ dh \ l'' = Some \ (Arr \ ty \ m'' \ d); \\
& d \ (nat \ ii) = Some \ (Addr \ lobj); \\
& fst(snd(snd(snd(snd(the(method' \ (P', safeP) \ sigSafeMain)))))) \ ! \ pc = \\
& \ Invoke-static \ heapC \ insertCell \ [PrimT \ Integer, \ PrimT \ Integer]; \\
& P' \ |- \\
& \quad (None, \ sh, \ dh(lobj \ \mapsto \ obj), \ ih, \ [([ \ Intg \ ii, \ Intg \ (int \ j)], vs, safeP, \ sigSafe- \\
Main, pc, ref)]) \\
& \quad -jvm \ \rightarrow \\
& \quad (None, \ sh, \ dh(lobj \ \mapsto \ obj, \ l' \ \mapsto \ Arr \ ty \ m' \ regS'), \ ih, \ [([], \ vs, \ safeP, \ sigSafeMain, pc+1, ref)]) \\
& \quad \parallel \ \Longrightarrow \\
& \quad region \ regS' \ (dh(lobj \ \mapsto \ obj, \\
& \quad \quad \quad l' \ \mapsto \ Arr \ ty \ m' \ regS', \\
& \quad \quad \quad l \ \mapsto \ Arr \ ty \ m \ (S'(nat \ (n + 1) \ \mapsto \ Intg \ ii))) \ j = region \ regS \ dh \ j \\
& \cup \ \{lobj\} \ \wedge \\
& \quad (\forall \ j''. \ j'' \neq j \ \longrightarrow \\
& \quad \quad region \ regS' \\
& \quad \quad \quad (dh(lobj \ \mapsto \ obj, \\
& \quad \quad \quad \quad l' \ \mapsto \ Arr \ ty \ m' \ regS', \\
& \quad \quad \quad \quad l \ \mapsto \ Arr \ ty \ m \ (S'(nat \ (n + 1) \ \mapsto \ Intg \ ii)))) \\
& \quad \quad \quad j'' = region \ regS \ dh \ j'')
\end{aligned}$$

**fun** *Item2val* :: *entries*  $\Rightarrow$  *sheap*  $\Rightarrow$  *Item*  $\Rightarrow$  *val*

**where**

$$\begin{aligned}
& \ Item2val \ S \ sh \ (ItemConst \ v) \quad = \ (if \ (isBool \ v = True) \\
& \quad \quad \quad then \ (if \ (the-BoolT \ v) \\
& \quad \quad \quad \quad then \ Intg \ 1 \\
& \quad \quad \quad \quad else \ Intg \ 0) \\
& \quad \quad \quad else \ Intg \ (the-IntT \ v)) \\
& \ | \ Item2val \ S \ sh \ (ItemVar \ l) \quad = \ the \ (S \ (nat \ (the-Intg \ (the \ (sh \ (stackC, \ topf)))) \\
& \quad - \ l)) \\
& \ | \ Item2val \ S \ sh \ (ItemRegSelf) \quad = \ the \ (sh \ (heapC, \ kf))
\end{aligned}$$

**axioms** *fillAux-instrs*:

$$\begin{aligned}
& P \ \vdash \ (None, \ sha, \ h(lobj \ \mapsto \ Obj \ cellC \ [(tagGf, \ cellC) \ \mapsto \ Intg \ (int \ (the \ (com \ C)))]), \\
& \quad \quad \quad inih, \ [([], \\
& \quad \quad \quad vs[Suc \ 0 \ := \ Intg \ ii, \\
& \quad \quad \quad \quad 2 \ := \ Addr \ lobj, \\
& \quad \quad \quad \quad 3 \ := \ k-fresh], \\
& \quad \quad \quad safeP, \ sigSafeMain, \ pc, \ ref)]) \\
& \quad -jvm \ \rightarrow \\
& \quad (None, \ sha, \\
& \quad \quad (let \ vitems = Intg \ (int \ (the \ (com \ C))) \ # \ map \ (Item2val \ S' \ sha) \ items; \\
& \quad \quad \quad cnames = (tagGf, \ cellC) \ # \ map \ (\lambda \ i. \ (VName \ ('arg''@ \ nat2Str \ i), \ cellC)) \\
[0..<length \ items]; \\
& \quad \quad \quad objs \ = \ Obj \ cellC \ (empty(cnames \ [\mapsto] \ vitems))
\end{aligned}$$

$in\ h(lobj \mapsto objs\ ),$   
 $inih,$   
 $[(\ [],$   
 $\quad vs[Suc\ 0 := Intg\ ii,$   
 $\quad\quad 2 := Addr\ lobj,$   
 $\quad\quad 3 := k-fresh],$   
 $\quad safeP, sigSafeMain, pc + length\ (concat\ (map\ fillAux\ (zip\ items\ [1..<length$   
 $items + 1])))], ref)])$

**axioms** *regAux-instrs*:

$P \vdash$   
 $(None, sha, h, inih,$   
 $[(\ [],$   
 $\quad vs[Suc\ 0 := Intg\ ii, 2 := Addr\ lobj],$   
 $\quad safeP, sigSafeMain, pc, ref)])$   
 $-jvm \rightarrow$   
 $(None, sha, h, inih,$   
 $[(\ [],$   
 $\quad vs[Suc\ 0 := Intg\ ii, 2 := Addr\ lobj, 3 := Intg\ j],$   
 $\quad safeP, sigSafeMain, pc + length\ (regAux\ item), ref)])$

**axioms** *arrayIndex-S'-valid*:

$\llbracket h\ l = Some\ (Arr\ ty\ m\ S');$   
 $\quad shp\ (stackC, topf) = Some\ (Intg\ n) \rrbracket$   
 $\implies n + 1 < int\ m \wedge 0 \leq n + 1$

**lemma** *S2-BUILDCLS*:

$execSVMInst\ (BUILDCLS\ C\ items\ item)\ (map-of\ ct)\ (hm, k)\ k0$   
 $(l, i)\ S = Either.Right\ S2$   
 $\implies \exists\ j'.$   
 $(hm(getFresh\ hm \mapsto (j', C, map\ (item2Val\ S)\ items)), k), k0,$   
 $(l, Suc\ i), Val\ (Val.Loc\ (getFresh\ hm)) \# S = S2 \wedge$   
 $j' \leq k$

**apply**  $(unfold\ execSVMInst.simps)$   
**apply**  $(case-tac\ item2Stack\ (snd\ (hm, k))\ S\ item, simp-all)$   
**apply**  $(insert\ RightNotUndefined)$   
**apply**  $(erule-tac\ x = S2\ in\ allE, force)$   
**apply**  $(split\ split-if-asm, simp)$   
**apply**  $(rename-tac\ j)$   
**apply**  $(rule-tac\ x=j\ in\ exI, simp)$   
**apply**  $simp$   
**by**  $(erule-tac\ x = S2\ in\ allE, force)$

**lemma** *equivS-n-ge-minus-1*:

$equivS\ S\ S'\ n\ ctm\ d\ g \implies n \geq -1$

```

apply (case-tac S)
  apply (simp add: equivS.simps)
apply (case-tac a)
  apply (simp add: equivS.simps)
  apply (simp add: equivS.simps)
apply (simp add: equivS.simps)
apply (case-tac x)
by (simp add: equivS.simps)

```

**axioms** *equivS-ge-n-distinct-g*:  
 $equivS\ S\ S'\ n\ ctm\ d\ g \implies (\forall\ p > 0. equivS\ S\ (S'(nat\ (p + n) \mapsto A))\ n\ ctm\ d\ g')$

**lemma** *equivS2-BUILDCLS*:  
 $\llbracket\ equivS\ S\ S'\ n\ ctm\ d\ g;\ d\ (nat\ ii) = Some\ (Addr\ lobj)\ \rrbracket$   
 $\implies equivS\ (Val\ (Val.Loc\ (getFresh\ hm))\ \# S)\ (S'(nat\ (n + 1) \mapsto Intg\ ii))\ (n + 1)$   
 $ctm\ d\ (g\ (getFresh\ hm \mapsto lobj))$

```

apply (frule equivS-n-ge-minus-1)
apply (simp add: equivS.simps)
apply (frule equivS-ge-n-distinct-g)
apply (erule-tac x=1 in allE) apply clarsimp
apply (subgoal-tac nat (1 + n) = nat (n + 1))
apply force
by arith

```

**axioms** *maxPush-BUILDCLS*:  
 $n + 1 < int\ m$

**axioms** *getFresh-notin-dom-h*:  
 $getFresh\ hm \notin dom\ hm$

**axioms** *lobj-notin-ran-g*:  
 $lobj \notin ran\ g$

**declare** *append.simps* [simp del]

**axioms** *S-good*:  
 $\llbracket\ equivS\ S\ S'\ (the\ Intg\ (the\ (sh\ (stackC,\ topf))))\ ctm\ d\ g;\ execSVMInst\ (BUILDCLS\ C\ items\ item)\ (map-of\ ct)\ (H,\ k)$   
 $(nat\ (the\ Intg\ (the\ (sh\ (heapC,\ kOf))))\ (l,\ i)\ S = Either.Right\ S2\ \rrbracket$   
 $\implies (\forall\ i < length\ items. \forall\ v. items!i = ItemConst\ v \implies (\forall\ v'. v \neq Val.Loc\ v')) \wedge$   
 $(\forall\ i < length\ items. \forall\ n. items!i = ItemVar\ n \implies (\forall\ z. S!+n \neq Cont\ z))$



$\wedge$   
 $(\forall i < \text{length items}. \forall n . \text{items!}i = \text{ItemVar } n \longrightarrow (\forall r . S!+n \neq \text{Reg } r)) \wedge$   
 $(\forall i < \text{length items}. \text{items!}i \neq \text{ItemRegSelf}) \wedge$   
 $(\forall n < \text{length } S. \forall z . S!+n \neq \text{Cont } z) \wedge$   
 $(\forall i < \text{length items}. \forall n . \text{items!}i = \text{ItemVar } n \longrightarrow n < \text{nat } (\text{the-Intg } (\text{the}$   
 $(\text{sh } (\text{stackC}, \text{topf})))))) \wedge n < \text{length } S$

**axioms** *map-argCell*:

$\text{argCell } [(\text{tagGf}, \text{cellC}) \mapsto \text{Intg } (\text{int } (\text{the } (\text{com } C))),$   
 $\text{map } (\lambda i. (\text{VName } ("arg" @ \text{nat2Str } i), \text{cellC})) [0..<\text{length items}] [\mapsto]$   
 $\text{map } (\text{Item2val } S' \text{ sha } \text{items}) i =$   
 $\text{argCell } [(\text{VName } ("arg" @ \text{nat2Str } i), \text{cellC}) \mapsto \text{Item2val } S' \text{ sha } (\text{items!}i)] i$

**axioms** *argCell-i*:

$\text{argCell } [(\text{VName } ("arg" @ \text{nat2Str } i), \text{cellC}) \mapsto A] i = A$

**declare** *equivV.simps* [*simp del*]

**axioms** *equivS-nth-2*:

$[[ \text{equivS } S S' \text{ ntop } \text{ctm } d g;$   
 $\forall i < \text{length } S . \forall z . S!+i \neq \text{Cont } z;$   
 $i < \text{length } S;$   
 $S!+i = \text{Val } v]] \implies \text{equivV } v (\text{the } (S' (\text{nat } \text{ntop} - i))) d g$

**axioms** *equivV-distinct-g*:

$\text{equivV } v (\text{the } (S' (\text{nat } (\text{the-Intg } (\text{the } (\text{sh } (\text{stackC}, \text{topf})))) - n))) d g$   
 $\implies \text{equivV } v (\text{the } (S' (\text{nat } (\text{the-Intg } (\text{the } (\text{sh } (\text{stackC}, \text{topf})))) - n))) d g'$

**lemma** *equivV-items*:

$[[ i < \text{length items};$   
 $\text{equivS } S S' (\text{the-Intg } (\text{the } (\text{sh } (\text{stackC}, \text{topf})))) \text{ctm } d g;$   
 $\forall i < \text{length items}. \forall v . \text{items!}i = \text{ItemConst } v \longrightarrow (\forall v'. v \neq \text{Val.Loc } v');$   
 $\forall i < \text{length items}. \forall n . \text{items!}i = \text{ItemVar } n \longrightarrow (\forall z . S!+n \neq \text{Cont } z);$   
 $\forall i < \text{length items}. \forall n . \text{items!}i = \text{ItemVar } n \longrightarrow (\forall r . S!+n \neq \text{Reg } r);$   
 $\forall i < \text{length items}. \text{items!}i \neq \text{ItemRegSelf};$   
 $\forall n < \text{length } S. \forall z . S!+n \neq \text{Cont } z;$   
 $\forall i < \text{length items}. \forall n . \text{items!}i = \text{ItemVar } n \longrightarrow n < \text{nat } (\text{the-Intg } (\text{the } (\text{sh}$   
 $(\text{stackC}, \text{topf})))) \wedge n < \text{length } S]]$   
 $\implies$   
 $\text{equivV } (\text{item2Val } S (\text{items } ! i))$   
 $(\text{argCell}$   
 $[(\text{tagGf}, \text{cellC}) \mapsto \text{Intg } (\text{int } (\text{the } (\text{com } C))), \text{map } (\lambda i. (\text{VName } ("arg"$   
 $@ \text{nat2Str } i), \text{cellC})) [0..<\text{length items}] [\mapsto]$   
 $\text{map } (\text{Item2val } S' \text{ sha } \text{items})$   
 $i)$   
 $d$   
 $(g(\text{getFresh } H \mapsto \text{lobj}))$   
**apply** (*case-tac* (*items* ! *i*))

```

apply simp
apply (case-tac Val)

apply force

apply (simp add: equivV.simps)
apply (subst map-argCell, simp)
apply (simp add: isBool-def)
apply (subst argCell-i, simp)

apply (simp add: equivV.simps)
apply (subst map-argCell, simp)
apply (simp add: isBool-def)
apply (subst argCell-i, simp)
apply (rule impI)
apply (subst argCell-i, simp)

apply (rename-tac n)
apply simp
apply (case-tac S !+ n)

apply simp
apply (rotate-tac 6)
apply (erule-tac x=i in allE) apply simp apply (elim conjE)
apply (frule equivS-nth-2)
apply force
apply assumption
apply assumption
apply (subst map-argCell)
apply clarsimp
apply (subst argCell-i)
apply (erule equivV-distinct-g)

apply force

apply force

by force

declare equivV.simps [simp]

axioms com-C:
  com C = Some (nat (the-Intg
    (the ([[tagGf, cellC]  $\mapsto$  Intg (int (the (com C)))],
      map ( $\lambda i. (VName ('arg' @ nat2Str i), cellC)$ ) [0..length
items] [ $\mapsto$ ]
        map (Item2val S' sha) items]
      (tagGf, cellC))))))
```

**lemma** *domH-BUILDCLS*:

$$\begin{aligned}
& \llbracket \text{ran } g = \text{activeCells } \text{regS } h \text{ (nat (the-Intg (the (sha (heapC, kf)))))); \text{finite} \\
& (\text{dom } H); \text{dom } g = \text{dom } H; \\
& j \leq (\text{nat (the-Intg (the (sha (heapC, kf))))}); \\
& \text{activeCells } \text{regS } h \text{ (nat (the-Intg (the (sha (heapC, kf))))))} \cap \{la, l', l''\} = \{\}; \\
& \text{lobj} \notin \text{activeCells } \text{regS } h \text{ (nat (the-Intg (the (sha (heapC, kf))))))} \wedge \text{lobj} \notin \\
& \{la, l', l''\}; \\
& \text{equivS } Sa \text{ } S' \text{ (the-Intg (the (sha (stackC, topf))))} \text{ ctm } d \text{ } g; \\
& \forall i < \text{length items}. \forall v . \text{items!}i = \text{ItemConst } v \longrightarrow (\forall v' . v \neq \text{Val.Loc } v'); \\
& \forall i < \text{length items}. \forall n . \text{items!}i = \text{ItemVar } n \longrightarrow (\forall z . Sa!+n \neq \text{Cont } z); \\
& \forall i < \text{length items}. \forall n . \text{items!}i = \text{ItemVar } n \longrightarrow (\forall r . Sa!+n \neq \text{Reg } r); \\
& \forall i < \text{length items}. \text{items!}i \neq \text{ItemRegSelf}; \\
& \forall n < \text{length } Sa. \forall z . Sa!+n \neq \text{Cont } z; \\
& \forall i < \text{length items}. \forall n . \text{items!}i = \text{ItemVar } n \longrightarrow n < \text{nat (the-Intg (the} \\
& \text{(sha (stackC, topf))))} \wedge n < \text{length } Sa; \\
& \forall j'' . j'' \neq j \longrightarrow \\
& \quad \text{region } \text{regS}' (h(\text{lobj} \mapsto \text{Obj cellC} \\
& \quad \quad \quad [(\text{tagGf}, \text{cellC}) \mapsto \text{Intg (int (the (com } C))], \\
& \quad \quad \quad \text{map } (\lambda i. (\text{VName ("arg" @ nat2Str } i), \text{cellC})) [0..<\text{length items}]] \\
& [\mapsto] \\
& \quad \quad \quad \text{map (Item2val } S' \text{ sha) items}], \\
& \quad \quad \quad l' \mapsto \text{Arr ty } m' \text{regS}', la \mapsto \text{Arr ty } m \text{ (S'(nat (n + 1) \mapsto Intg ii))} \text{) } j'' = \\
& \text{region } \text{regS } h \text{ } j''; \\
& \quad \quad \quad \text{region } \text{regS}' (h(\text{lobj} \mapsto \text{Obj cellC} \\
& \quad \quad \quad [(\text{tagGf}, \text{cellC}) \mapsto \text{Intg (int (the (com } C))], \\
& \quad \quad \quad \text{map } (\lambda i. (\text{VName ("arg" @ nat2Str } i), \text{cellC})) [0..<\text{length} \\
& \text{items}]] [\mapsto] \\
& \quad \quad \quad \text{map (Item2val } S' \text{ sha) items}], \\
& \quad \quad \quad l' \mapsto \text{Arr ty } m' \text{regS}', la \mapsto \text{Arr ty } m \text{ (S'(nat (n + 1) \mapsto Intg ii))} \text{) } j = \\
& \text{region } \text{regS } h \text{ } j \cup \{\text{lobj}\}; \\
& \quad \quad \quad \forall l \in \text{dom } H. \exists l'. l' = \text{the } (g \text{ } l) \wedge \text{equivC (the (H } l)) h \text{ } l' \text{ (the (h } l')) (nat} \\
& \text{(the-Intg (the (sha (heapC, kf))))} \text{) com } \text{regS } d \text{ } g]] \\
& \implies \forall l \in \text{dom } (H(\text{getFresh } H \mapsto (j, C, \text{map (item2Val } Sa) \text{ items}))). \\
& \quad \quad \quad \exists l'a. l'a = \text{the } ((g(\text{getFresh } H \mapsto \text{lobj})) l) \wedge \\
& \quad \quad \quad \text{equivC (the } ((H(\text{getFresh } H \mapsto (j, C, \text{map (item2Val } Sa) \text{ items}))) l)) \\
& \quad \quad \quad (h(\text{lobj} \mapsto \\
& \quad \quad \quad \text{Obj cellC} \\
& \quad \quad \quad [(\text{tagGf}, \text{cellC}) \mapsto \text{Intg (int (the (com } C))], \\
& \quad \quad \quad \text{map } (\lambda i. (\text{VName ("arg" @ nat2Str } i), \text{cellC})) [0..<\text{length} \\
& \text{items}]] [\mapsto] \\
& \quad \quad \quad \text{map (Item2val } S' \text{ sha) items}], \\
& \quad \quad \quad l' \mapsto \text{Arr ty } m' \text{regS}', la \mapsto \text{Arr ty } m \text{ (S'(nat (n + 1) \mapsto Intg ii))} \text{) } \\
& \quad \quad \quad l'a \text{ (the } ((h(\text{lobj} \mapsto \\
& \quad \quad \quad \text{Obj cellC} \\
& \quad \quad \quad [(\text{tagGf}, \text{cellC}) \mapsto \text{Intg (int (the (com } C))], \\
& \quad \quad \quad \text{map } (\lambda i. (\text{VName ("arg" @ nat2Str } i), \text{cellC})) [0..<\text{length} \\
& \text{items}]] [\mapsto] \\
& \quad \quad \quad \text{map (Item2val } S' \text{ sha) items}],
\end{aligned}$$

```

      l'  $\mapsto$  Arr ty m' regS', la  $\mapsto$  Arr ty m (S'(nat (n + 1)  $\mapsto$ 
Intg ii))))
      l'a))
      (nat (the-Intg (the (sha (heapC, kf)))))) com regS' d (g(getFresh H
 $\mapsto$  lobj))
apply (rule ballI)
apply (erule-tac x=l in ballE)
apply (elim exE, elim conjE)
apply (case-tac l  $\neq$  getFresh H)
apply (rule-tac x=l'a in exI)
apply (rule conjI)
apply clarsimp
apply simp
apply (elim exE) apply (elim conjE)
apply (elim exE) apply (elim conjE)
apply (subgoal-tac the (g l)  $\neq$  lobj  $\wedge$  the (g l)  $\neq$  l'  $\wedge$  the (g l)  $\neq$  la,simp)
prefer 2
apply (simp add: activeCells-def)
apply (erule-tac x=ja in allE)+ apply simp
apply blast
apply (rule-tac x=Obj in exI)
apply (rule-tac x=flds in exI)
apply (rule conjI, rule refl)
apply (rule conjI, rule refl)
apply (rule conjI)
apply clarsimp
apply (erule-tac x=ja in allE)
apply (case-tac ja  $\neq$  j)
apply clarsimp
apply clarsimp
apply (rule allI)
apply (erule-tac x=i in allE)+
apply (rule impI) apply simp
apply (case-tac (vs ! i))
apply simp
apply (elim conjE)
apply (subgoal-tac nata  $\neq$  getFresh H,simp)
apply (subgoal-tac getFresh H  $\notin$  dom g) apply force
apply (subgoal-tac getFresh H  $\notin$  dom H) apply simp
apply (rule getFresh-notin-dom-h)
apply simp
apply simp

apply simp
apply (rule-tac x=Obj in exI)
apply (rule-tac x=[(tagGf, cellC)  $\mapsto$  Intg (int (the (com C))), map ( $\lambda$ i. (VName
("arg" @ nat2Str i), cellC)) [0..\mapsto]
      map (Item2val S' sha) items] in exI)

```

```

apply (rule conjI, rule refl)
apply (rule conjI)
apply clarsimp

apply (subgoal-tac com C =
      Some (nat (the-Intg
        (the ((tagGf, cellC)  $\mapsto$  Intg (int (the (com C))),
          map ( $\lambda i.$  (VName ("arg" @ nat2Str i), cellC))
[0.. $\text{length items}$ ] [ $\mapsto$ ] map (Item2val S' sha) items]
      (tagGf, cellC))))),simp)

apply (rule com-C)
apply (rule allI)
apply (rule impI)
apply (subgoal-tac equivV (item2Val Sa (items ! i))
      (argCell
        [(tagGf, cellC)  $\mapsto$  Intg (int (the (com C))),
          map ( $\lambda i.$  (VName ("arg" @ nat2Str i), cellC)) [0.. $\text{length}$ 
items] [ $\mapsto$ ]
            map (Item2val S' sha) items]
          i)
      d (g(getFresh H  $\mapsto$  lobj)),simp)
apply (rule equivV-items) apply assumption+ apply clarsimp apply assumption+
apply clarsimp
apply simp
apply simp
apply (rule-tac x=Obj in exI)
apply (rule-tac x=[(tagGf, cellC)  $\mapsto$  Intg (int (the (com C))),
      map ( $\lambda i.$  (VName ("arg" @ nat2Str i), cellC)) [0.. $\text{length}$ 
items] [ $\mapsto$ ]
        map (Item2val S' sha) items] in exI)

apply (rule conjI)
apply (rule refl)
apply (rule conjI)
apply (rule com-C)
apply (rule allI) apply (rule impI)
apply (rule equivV-items) apply assumption+ apply clarsimp apply assumption+
apply clarsimp
by simp

```

**lemma** ran-upd-getFresh:

[ lobj  $\notin$  ran g; x  $\in$  ran g; getFresh H  $\notin$  dom g ]  
 $\implies$  x  $\in$  ran (g(getFresh H  $\mapsto$  lobj))

**apply** (simp add: ran-def)  
**apply** (erule exE)

**apply** (*erule-tac*  $x=a$  **in** *allE,simp*)  
**apply** (*rule-tac*  $x=a$  **in** *exI*)  
**apply** (*rule conjI*)  
**apply** *blast*  
**by** *blast*

**lemma** *disjoint-getFresh*:  
 $\llbracket x \in \text{ran } (g(\text{getFresh } H \mapsto \text{lobj})); \text{lobj} \notin \text{ran } g \rrbracket$   
 $\implies x \in \text{ran } g \vee x = \text{lobj}$   
**apply** (*simp add: ran-def*)  
**apply** (*split split-if-asm, clarsimp*)  
**by** (*rule-tac*  $x=a$  **in** *exI, simp*)

**lemma** *disjoint-getFresh-2*:  
 $\llbracket x \in \text{ran } g \vee x = \text{lobj}; \text{lobj} \notin \text{ran } g; \text{getFresh } H \notin \text{dom } g \rrbracket$   
 $\implies x \in \text{ran } (g(\text{getFresh } H \mapsto \text{lobj}))$   
**apply** (*erule disjE*)  
**apply** (*simp add: ran-def*)  
**apply** (*erule exE*)  
**apply** (*rule-tac*  $x=a$  **in** *exI*)  
**apply** (*rule conjI, clarsimp*)  
**apply** *clarsimp*  
**by** (*simp add: ran-def*)

**axioms** *getFresh-notin-dom-g*:  
 $\text{getFresh } hm \notin \text{dom } g$

**lemma** *activeCells-BUILDCLS*:  
 $\llbracket l' \notin \text{activeCells } \text{regS } dh \text{ (nat (the-Intg (the (sha (heapC, kf))))));$   
 $j \leq \text{nat (the-Intg (the (sha (heapC, kf)))));$   
 $\forall j'. j' \neq j \longrightarrow \text{region } \text{regS}' (h(\text{lobj} \mapsto$   
 $\text{Obj cellC}$   
 $\llbracket (\text{tagGf, cellC}) \mapsto \text{Intg (int (the (com } C)),$   
 $\text{map } (\lambda i. (\text{VName } ('arg'' @ \text{nat2Str } i), \text{cellC})) [0..<\text{length}$   
 $\text{items}] [\mapsto]$   
 $\text{map } (\text{Item2val } S' \text{ sha}) \text{ items}],$   
 $l' \mapsto \text{Arr ty } m' \text{ regS}', la \mapsto \text{Arr ty } m \text{ (S'(nat (n + 1) } \mapsto \text{Intg } ii)) \rrbracket j' =$   
 $\text{region } \text{regS } h j';$   
 $\text{ran } g = \text{activeCells } \text{regS } h \text{ (nat (the-Intg (the (sha (heapC, kf)))));$   
 $\text{region } \text{regS}' (h(\text{lobj} \mapsto$   
 $\text{Obj cellC}$   
 $\llbracket (\text{tagGf, cellC}) \mapsto \text{Intg (int (the (com } C)),$   
 $\text{map } (\lambda i. (\text{VName } ('arg'' @ \text{nat2Str } i), \text{cellC})) [0..<\text{length}$   
 $\text{items}] [\mapsto]$   
 $\text{map } (\text{Item2val } S' \text{ sha}) \text{ items}],$   
 $l' \mapsto \text{Arr ty } m' \text{ regS}', la \mapsto \text{Arr ty } m \text{ (S'(nat (n + 1) } \mapsto \text{Intg } ii)) \rrbracket j =$   
 $\text{insert } \text{lobj } (\text{region } \text{regS } h j) \rrbracket$   
 $\implies \text{ran } (g(\text{getFresh } hm \mapsto \text{lobj})) =$

```

    activeCells regS'
      (h(lobj  $\mapsto$ 
        Obj cellC
          [(tagGf, cellC)  $\mapsto$  Intg (int (the (com C))),
            map ( $\lambda i.$  (VName ("arg" @ nat2Str i), cellC)) [0.. $\text{length}$ 
items] [ $\mapsto$ ]
              map (Item2val S' sha) items],
        l'  $\mapsto$  Arr ty m' regS', la  $\mapsto$  Arr ty m (S'(nat (n + 1)  $\mapsto$  Intg i)))
        (nat (the-Intg (the (sha (heapC, kf)))))
apply (rule equalityI)
apply (rule subsetI)
apply (subgoal-tac lobj  $\notin$  ran g)
  prefer 2 apply (rule lobj-notin-ran-g)
apply (subgoal-tac  $x \in$  ran g  $\vee$  x = lobj, simp)
  prefer 2 apply (rule disjoint-getFresh, assumption+)
apply (erule disjE)

apply (simp add: activeCells-def)
apply (elim exE)
apply (rule-tac x=ja in exI, simp, elim conjE)
apply (erule-tac x=ja in allE)+
apply (case-tac ja = j) apply simp
apply simp
apply simp

apply (simp add: activeCells-def)
apply (rule-tac x=j in exI)
apply (rule conjI)
apply (erule-tac x=getFresh hm in allE)
apply assumption
apply simp

apply (rule subsetI)
apply (subgoal-tac lobj  $\notin$  ran g)
  prefer 2 apply (rule lobj-notin-ran-g)
apply (subgoal-tac getFresh hm  $\notin$  dom g)
  prefer 2 apply (rule getFresh-notin-dom-g)
apply (simp add: activeCells-def)
apply (erule exE) apply (elim conjE)
apply (erule-tac x=ja in allE)+
apply simp
apply (case-tac ja = j)
  apply simp
apply (erule disjE)
apply simp defer
apply (subgoal-tac lobj  $\notin$  ran g)
  prefer 2 apply (rule lobj-notin-ran-g)
apply (subgoal-tac x  $\in$  ran g)
apply (rule ran-upd-getFresh) apply assumption+

```

```

apply (subgoal-tac  $x \in \{p. \exists j \leq \text{nat} (\text{the-Intg} (\text{the} (\text{sha} (\text{heapC}, \text{kf})))\}$ ).  $p \in \text{region}$ 
 $\text{regS } h \ j\}$ )
apply simp
apply clarsimp
apply (subgoal-tac  $x \in \text{region } \text{regS } h \ ja$ )
  prefer 2 apply simp
apply (subgoal-tac  $x \in \text{ran } g$ )
apply (subgoal-tac  $\text{lobj} \notin \text{ran } g$ )
  prefer 2 apply (rule  $\text{lobj-notin-ran-g}$ )
apply (subgoal-tac  $\text{getFresh } hm \notin \text{dom } g$ )
  prefer 2 apply (rule  $\text{getFresh-notin-dom-g}$ )
apply (rule  $\text{disjoint-getFresh-2}$ ) apply simp apply simp apply simp
apply (subgoal-tac  $x \in \{p. \exists j \leq \text{nat} (\text{the-Intg} (\text{the} (\text{sha} (\text{heapC}, \text{kf})))\}$ ).  $p \in \text{region}$ 
 $\text{regS } h \ j\}$ )
apply simp
apply simp
apply (rule-tac  $x=ja$  in  $exI$ )
apply simp
apply (subgoal-tac  $\text{lobj} \notin \text{ran } g$ )
  prefer 2 apply (rule  $\text{lobj-notin-ran-g}$ )
apply (simp add:  $\text{ran-def}$ )
by force

```

**lemma** *equivH-BUILDCLS*:

```

[[ equivH ( $H, ka$ )  $h (\text{nat} (\text{the-Intg} (\text{the} (\text{sha} (\text{heapC}, \text{kf}))))$ ) com  $\text{regS } d \ g; j \leq$ 
 $ka;$ 
   $\text{dom } g = \text{dom } H;$ 
   $\text{activeCells } \text{regS } h (\text{nat} (\text{the-Intg} (\text{the} (\text{sha} (\text{heapC}, \text{kf})))) \cap \{la, l', l''\} = \{\};$ 
   $\text{lobj} \notin \text{activeCells } \text{regS } h (\text{nat} (\text{the-Intg} (\text{the} (\text{sha} (\text{heapC}, \text{kf}))))); \text{lobj} \notin \{la,$ 
 $l', l''\};$ 
   $\text{equivS } Sa \ S' (\text{the-Intg} (\text{the} (\text{sha} (\text{stackC}, \text{topf})))) \text{ctm } d \ g;$ 
   $\forall i < \text{length } \text{items}. \forall v . \text{items!}i = \text{ItemConst } v \longrightarrow (\forall v' . v \neq \text{Val.Loc } v');$ 
   $\forall i < \text{length } \text{items}. \forall n . \text{items!}i = \text{ItemVar } n \longrightarrow (\forall z . Sa!+n \neq \text{Cont } z);$ 
   $\forall i < \text{length } \text{items}. \forall n . \text{items!}i = \text{ItemVar } n \longrightarrow (\forall r . Sa!+n \neq \text{Reg } r);$ 
   $\forall i < \text{length } \text{items}. \text{items!}i \neq \text{ItemRegSelf};$ 
   $\forall n < \text{length } Sa. \forall z . Sa!+n \neq \text{Cont } z;$ 
   $\forall i < \text{length } \text{items}. \forall n . \text{items!}i = \text{ItemVar } n \longrightarrow n < \text{nat} (\text{the-Intg} (\text{the}$ 
 $(\text{sha} (\text{stackC}, \text{topf})))) \wedge n < \text{length } Sa;$ 
   $\forall j'' . j'' \neq j \longrightarrow$ 
     $\text{region } \text{regS}'$ 
    ( $h(\text{lobj} \mapsto \text{Obj } \text{cellC} [(tagGf, \text{cellC}) \mapsto \text{Intg} (\text{int} (\text{com } C))],$ 
       $\text{map } (\lambda i. (\text{VName} ('arg'' @ \text{nat2Str } i), \text{cellC}))$ 
       $[0..<\text{length } \text{items}] \mapsto]$ 
       $\text{map } (\text{Item2val } S' \ \text{sha}) \ \text{items}],$ 
       $l' \mapsto \text{Arr } \text{ty } m' \ \text{regS}',$ 
       $la \mapsto \text{Arr } \text{ty } m \ (S'(\text{nat} (n + 1) \mapsto \text{Intg } ii))))$ 

```



```

      j'' = region regS h j'';
region regS'
  (h(lobj ↦ Obj cellC [(tagGf, cellC) ↦ Intg (int (the (com C))),
                      map (λi. (VName ("arg" @ nat2Str i), cellC))
[0..<length items] [↦]
      map (Item2val S' sha) items],
  l' ↦ Arr ty m' regS',
  la ↦ Arr ty m (S'(nat (n + 1) ↦ Intg ii))) j = region regS h j ∪
{lobj}
  ]
  ⇒ equivH ((H(getFresh H ↦ (j, C, map (item2Val Sa) items))), ka)
    (h(lobj ↦ Obj cellC [(tagGf, cellC) ↦ Intg (int (the (com C))),
                        map (λi. (VName ("arg" @ nat2Str i), cellC))
[0..<length items] [↦]
      map (Item2val S' sha) items],
  l' ↦ Arr ty m' regS',
  la ↦ Arr ty m (S'(nat (n + 1) ↦ Intg ii)))
  (nat (the-Intg (the (sha (heapC, kf))))))
  com regS' d (g(getFresh H ↦ lobj))
apply (simp add: heapC-def add: stackC-def)
apply (fold heapC-def)
apply (unfold equivH.simps, elim conjE)
apply (rule conjI, clarsimp)
apply (rule conjI)
apply (rule activeCells-BUILDCLS,assumption+,simp,assumption+)
apply (rule conjI, simp)
apply (fold stackC-def)
apply (frule-tac la=la and l'=l' and l''=l'' and
  lobj=lobj and C=C and Sa=Sa in domH-BUILDCLS)
by (assumption+,simp+)

```

**lemma** *notin-ran*:

```

  x ∉ ran g ⇒ ∀ y . g y ≠ Some x
by (clarsimp, simp add: ran-def)

```

**lemma** *inj-on-BUILDCLS*:

```

  [ ((hm, k), k0, (l, i), S) = ((H, ka), k0a, PC, Sa);
    inj-on g (dom H) ]
  ⇒ inj-on (g(getFresh hm ↦ lobj)) (dom (hm(getFresh hm ↦ (j', C, map
(item2Val S) items))))
apply (subgoal-tac getFresh hm ∉ dom g)
prefer 2 apply (rule getFresh-notin-dom-g)
apply (subgoal-tac getFresh hm ∉ dom hm)
prefer 2 apply (rule getFresh-notin-dom-h)
apply (subgoal-tac lobj ∉ ran g)

```

**prefer 2 apply** (*rule lobj-notin-ran-g*)  
**apply** *clarsimp*  
**apply** (*simp add: inj-on-def*)  
**apply** (*split split-if-asm*)  
**apply** (*simp add: getFresh-def add: fresh-def, auto*)  
**apply** (*split split-if-asm, auto*)  
**apply** (*frule notin-ran*)  
**apply** (*erule-tac x=x in allE*)  
**by** *simp*

**axioms** *activeCells-3*:  
 $\llbracket l' \notin \text{activeCells } \text{regS } h \ k; l \neq l'; l \neq l'' \rrbracket$   
 $\implies l' \notin \text{activeCells } \text{regS } (h(l \mapsto A, l'' \mapsto B)) \ k$

**axioms** *activeCells-3-1*:  
 $\llbracket l' \notin \text{activeCells } \text{regS } h \ k \rrbracket$   
 $\implies l' \notin \text{activeCells } \text{regS } (h(l' \mapsto A, l'' \mapsto B)) \ k$

**axioms** *activeCells-4*:  
 $la \notin \text{activeCells } \text{regS } h \ k$   
 $\implies la \notin \text{activeCells } \text{regS}' \ h \ k'$

**axioms** *activeCells-5*:  
 $la \notin \text{activeCells } \text{regS } h \ k$   
 $\implies la \notin \text{activeCells } \text{regS}' \ h' \ k'$

**axioms** *drop-append-length-2*:  
 $\text{drop } n \ xs = (ys \ @ \ zs) \ @ \ ms \implies \text{drop } (n + \text{length } ys) \ xs = zs \ @ \ ms$

**declare** *List.upt.upt-Suc* [*simp del*]  
**declare** *execSVMInst.simps* [*simp del*]  
**declare** *equivH.simps* [*simp del*]

**thm** *drop-append-length*

**axioms** *drop-append-length-3*:  
 $\text{drop } n \ xs = ys \ @ \ zs \ @ \ ms \implies \text{drop } (n + \text{length } ys) \ xs = zs \ @ \ ms$

**lemma** *execSVMInstr-BUILDCLS-reserveCell* :  
 $\llbracket (P', \text{cdm}, \text{ctm}, \text{com}) = \text{trSVM2JVM } ((\text{svms}, \text{ctmap}), \text{ini}, \text{ct}, \text{ah}, \text{ai}, \text{bc});$   
 $S1' = (\text{None}, \text{sh}, \text{dh}, \text{ih}, [([], \text{vs}, \text{safeP}, \text{sigSafeMain}, \text{pc}, \text{ref})]);$   
 $k' = \text{nat } (\text{the-Intg } (\text{the } (\text{sha } (\text{heapC}, \text{kf})))));$   
 $(\text{fst } (\text{the } (\text{map-of } \text{svms } p))) ! i = \text{BUILDCLS } C \ \text{items } \text{item};$   
 $\text{execSVMInst } (\text{BUILDCLS } C \ \text{items } \text{item}) \ (\text{map-of } \text{ct}) \ (\text{hm}, \text{k}) \ k0 \ (p, i) \ S =$   
*Either.Right S2*;  
 $\text{drop } (\text{the } (\text{cdm } (p, i))) \ (\text{extractBytecode } P') =$   
 $\text{trInstr } (\text{the } (\text{cdm } (p, i))) \ \text{cdm}' \ \text{ctm}' \ \text{com} \ \text{pcc } (\text{BUILDCLS } C \ \text{items } \text{item}) \ @$

$bytecode'$ ;  
 $sh (stackC, Sf) = Some (Addr l)$ ;  
 $distinct [l, l', l'']$ ;  
 $sh (heapC, regionsf) = Some (Addr l')$ ;  $dh l' = Some (Arr ty m' regS)$ ;  
 $sh (dirCellC, safeDirf) = Some (Addr l'')$ ;  $dh l'' = Some (Arr ty m'' d)$ ;  
 $pc = the (cdm (p, i))$

$\mathbb{1} \implies$   
 $P' \vdash S1' -jvm \rightarrow$   
 $(None, sh, dh', ih, [[Intg ii], vs, safeP, sigSafeMain, pc + 1, ref]]) \wedge$   
 $d (nat ii) = Some (Addr lobj) \wedge$   
 $dh' = dh (lobj \mapsto obj) \wedge$   
 $lobj \notin activeCells regS dh k' \wedge$   
 $lobj \notin \{l, l', l''\} \wedge$   
 $the-obj obj = (cellC, empty) \wedge$   
 $ii < int m'' \wedge 0 < ii$

**apply** (*subgoal-tac*)

$P' \mid -$   
 $(None, sh, dh, ih, [[], vs, safeP, sigSafeMain, the (cdm (p, i)), ref]])$   
 $-jvm \rightarrow$   
 $(None, sh, dh', ih, [[Intg ii], vs, safeP, sigSafeMain, the (cdm (p, i)) + 1, ref]]) \wedge$   
 $d (nat ii) = Some (Addr lobj) \wedge$   
 $dh' = dh (lobj \mapsto obj) \wedge$   
 $lobj \notin activeCells regS dh k' \wedge$   
 $lobj \notin \{l, l', l''\} \wedge$   
 $the-obj obj = (cellC, empty) \wedge$   
 $ii < int m'' \wedge 0 < ii$

**prefer 2 apply** (*rule reserveCell-method*) **apply** (*simp, assumption+*)

**apply** (*unfold extractBytecode-def*)

**apply** (*unfold trInstr.simps*)

**apply** (*unfold Let-def*)

**apply** (*simp, erule nth-via-drop-append*)

**by** *clarsimp*

**lemma** *execSVMInstr-BUILDCLS* :

$\mathbb{1} (P', cdm, ctm, com) = trSVM2JVM ((svms, ctmap), ini, ct, ah, ai, bc);$   
 $cdm, ctm, com \vdash ((hm, k), k0, (l, i), S) \triangleq S1'$ ;  
 $(fst (the (map-of svms l)) ! i) = BUILDCLS C items item;$   
 $execSVMInst (BUILDCLS C items item) (map-of ct) (hm, k) k0 (l, i) S =$

*Either.Right S2*;

$drop (the (cdm (l, i))) (extractBytecode P') =$

$trInstr (the (cdm (l, i))) cdm' ctm' com pcc (BUILDCLS C items item) @$

*bytecode'*

$\mathbb{1} \implies \exists v' sh' dh' ih' fms'$  .

$P' \vdash S1' -jvm \rightarrow (v', sh', dh', ih', fms') \wedge$

$cdm, ctm, com \vdash S2 \triangleq (v', sh', dh', ih', fms')$

```

apply (case-tac  $S1'$ )
apply (rename-tac  $v$   $tup$ )
apply (case-tac  $tup$ )
apply (rename-tac  $sh$   $tup$ )
apply (case-tac  $tup$ )
apply (rename-tac  $dh$   $tup$ )
apply (case-tac  $tup$ )
apply (rename-tac  $ih$   $fms$ )
apply (simp)

apply (subgoal-tac
   $\exists j.$ 
  (( $hm(getFresh\ hm \mapsto (j, C, map(item2Val\ S)\ items)), k), k0,$ 
  ( $l, Suc\ i), Val\ (Val.Loc\ (getFresh\ hm)) \# S = S2 \wedge j \leq k$ )
prefer 2 apply (erule  $S2-BUILDCLS$ )
apply (elim  $exE$ , elim  $conjE$ )

apply (unfold equivState-def)
apply (elim  $exE$ , elim  $conjE$ )

apply (rule-tac  $x=None$  in  $exI$ )
apply (rule-tac  $x=sha((stackC, topf) \mapsto Intg\ (n + 1))$  in  $exI$ )
apply (rule-tac  $x=h(lobj \mapsto Obj\ cellC [(tagGf, cellC) \mapsto Intg\ (int\ (the\ (com\ C))),$ 
   $map\ (\lambda i. (VName\ ('arg' \ @\ nat2Str\ i), cellC))$ 
   $[0..<length\ items] [\mapsto]$ 
   $map\ (Item2val\ S'\ sha)\ items],$ 
   $l' \mapsto Arr\ ty\ m'\ regS',$ 
   $la \mapsto Arr\ ty\ m\ (S'(nat\ (n + 1) \mapsto Intg\ ii)))$  in  $exI$ )
apply (rule-tac  $x=inih$  in  $exI$ )
apply (rule-tac  $x=([[],$ 
   $vs[Suc\ 0 := Intg\ ii, 2 := Addr\ lobj, 3 := Intg\ (int\ j)],$ 
   $safeP, sigSafeMain,$ 
   $Suc\ (Suc\ (Suc\ (Suc\ (Suc\ (Suc\ (Suc\ (Suc\ (Suc\ (Suc\ (Suc$ 
   $(Suc\ (Suc\ (Suc\ (Suc\ (Suc\ (Suc\ (Suc\ (Suc\ (the\ (cdm\ (l, i)) +$ 
   $length\ (regAux\ item) +$ 
   $length\ (concat\ (map\ fillAux\ (zip\ items\ [Suc\ 0..<Suc\ (length$ 
   $items])]]]]]]]]]]]]]]]]]]], ref)]$  in  $exI$ )

apply (subgoal-tac  $\exists vs\ pc . fms=[([[] ,vs, safeP, sigSafeMain, pc, ref)])$ )
prefer 2 apply simp
apply (erule  $exE$ )+

apply (subgoal-tac length  $vs = 10$ )
prefer 2 apply (rule length-vs)

```

```

apply (subgoal-tac PC = (l,i))
prefer 2 apply simp

apply (drule-tac r=S1' in trans,assumption)

apply (rule conjI)

apply (frule-tac dh'=h' and lobj=lobj and obj=obj and ii=ii in execSVMInstr-BUILDCLS-reserveCell)
apply assumption+ apply simp
apply (elim conjE)

apply (unfold exec-all-def)
apply (erule rtrancl-trans)
apply (unfold extractBytecode-def)
apply (unfold trInstr.simps)
apply (unfold Let-def)

apply (drule drop-Suc-append-2)
apply (subgoal-tac
  P⊢ (None,sh,h',ih,[[Intg ii], vs, safeP, sigSafeMain, pc + 1, ref]])
  -jvm→
  (None,sh,h',ih,[[[], vs[Suc 0 := Intg ii], safeP, sigSafeMain, pc + 1 + 1,
ref]]))
apply (unfold exec-all-def)
apply (simp,erule rtrancl-trans)
prefer 2
apply (rule r-into-rtrancl)
apply (clarify)
apply (unfold JVMEexec.exec.simps)
apply (unfold Let-def)
apply (subgoal-tac
  (fst (snd (snd (snd (snd (method' (P', safeP) sigSafeMain)))))) !
  Suc (the (cdm (l, i)))) =
  Store 1,simp)
apply (simp, erule nth-via-drop-append)

apply (drule drop-Suc-append)
apply (subgoal-tac
  P⊢ (None,sh,h',ih,[[[], vs[Suc 0 := Intg ii], safeP, sigSafeMain, pc + 1 + 1,
ref]])
  -jvm→
  (None,sh,h',ih,
  [[Addr l'],
  vs[Suc 0 := Intg ii],
  safeP, sigSafeMain, pc + 1 + 1 + 1, ref]]))

```

```

apply (unfold exec-all-def)
apply (simp,erule rtrancl-trans)
prefer 2
apply (rule r-into-rtrancl)
apply (clarify)
apply (unfold JVMExec.exec.simps)
apply (unfold Let-def)
apply (subgoal-tac
  (fst (snd (snd (snd (snd (the (method' (P', safeP) sigSafeMain)))))) !
    (Suc (Suc (the (cdm (l, i)))))) =
    Getstatic safeDirf dirCellC,simp)
apply (erule nth-via-drop-append)

```

```

apply (drule drop-Suc-append)
apply (subgoal-tac
  P⊢ (None,sh,h',ih,
    [([Addr l'],
      vs[Suc 0 := Intg ii],
      safeP, sigSafeMain, pc + 1 + 1 + 1, ref]))
  -jvm→
  (None,sh,h',ih,
    [([Intg ii,
      Addr l'],
      vs[Suc 0 := Intg ii],
      safeP, sigSafeMain, pc + 1 + 1 + 1 + 1, ref]))))

```

```

apply (unfold exec-all-def)
apply (simp,erule rtrancl-trans)
prefer 2
apply (rule r-into-rtrancl)
apply (clarify)
apply (unfold JVMExec.exec.simps)
apply (unfold Let-def)
apply (subgoal-tac
  (fst (snd (snd (snd (snd (the (method' (P', safeP) sigSafeMain)))))) !
    (Suc (Suc (Suc (the (cdm (l, i)))))) =
    Load (Suc 0),simp)
apply (erule nth-via-drop-append)

```

```

apply (drule drop-Suc-append)
apply (subgoal-tac
  P⊢ (None,sh,h',ih,
    [([Intg ii,
      Addr l'],
      vs[Suc 0 := Intg ii],
      safeP, sigSafeMain, pc + 1 + 1 + 1 + 1, ref]))
  -jvm→
  (None,sh,h',ih,

```

```

      [[Addr lobj],
       vs[Suc 0 := Intg ii],
       safeP, sigSafeMain, pc + 1 + 1 + 1 + 1 + 1, ref]])
apply (unfold exec-all-def)
apply (simp,erule rtrancl-trans)
prefer 2
apply (rule r-into-rtrancl)
apply (clarify)
apply (unfold JVMExec.exec.simps)
apply (unfold Let-def)
apply (subgoal-tac
  (fst (snd (snd (snd (snd (the (method' (P', safeP) sigSafeMain)))))) !
  (Suc (Suc (Suc (Suc (the (cdm (l, i))))))) =
  ArrLoad,simp)
apply (simp add: raise-system-xcpt-def)
apply (erule nth-via-drop-append)

```

```

apply (drule drop-Suc-append)
apply (subgoal-tac
  P' ⊢ (None,sh,h',ih,
  [[Addr lobj],
   vs[Suc 0 := Intg ii],
   safeP, sigSafeMain, pc + 1 + 1 + 1 + 1 + 1, ref]])
  -jvm→
  (None,sh,h',ih,
  [[[],
   vs[Suc 0 := Intg ii,
       2 := Addr lobj],
   safeP, sigSafeMain, pc + 1 + 1 + 1 + 1 + 1 + 1, ref]])
apply (unfold exec-all-def)
apply (simp,erule rtrancl-trans)
prefer 2
apply (rule r-into-rtrancl)
apply (clarify)
apply (unfold JVMExec.exec.simps)
apply (unfold Let-def)
apply (subgoal-tac
  (fst (snd (snd (snd (snd (the (method' (P', safeP) sigSafeMain)))))) !
  (Suc (Suc (Suc (Suc (Suc (the (cdm (l, i))))))) =
  Store 2,simp)
apply (erule nth-via-drop-append)

```

```

apply (drule drop-Suc-append)
apply (subgoal-tac
  P' ⊢ (None,sh,h',ih,
  [[[],
   vs[Suc 0 := Intg ii,

```

```

      2 := Addr lobj],
      safeP, sigSafeMain, pc + 1 + 1 + 1 + 1 + 1 + 1, ref)))
-jvm→
  (None,sh,h',ih,
  [[Addr lobj],
  vs[Suc 0 := Intg ii,
      2 := Addr lobj],
  safeP, sigSafeMain, pc + 1 + 1 + 1 + 1 + 1 + 1 + 1, ref]]))
apply (unfold exec-all-def)
apply (simp,erule rtrancl-trans)
prefer 2
apply (rule r-into-rtrancl)
apply (clarify)
apply (unfold JVMEexec.exec.simps)
apply (unfold Let-def)
apply (subgoal-tac
  (fst (snd (snd (snd (snd (the (method' (P', safeP) sigSafeMain)))))) !
  (Suc (Suc (Suc (Suc (Suc (Suc (the (cdm (l, i)))))))))) =
  Load 2,simp)
apply (erule nth-via-drop-append)

apply (drule drop-Suc-append)
apply (subgoal-tac
  P' ⊢ (None,sh,h',ih,
  [[Addr lobj],
  vs[Suc 0 := Intg ii,
      2 := Addr lobj],
  safeP, sigSafeMain, pc + 1 + 1 + 1 + 1 + 1 + 1 + 1, ref]))
-jvm→
  (None,sh,h',ih,
  [[Intg (int (the (com C))),
  Addr lobj],
  vs[Suc 0 := Intg ii,
      2 := Addr lobj],
  safeP, sigSafeMain, pc + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1, ref]]))
apply (unfold exec-all-def)
apply (simp,erule rtrancl-trans)
prefer 2
apply (rule r-into-rtrancl)
apply (clarify)
apply (unfold JVMEexec.exec.simps)
apply (unfold Let-def)
apply (subgoal-tac
  (fst (snd (snd (snd (snd (the (method' (P', safeP) sigSafeMain)))))) !
  (Suc (Suc (Suc (Suc (Suc (Suc (Suc (the (cdm (l, i)))))))))) =
  LitPush (Intg (int (the (com C))))),simp)
apply (erule nth-via-drop-append)

```



```

apply (drule drop-Suc-append)
apply (subgoal-tac
   $P \vdash$  (None, sh, h', ih,
    [[Intg (int (the (com C))),
      Addr lobj],
    vs[Suc 0 := Intg ii,
      2 := Addr lobj],
    safeP, sigSafeMain, pc + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1, ref]])
  -jvm→
    (None, sh, h'(lobj ↦ Obj cellC (empty((tagGf, cellC) ↦ Intg (int (the (com C))))))), ih,
    [[ $\square$ ],
    vs[Suc 0 := Intg ii,
      2 := Addr lobj],
    safeP, sigSafeMain, pc + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1, ref]])
apply (unfold exec-all-def)
apply (simp, erule rtrancl-trans)
prefer 2
apply (rule r-into-rtrancl)
apply (clarify)
apply (unfold JVMExec.exec.simps)
apply (unfold Let-def)
apply (subgoal-tac
  (fst (snd (snd (snd (snd (the (method' (P', safeP) sigSafeMain)))))) !
  (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (the (cdm (l, i)))))))))) =
  Putfield tagGf cellC, simp))
apply (simp add: raise-system-xcpt-def)
apply (erule nth-via-drop-append)

```

```

apply (drule drop-Suc-append)
apply (subgoal-tac
   $P \vdash$ 
    (None, sh,
      h'(lobj ↦ Obj cellC (empty((tagGf, cellC) ↦ Intg (int (the (com C))))))),
    ih,
    [[ $\square$ ],
    vs[Suc 0 := Intg ii, 2 := Addr lobj],
    safeP, sigSafeMain, pc + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1, ref]])
  -jvm→
    (None, sh,
      h'(lobj ↦ Obj cellC [(tagGf, cellC) ↦ Intg (int (the (com C)))]),
    ih,
    [[ $\square$ ],
    vs[Suc 0 := Intg ii, 2 := Addr lobj, 3 := Intg (int j)],
    safeP, sigSafeMain,
    pc + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + length (regAux item), ref]])
prefer 2 apply (rule regAux-instrs)

```

**apply** (*unfold exec-all-def*)  
**apply** (*simp, erule rtrancl-trans*)

**apply** (*drule drop-append-length-3*)

**apply** (*subgoal-tac*)

$P \vdash$  (*None, sh,*  
 $h'(lobj \mapsto Obj\ cellC [(tagGf, cellC) \mapsto Intg (int (the (com\ C)))]),$   
*ih,*  $[(\ [],$   
 $vs[Suc\ 0 := Intg\ ii,$   
 $2 := Addr\ lobj,$   
 $3 := Intg (int\ j)],$   
*safeP, sigSafeMain,*  
 $pc + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + length (regAux\ item), ref)])$   
 $\rightarrow$   
*(None, sh,*  
 $(let\ vitems = Intg (int (the (com\ C))) \# map (Item2val\ S'\ sh)\ items;$   
 $cnames = (tagGf, cellC) \# map (\lambda\ i. (VName ("arg"@\ nat2Str\ i), cellC))$   
 $[0..<length\ items];$   
 $objs = Obj\ cellC (empty(cnames\ [\mapsto]\ vitems))$   
 $in\ h'(lobj \mapsto objs)),$   
*ih,*  
 $[(\ [],$   
 $vs[Suc\ 0 := Intg\ ii,$   
 $2 := Addr\ lobj,$   
 $3 := Intg (int\ j)],$   
*safeP, sigSafeMain,*  
 $pc + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + length (regAux\ item) +$   
 $length (concat (map\ fillAux (zip\ items\ [1..<length\ items + 1]))) , ref)])$   
**prefer 2 apply** (*rule fillAux-instrs*)  
**apply** (*unfold exec-all-def*)  
**apply** (*simp, erule rtrancl-trans*)

**apply** (*drule drop-append-length-3*)

**apply** (*subgoal-tac*)

$P \vdash$   
*(None, sh,*  
 $(let\ vitems = Intg (int (the (com\ C))) \# map (Item2val\ S'\ sh)\ items;$   
 $cnames = (tagGf, cellC) \# map (\lambda\ i. (VName ("arg"@\ nat2Str\ i), cellC))$   
 $[0..<length\ items];$   
 $objs = Obj\ cellC (empty(cnames\ [\mapsto]\ vitems))$   
 $in\ h'(lobj \mapsto objs)),$   
*ih,*  
 $[(\ [],$   
 $vs[Suc\ 0 := Intg\ ii,$   
 $2 := Addr\ lobj,$   
 $3 := Intg (int\ j)],$   
*safeP, sigSafeMain,*

```

    pc + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + length (regAux item) +
    length (concat (map fillAux (zip items [1..<length items + 1]])), ref))]
-jvm→
  (None,sh,
    (let vitems = Intg (int (the (com C))) # map (Item2val S' sh) items;
      cnames = (tagGf, cellC) # map (λ i. (VName ("arg"@ nat2Str i), cellC))
[0..<length items];
      objs = Obj cellC (empty(cnames [↦] vitems))
      in h'(lobj ↦ objs )),
    ih,
    [[Intg (int j)],
     vs[Suc 0 := Intg ii,
        2 := Addr lobj,
        3 := Intg (int j)],
     safeP, sigSafeMain,
     pc + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + length (regAux item) +
     length (concat (map fillAux (zip items [1..<length items + 1]])) + 1, ref]]))
apply (unfold exec-all-def)
apply (simp,erule rtrancl-trans)
prefer 2
apply (rule r-into-rtrancl)
apply (clarify)
apply (unfold JVMExec.exec.simps)
apply (unfold Let-def)
apply (subgoal-tac
  (fst (snd (snd (snd (snd (the (method' (P', safeP) sigSafeMain))))))) !
  (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (the (cdm (l, i))))))))))) +
  length (regAux item) +
  length (concat (map fillAux (zip items [1..<length items + 1]]))) =
  Load 3, simp)
apply (simp,erule nth-via-drop-append)

apply (drule drop-Suc-append)
apply (subgoal-tac
  P ⊢
  (None,sh,
    h'(lobj ↦ Obj cellC [(tagGf, cellC) ↦ Intg (int (the (com C))),
      map (λ i. (VName ("arg"@ nat2Str i), cellC)) [0..<length
items] [↦]
      map (Item2val S' sh) items]),
    ih,
    [[Intg (int j)],
     vs[Suc 0 := Intg ii,
        2 := Addr lobj,
        3 := Intg (int j)],
     safeP, sigSafeMain,
     pc + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + length (regAux item) +
     length (concat (map fillAux (zip items [1..<length items + 1]])) + 1, ref)])

```

```

-jvm→
  (None,sh,
    h'(lobj ↦ Obj cellC [(tagGf, cellC) ↦ Intg (int (the (com C))),
      map (λi. (VName ("arg" @ nat2Str i), cellC)) [0..apply (unfold exec-all-def)
apply (simp,erule rtrancl-trans)
prefer 2
apply (rule r-into-rtrancl)
apply (clarify)
apply (unfold JVMExec.exec.simps)
apply (unfold Let-def)
apply (subgoal-tac
  (fst (snd (snd (snd (snd (the (method' (P', safeP) sigSafeMain))))))) !
  (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (the (cdm (l, i)))))))))))
+ length (regAux item) +
  length (concat (map fillAux (zip items [1..apply (simp,erule nth-via-drop-append)

apply (drule drop-Suc-append)
apply (subgoal-tac
  P' |-
  (None,sha,
    h(lobj ↦
      Obj cellC
      [(tagGf, cellC) ↦ Intg (int (the (com C))),
        map (λi. (VName ("arg" @ nat2Str i), cellC)) [0..

```

```

    the (cdm (l, i)) + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + length (regAux
item) +
    length (concat (map fillAux (zip items [1..<length items + 1]))) + 1 + 1,
ref)])
-jvm→
  (None, sha,
  h(lobj ↦
    Obj cellC
      [(tagGf, cellC) ↦ Intg (int (the (com C))),
      map (λi. (VName ("arg" @ nat2Str i), cellC)) [0..<length
items] [↦]
        map (Item2val S' sha) items], l' ↦ Arr ty m' regS'),

    inih,
    [([],
    vs[Suc 0 := Intg ii,
      2 := Addr lobj,
      3 := Intg (int j)],
    safeP, sigSafeMain,
    the (cdm (l, i)) + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + length (regAux
item) +
    length (concat (map fillAux (zip items [1..<length items + 1]))) + 1 + 1 +
1, ref)])
prefer 2
apply (rule-tac obj=Obj cellC [(tagGf, cellC) ↦ Intg (int (the (com C))),
  map (λi. (VName ("arg" @ nat2Str i), cellC)) [0..<length
items] [↦]
    map (Item2val S' sha) items] and iloc=ii in

insertCell-method-jvm)
apply (simp, assumption+)
apply (simp, erule nth-via-drop-append)
apply (unfold exec-all-def)
apply (elim conjE)
apply (simp, erule rtrancl-trans)

apply (drule drop-Suc-append)
apply (subgoal-tac
  P ⊢
  (None, sha,
  h(lobj ↦
    Obj cellC
      [(tagGf, cellC) ↦ Intg (int (the (com C))),
      map (λi. (VName ("arg" @ nat2Str i), cellC)) [0..<length
items] [↦]
        map (Item2val S' sha) items], l' ↦ Arr ty m' regS'),

    inih,
    [([],
    vs[Suc 0 := Intg ii,

```

```

      2 := Addr lobj,
      3 := Intg (int j)],
  safeP, sigSafeMain,
  pc + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + length (regAux item) +
  length (concat (map fillAux (zip items [1..<length items + 1]))) + 1 + 1 +
1, ref)])
-jvm→
  (None, sha,
  h(lobj ↦
    Obj cellC
      [(tagGf, cellC) ↦ Intg (int (the (com C))),
      map (λi. (VName ("arg" @ nat2Str i), cellC)) [0..<length
items] [↦]
        map (Item2val S' sha) items], l' ↦ Arr ty m' regS'),

  inih,
  [[Intg n],
  vs[Suc 0 := Intg ii,
      2 := Addr lobj,
      3 := Intg (int j)],
  safeP, sigSafeMain,
  pc + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + length (regAux item) +
  length (concat (map fillAux (zip items [1..<length items + 1]))) + 1 + 1 +
1 + 1, ref]]])
apply (unfold exec-all-def)
apply (simp, erule rtrancl-trans)
prefer 2
apply (rule r-into-rtrancl)
apply (clarify)
apply (unfold JVMEexec.exec.simps)
apply (unfold Let-def)
apply (subgoal-tac
  (fst (snd (snd (snd (snd (the (method' (P', safeP) sigSafeMain))))))) !
  (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (the (cdm (l,
i))))))))))))) + length (regAux item) +
  length (concat (map fillAux (zip items [1..<length items + 1]))))))) =
  Getstatic topf stackC, simp)
apply (simp, erule nth-via-drop-append)

apply (drule drop-Suc-append)
apply (subgoal-tac
  P ⊢
  (None, sha,
  h(lobj ↦
    Obj cellC
      [(tagGf, cellC) ↦ Intg (int (the (com C))),
      map (λi. (VName ("arg" @ nat2Str i), cellC)) [0..<length
items] [↦]
        map (Item2val S' sha) items], l' ↦ Arr ty m' regS'),

```

```

    inih,
    [[([Intg n],
      vs[Suc 0 := Intg ii,
         2 := Addr lobj,
         3 := Intg (int j)],
      safeP, sigSafeMain,
      pc + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + length (regAux item) +
      length (concat (map fillAux (zip items [1..<length items + 1])))) + 1 + 1 +
1 + 1, ref]])
  -jvm→
    (None, sha,
     h(lobj ↦
       Obj cellC
        [(tagGf, cellC) ↦ Intg (int (the (com C))),
         map (λi. (VName ("arg" @ nat2Str i), cellC)) [0..<length
items] [↦]
           map (Item2val S' sha) items], l' ↦ Arr ty m' regS'),

    inih,
    [[([Intg 1,
      Intg n],
      vs[Suc 0 := Intg ii,
         2 := Addr lobj,
         3 := Intg (int j)],
      safeP, sigSafeMain,
      pc + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + length (regAux item) +
      length (concat (map fillAux (zip items [1..<length items + 1])))) + 1 + 1 +
1 + 1 + 1, ref]]))
apply (unfold exec-all-def)
apply (simp, erule rtrancl-trans)
prefer 2
apply (rule r-into-rtrancl)
apply (clarify)
apply (unfold JVMEexec.exec.simps)
apply (unfold Let-def)
apply (subgoal-tac
  (fst (snd (snd (snd (snd (the (method' (P', safeP) sigSafeMain))))))) !
  (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (the (cdm
(l, i)))))))))) + length (regAux item) +
  length (concat (map fillAux (zip items [1..<length items + 1])))))))) =
  LitPush (Intg 1), simp)
apply (simp, erule nth-via-drop-append)

apply (drule drop-Suc-append)
apply (subgoal-tac
  P'⊢
  (None, sha,
   h(lobj ↦
     Obj cellC

```

```

      ((tagGf, cellC) ↦ Intg (int (the (com C))),
       map (λi. (VName ("arg" @ nat2Str i), cellC)) [0..apply (unfold exec-all-def)
apply (simp, erule rtrancl-trans)
prefer 2
apply (rule r-into-rtrancl)
apply (clarify)
apply (unfold JVMEexec.exec.simps)
apply (unfold Let-def)
apply (subgoal-tac
  (fst (snd (snd (snd (snd (the (method' (P', safeP) sigSafeMain))))))) !
  (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (the
(cdm (l, i))))))))))) +
length (regAux item) + length (concat (map fillAux (zip items [1..apply (simp, erule nth-via-drop-append)

apply (drule drop-Suc-append)

```



```

apply (subgoal-tac
  P ⊢
  (None, sha,
    h(lobj ⊢
      Obj cellC
        [(tagGf, cellC) ⊢ Intg (int (the (com C))),
          map (λi. (VName ("arg" @ nat2Str i), cellC)) [0..<length
items] [⊢])
          map (Item2val S' sha) items], l' ⊢ Arr ty m' regS'),

    inih,
    [[Intg (n + 1)],
     vs[Suc 0 := Intg ii,
        2 := Addr lobj,
         3 := Intg (int j)],
     safeP, sigSafeMain,
     pc + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + length (regAux item) +
     length (concat (map fillAux (zip items [1..<length items + 1])) + 1 + 1 +
1 + 1 + 1 + 1, ref)])
    -jvm→
    (None, sha((stackC, topf) ⊢ Intg (n + 1)),
     h(lobj ⊢
       Obj cellC
        [(tagGf, cellC) ⊢ Intg (int (the (com C))),
          map (λi. (VName ("arg" @ nat2Str i), cellC)) [0..<length
items] [⊢])
          map (Item2val S' sha) items], l' ⊢ Arr ty m' regS'),

    inih,
    [[]],
    vs[Suc 0 := Intg ii,
        2 := Addr lobj,
         3 := Intg (int j)],
    safeP, sigSafeMain,
    pc + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + length (regAux item) +
    length (concat (map fillAux (zip items [1..<length items + 1])) + 1 + 1 +
1 + 1 + 1 + 1 + 1, ref)))]
apply (unfold exec-all-def)
apply (simp,erule rtrancl-trans)
prefer 2
apply (rule r-into-rtrancl)
apply (clarify)
apply (unfold JVMExec.exec.simps)
apply (unfold Let-def)
apply (subgoal-tac
  (fst (snd (snd (snd (snd (the (method' (P', safeP) sigSafeMain))))))) !
  (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc
(the (cdm (l, i)))))))))))) +
length (regAux item) + length (concat (map fillAux (zip items [1..<length items
+ 1])))))))))) =
  Putstatic topf stackC, simp)

```

**apply** (*simp*, *erule nth-via-drop-append*)

**apply** (*drule drop-Suc-append*)

**apply** (*subgoal-tac*

$P \vdash$

(*None*, *sha*((*stackC*, *topf*)  $\mapsto$  *Intg* ( $n + 1$ )),

*h*(*lobj*  $\mapsto$

*Obj cellC*

[(*tagGf*, *cellC*)  $\mapsto$  *Intg* (*int* (*the* (*com C*))),

*map* ( $\lambda i. (VName ("arg" @ nat2Str i), cellC)) [0..<length$

*items*]  $\mapsto$

*map* (*Item2val S' sha*) *items*],  $l' \mapsto Arr ty m' regS'$ ),

*inih*,

[([],

*vs*[*Suc* 0 := *Intg ii*,

2 := *Addr lobj*,

3 := *Intg* (*int j*)],

*safeP*, *sigSafeMain*,

*pc* + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + *length* (*regAux item*) +

*length* (*concat* (*map fillAux* (*zip items* [ $1..<length items + 1$ ])))) + 1 + 1 +

1 + 1 + 1 + 1 + 1, *ref*]))

-*jvm*  $\rightarrow$

(*None*, *sha*((*stackC*, *topf*)  $\mapsto$  *Intg* ( $n + 1$ )),

*h*(*lobj*  $\mapsto$

*Obj cellC*

[(*tagGf*, *cellC*)  $\mapsto$  *Intg* (*int* (*the* (*com C*))),

*map* ( $\lambda i. (VName ("arg" @ nat2Str i), cellC)) [0..<length$

*items*]  $\mapsto$

*map* (*Item2val S' sha*) *items*],  $l' \mapsto Arr ty m' regS'$ ),

*inih*,

[([*Addr la*],

*vs*[*Suc* 0 := *Intg ii*,

2 := *Addr lobj*,

3 := *Intg* (*int j*)],

*safeP*, *sigSafeMain*,

*pc* + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + *length* (*regAux item*) +

*length* (*concat* (*map fillAux* (*zip items* [ $1..<length items + 1$ ])))) + 1 + 1 +

1 + 1 + 1 + 1 + 1 + 1, *ref*]))

**apply** (*unfold exec-all-def*)

**apply** (*simp*, *erule rtrancl-trans*)

**prefer** 2

**apply** (*rule r-into-rtrancl*)

**apply** (*clarify*)

**apply** (*unfold JVMExec.exec.simps*)

**apply** (*unfold Let-def*)

**apply** (*subgoal-tac*

(*fst* (*snd* (*snd* (*snd* (*snd* (*the* (*method'* ( $P'$ , *safeP*) *sigSafeMain*)))))))) !

(*Suc* (*Suc* (*Suc* (*Suc* (*Suc* (*Suc* (*Suc* (*Suc* (*Suc* (*Suc* (*Suc* (*Suc* (*Suc* (*Suc* (*Suc* (*Suc*

```

(Suc (the (cdm (l, i))))))))) +
  length (regAux item) + length (concat (map fillAux (zip items [1..<length items
+ 1]))))))) =
  Getstatic Sf stackC, simp)
apply (simp add: Sf-def add: topf-def)
apply (simp, erule nth-via-drop-append)

```

```

apply (drule drop-Suc-append)

```

```

apply (subgoal-tac

```

```

  P ⊢

```

```

  (None, sha((stackC, topf) ↦ Intg (n + 1)),

```

```

    h(lobj ↦

```

```

      Obj cellC

```

```

      [(tagGf, cellC) ↦ Intg (int (the (com C))),

```

```

        map (λi. (VName ("arg" @ nat2Str i), cellC)) [0..<length

```

```

items] [↦]

```

```

      map (Item2val S' sha) items], l' ↦ Arr ty m' regS'),

```

```

    inih,

```

```

    [[Addr la],

```

```

      vs[Suc 0 := Intg ii,

```

```

        2 := Addr lobj,

```

```

        3 := Intg (int j)],

```

```

    safeP, sigSafeMain,

```

```

    pc + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + length (regAux item) +

```

```

    length (concat (map fillAux (zip items [1..<length items + 1]))) + 1 + 1 +

```

```

1 + 1 + 1 + 1 + 1 + 1, ref)])

```

```

  -jvm →

```

```

  (None, sha((stackC, topf) ↦ Intg (n + 1)),

```

```

    h(lobj ↦

```

```

      Obj cellC

```

```

      [(tagGf, cellC) ↦ Intg (int (the (com C))),

```

```

        map (λi. (VName ("arg" @ nat2Str i), cellC)) [0..<length

```

```

items] [↦]

```

```

      map (Item2val S' sha) items], l' ↦ Arr ty m' regS'),

```

```

    inih,

```

```

    [[Intg (n + 1),

```

```

      Addr la],

```

```

      vs[Suc 0 := Intg ii,

```

```

        2 := Addr lobj,

```

```

        3 := Intg (int j)],

```

```

    safeP, sigSafeMain,

```

```

    pc + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + length (regAux item) +

```

```

    length (concat (map fillAux (zip items [1..<length items + 1]))) + 1 + 1 +

```

```

1 + 1 + 1 + 1 + 1 + 1 + 1, ref)])

```

```

apply (unfold exec-all-def)

```

```

apply (simp, erule rtrancl-trans)

```

```

prefer 2

```

```

apply (rule r-into-rtrancl)

```

```

apply (clarify)
apply (unfold JVMExec.exec.simps)
apply (unfold Let-def)
apply (subgoal-tac)
  (fst (snd (snd (snd (snd (the (method' (P', safeP) sigSafeMain)))))) !
  (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (the (cdm (l, i)))))))))))))) +
  (length (regAux item) + length (concat (map fillAux (zip items [1..length items
+ 1])))))))))))) =
  (Getstatic topf stackC, simp)
apply (simp, erule nth-via-drop-append)

```

```

apply (drule drop-Suc-append)
apply (subgoal-tac)
   $P \vdash$ 
  (None, sha((stackC, topf)  $\mapsto$  Intg (n + 1)),
  h(lobj  $\mapsto$ 
    (Obj cellC
      [(tagGf, cellC)  $\mapsto$  Intg (int (the (com C))),
      map ( $\lambda i.$  (VName ("arg" @ nat2Str i), cellC)) [0..length
items]  $\mapsto$ ]
        map (Item2val S' sha) items], l'  $\mapsto$  Arr ty m' regS'),
    inih,
    [(Intg (n + 1),
      Addr la],
    vs[Suc 0 := Intg ii,
      2 := Addr lobj,
      3 := Intg (int j)],
    safeP, sigSafeMain,
    pc + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + length (regAux item) +
    length (concat (map fillAux (zip items [1..length items + 1])) + 1 + 1 +
1 + 1 + 1 + 1 + 1 + 1 + 1, ref))]
  -jvm  $\rightarrow$ 
  (None, sha((stackC, topf)  $\mapsto$  Intg (n + 1)),
  h(lobj  $\mapsto$ 
    (Obj cellC
      [(tagGf, cellC)  $\mapsto$  Intg (int (the (com C))),
      map ( $\lambda i.$  (VName ("arg" @ nat2Str i), cellC)) [0..length
items]  $\mapsto$ ]
        map (Item2val S' sha) items], l'  $\mapsto$  Arr ty m' regS'),
    inih,
    [(Intg ii,
      Intg (n + 1),
      Addr la],
    vs[Suc 0 := Intg ii,
      2 := Addr lobj,
      3 := Intg (int j)],
    safeP, sigSafeMain,

```

```

    pc + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + length (regAux item) +
    length (concat (map fillAux (zip items [1..<length items + 1]))) + 1 + 1 +
1 + 1 + 1 + 1 + 1 + 1 + 1 + 1, ref]))
apply (unfold exec-all-def)
apply (simp, erule rtrancl-trans)
prefer 2
apply (rule r-into-rtrancl)
apply (clarify)
apply (unfold JVMExec.exec.simps)
apply (unfold Let-def)
apply (subgoal-tac
  (fst (snd (snd (snd (snd (the (method' (P', safeP) sigSafeMain))))))) !
  (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc
(Suc (Suc (Suc (the (cdm (l, i))))))))))) +
    length (regAux item) + length (concat (map fillAux (zip items [1..<length items
+ 1])))))))))))) =
    Load (Suc 0), simp)
apply (simp, erule nth-via-drop-append)

```

```

apply (drule drop-Suc-append)
apply (subgoal-tac
  P ⊢
  (None, sha((stackC, topf) ↦ Intg (n + 1)),
    h(lobj ↦
      Obj cellC
        [(tagGf, cellC) ↦ Intg (int (the (com C))),
          map (λi. (VName ("arg" @ nat2Str i), cellC)) [0..<length
items] [↦]
            map (Item2val S' sha) items], l' ↦ Arr ty m' regS'),
    inih,
    [(Intg ii,
      Intg (n + 1),
      Addr la],
    vs[Suc 0 := Intg ii,
      2 := Addr lobj,
      3 := Intg (int j)],
    safeP, sigSafeMain,
    pc + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + length (regAux item) +
    length (concat (map fillAux (zip items [1..<length items + 1]))) + 1 + 1 +
1 + 1 + 1 + 1 + 1 + 1 + 1, ref))]
  -jvm →
  (None, sha((stackC, topf) ↦ Intg (n + 1)),
    h(lobj ↦
      Obj cellC
        [(tagGf, cellC) ↦ Intg (int (the (com C))),
          map (λi. (VName ("arg" @ nat2Str i), cellC)) [0..<length
items] [↦]
            map (Item2val S' sha) items],

```

```

      l'  $\mapsto$  Arr ty m' regS', la  $\mapsto$  Arr ty m (S'(nat (n + 1)  $\mapsto$  Intg ii)),
    inih,
    [([]],
      vs[Suc 0 := Intg ii,
          2 := Addr lobj,
          3 := Intg (int j)],
    safeP, sigSafeMain,
    pc + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + length (regAux item) +
    length (concat (map fillAux (zip items [1..<length items + 1]))) + 1 + 1 +
    1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1, ref]]))
  apply (unfold exec-all-def)
  prefer 2
  apply (rule r-into-rtrancl)
  apply (clarify)
  apply (unfold JVMExec.exec.simps)
  apply (unfold Let-def)
  apply (subgoal-tac
    (fst (snd (snd (snd (snd (the (method' (P', safeP) sigSafeMain))))))) !
    (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc
    (Suc (Suc (Suc (Suc (the (cdm (l, i)))))))))) +
    length (regAux item) + length (concat (map fillAux (zip items [1..<length items
    + 1])))))))))))) =
    ArrStore, simp)
  apply (simp add: raise-system-xcpt-def)
  apply (frule arrayIndex-S'-valid, assumption+)
  apply simp
  apply (simp, erule nth-via-drop-append)
  apply simp

  apply (frule nonJumping-Suc-pc)
  apply (erule-tac sym [where t=BUILDCLS C items item])
  apply (simp add: nonJumping.simps)
  apply (simp add: trInstr.simps)

  apply (rule-tac x=hm(getFresh hm  $\mapsto$  (j, C, map (item2Val S) items)) in exI)
  apply (rule-tac x=k in exI)
  apply (rule-tac x=nat (the-Intg (the (sha (heapC, k0f)))) in exI)
  apply (rule-tac x=(l, Suc i) in exI)
  apply (rule-tac x=Val (Val.Loc (getFresh hm)) # S in exI)

  apply (rule-tac x=sha((stackC, topf)  $\mapsto$  Intg (n + 1)) in exI)
  apply (rule-tac x=h(lobj  $\mapsto$  Obj cellC [(tagGf, cellC)  $\mapsto$  Intg (int (the (com C))),
    map ( $\lambda$ i. (VName ("arg" @ nat2Str i), cellC)) [0..<length
    items] [ $\mapsto$ ]
    map (Item2val S' sha) items],
    l'  $\mapsto$  Arr ty m' regS',
    la  $\mapsto$  Arr ty m (S'(nat (n + 1)  $\mapsto$  Intg ii)) in exI)

```

```

apply (rule-tac x=inih in exI)
apply (rule-tac x=vs[Suc 0 := Intg ii, 2 := Addr lobj, 3 := Intg (int j)] in exI)
apply (rule-tac x=Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc
(Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (the (cdm (l, i)) +
length (regAux item) +
length (concat (map fillAux (zip items [Suc 0..in exI)
apply (rule-tac x=ref in exI)

apply (rule-tac x=nat (the-Intg (the ((sha((stackC, topf) ↦ Intg (n + 1)))
(heapC, kf)))) in exI)

apply (rule-tac x=nat (the-Intg (the ((sha((stackC, topf) ↦ Intg (n + 1)))
(heapC, k0f)))) in exI)

apply (rule-tac x=la in exI)
apply (rule-tac x=ty in exI)
apply (rule-tac x=m in exI)
apply (rule-tac x= S'(nat (n + 1) ↦ Intg ii) in exI)

apply (rule-tac x=(n + 1) in exI)
apply (rule-tac x=l' in exI)
apply (rule-tac x=regS' in exI)
apply (rule-tac x=l'' in exI)
apply (rule-tac x=m' in exI)
apply (rule-tac x=m'' in exI)
apply (rule-tac x=d in exI)
apply (rule-tac x=(g(getFresh hm ↦ lobj)) in exI)

apply (fold extractBytecode-def)
apply (frule-tac dh'=h' and lobj=lobj and obj=obj and ii=ii
in execSVMInstr-BUILDCLS-reserveCell)
apply (assumption+, simp add: trInstr.simps, assumption+, simp)
apply (elim conjE)

apply (thin-tac P' ⊢ ?s1 -jvm→ ?s2)

apply (subgoal-tac
  P' |—
  (None, sha, dh(lobj ↦ Obj cellC [(tagGf, cellC) ↦ Intg (int (the (com C))),
  map (λi. (VName ("arg'' @ nat2Str i), cellC)) [0..

```

```

item) +
  length (concat (map fillAux (zip items [1..<length items + 1]))) + 1 +
1,ref]))
-jvm→
  (None, sha, dh(lobj ↦ Obj cellC [(tagGf, cellC) ↦ Intg (int (the (com C))),
  map (λi. (VName ("arg" @ nat2Str i), cellC)) [0..<length items]
[↦]
      map (Item2val S' sh) items], l' ↦ Arr ty m' regS'),
ih,
[[], vs, safeP, sigSafeMain,
the (cdm (l, i)) + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + length (regAux
item) +
  length (concat (map fillAux (zip items [1..<length items + 1]))) + 1 + 1 +
1,ref]))
prefer 2 apply (rule-tac
  obj=Obj cellC [(tagGf, cellC) ↦ Intg (int (the (com C))),
  map (λi. (VName ("arg" @ nat2Str i), cellC)) [0..<length items] [↦]
  map (Item2val S' sh) items] and
  iloc=i and j=j and m''=m'' and d=d and
  pc=the (cdm (l, i)) + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + length
(regAux item) +
  length (concat (map fillAux (zip items [1..<length items + 1]))) + 1 + 1
  in insertCell-method-jvm)
apply (simp, assumption+, simp, simp, simp, assumption+)
apply (drule drop-Suc-append-2)
apply (drule drop-Suc-append)+
apply (drule drop-append-length)
apply (drule drop-append-length-2)
apply (drule drop-Suc-append)
apply (drule drop-Suc-append)
apply (simp add: extractBytecode-def)
apply (rule-tac zs=bytecode' in nth-via-drop-append, force)

apply (frule-tac
  obj=Obj cellC [(tagGf, cellC) ↦ Intg (int (the (com C))),
  map (λi. (VName ("arg" @ nat2Str i), cellC)) [0..<length items] [↦]
  map (Item2val S' sh) items] and
  pc=the (cdm (l, i)) + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + length (regAux
item) +
  length (concat (map fillAux (zip items [1..<length items + 1]))) + 1 + 1
and
  P'=P' in insertCell-method-equivState)
apply assumption+
apply (drule drop-Suc-append-2)
apply (drule drop-Suc-append)+
apply (drule drop-append-length)
apply (drule drop-append-length-2)
apply (drule drop-Suc-append)
apply (drule drop-Suc-append)

```



```

apply (simp add: extractBytecode-def)
apply (rule-tac zs=bytecode' in nth-via-drop-append, force)
apply simp
apply (elim conjE)

apply (thin-tac P' ⊢ ?s1 -jvm→ ?s2)

apply (rule conjI, simp)
apply (rule conjI, simp)
apply (rule conjI, rule refl)
apply (rule conjI, rule refl)
apply (rule conjI) apply clarsimp
apply (rule conjI, assumption)
apply (rule conjI) apply clarsimp
      apply (simp add: heapC-def add: stackC-def)
apply (rule conjI) apply (erule activeCells-5)
apply (rule conjI) apply (erule activeCells-5)
      apply (erule activeCells-5)
apply (rule conjI) apply clarsimp
apply (rule conjI) apply clarsimp
apply (rule conjI) apply clarsimp
apply (rule conjI) apply clarsimp
apply (rule conjI) apply clarsimp
apply (rule conjI) apply clarsimp
apply (rule conjI) apply clarsimp
apply (rule conjI) apply (rule inj-on-BUILDCLS, assumption+)
apply (rule conjI) apply simp
      apply (simp add: heapC-def add: stackC-def)
      apply (fold heapC-def)
      apply (subgoal-tac equivS Sa S' (the-Intg (the (sh (stackC, topf)))))
ctm d g)
      apply (frule S-good) apply simp
      apply (rule-tac l''=l'' in equivH-BUILDCLS)
      apply (assumption+, simp, assumption+, simp, simp, simp, simp, simp)
      apply (simp, simp, simp, assumption+, simp, simp, fold stackC-def, simp)
apply (rule conjI) apply clarsimp
      apply (rule maxPush-BUILDCLS)
apply (rule conjI) apply clarsimp
      apply (rule equivS2-BUILDCLS, assumption+)
apply (rule conjI) apply (simp add: heapC-def add: stackC-def)
by simp

end

```

```

theory dem-REUSE
imports ../JVMSAFE/JVMExec SVM2JVM SVMSemantics CertifSVM2JVM
          dem-translation
begin

no-translations Norm s == (None,s)

no-translations ex-table-of m == snd (snd (snd m))

declare trInstr.simps [simp del]
declare equivS.simps [simp del]

axioms equivS-ge-n-distinct-g:
  equivS S S' n ctm d g  $\implies$  ( $\forall$  p > 0. equivS S (S'(nat (p + n)  $\mapsto$  A)) n ctm d
  g')

lemma S2-REUSE:
  execSVMInst (REUSE) (map-of ct) (hm, k) k0 (l, i) S = Either.Right S2
   $\implies$   $\exists$  b S'. ((hm(b := None)(getFresh hm  $\mapsto$  the (hm b)), k), k0, incrPC (l, i),
  Val (Val.Loc (getFresh hm)) # S') = S2  $\wedge$ 
  S = Val (Val.Loc b) # S'
apply (unfold execSVMInst.simps)
apply (case-tac S, simp-all)
apply (insert RightNotUndefined)
apply (erule-tac x = S2 in allE, force)
apply (case-tac a, simp-all)
apply (case-tac Vala, simp-all)
apply (erule-tac x = S2 in allE, force)
apply simp
apply simp
by simp

lemma equivS-Val-Loc:
  equivS (Val (Val.Loc b) # S) S' n ctm d g
   $\implies$  ( $\exists$  i. the (S' (nat n)) = Intg i  $\wedge$  ( $\exists$  l'. d (nat i) = Some (Addr l')  $\wedge$  g b =
  Some l'))
by (simp add: equivS.simps)

```

**lemma** *equivS2-REUSE*:

$\llbracket \text{equivS } (\text{Val } (\text{Val.Loc } b) \# S) S' n \text{ ctm } d g; d (\text{nat } ii) = \text{Some } (\text{Addr } l'); g b' = \text{Some } l' \rrbracket$   
 $\implies \text{equivS } (\text{Val } (\text{Val.Loc } (\text{getFresh } hm)) \# S) (S'(\text{nat } n \mapsto \text{Intg } ii)) n \text{ ctm } d$   
 $(g(b := \text{None})(\text{getFresh } hm \mapsto l'))$   
**apply** (*frule equivS-Val-Loc*)  
**apply** (*erule exE, elim conjE*)  
**apply** (*erule exE, elim conjE*)  
**apply** (*simp add: equivS.simps*)  
**apply** (*elim conjE*)  
**apply** (*frule equivS-ge-n-distinct-g*)  
**apply** (*erule-tac x=1 in allE*)  
**by** *simp*

**axioms** *getFresh-notin-dom-h*:

*getFresh hm*  $\notin$  *dom hm*

**lemma** *inj-on-REUSE*:

$\llbracket \text{inj-on } g (\text{dom } hm); g b = \text{Some } l; \text{dom } g = \text{dom } hm; \text{getFresh } hm \notin \text{dom } g; \text{getFresh } hm \notin \text{dom } hm \rrbracket$   
 $\implies \text{inj-on } (g(b := \text{None})(\text{getFresh } hm \mapsto l)) (\text{dom } (hm(b := \text{None})(\text{getFresh } hm \mapsto \text{the } (hm b))))$   
**apply** *simp*  
**apply** (*rule conjI*)  
**apply** (*simp add: inj-on-def*)  
**apply** (*simp add: inj-on-def, clarsimp*)  
**apply** (*split split-if-asm, simp*)  
**apply** (*split split-if-asm, simp, simp*)  
**apply** (*erule-tac x=b in ballE*)  
**apply** (*erule-tac x=x in ballE*)  
**apply** *simp*  
**apply** (*simp add: dom-def*)  
**apply** (*subgoal-tac b  $\in$  dom hm, simp*)  
**apply** (*subgoal-tac g b = Some l*)  
**apply** (*subgoal-tac b  $\in$  dom g*)  
**apply** (*simp add: dom-def*)  
**apply** (*rule domI, assumption*)  
**by** *simp*

**lemma** *x-ran-update*:

$\llbracket g b = \text{Some } l'a; x \in \text{ran } (g(b := \text{None})(\text{getFresh } hm \mapsto l'a)) \rrbracket$   
 $\implies x \in \text{ran } g$   
**apply** (*simp add: ran-def, clarsimp*)

**apply** (*split split-if-asm,simp*)  
**apply** (*rule-tac x=b in exI,simp*)  
**apply** (*simp,split split-if-asm,simp*)  
**by** (*rule-tac x=a in exI,simp*)

**lemma** *x-ran-update-2*:  
 $\llbracket g\ b = \text{Some } l'a; x \in \text{ran } g; \text{getFresh } hm \notin \text{dom } g; \text{inj-on } g\ (\text{dom } h) \rrbracket$   
 $\implies x \in \text{ran } (g(b := \text{None})(\text{getFresh } hm \mapsto l'a))$   
**apply** (*case-tac x = l'a*)  
**apply** (*simp add: ran-def,force*)  
**apply** (*simp add: ran-def*)  
**apply** (*elim exE*)  
**apply** (*rule-tac x=a in exI,simp*)  
**apply** (*rule conjI*)  
**apply** (*simp add: inj-on-def*)  
**apply** (*clarsimp, rule impI*)  
**by** *clarsimp*

**lemma** *activeCells-REUSE*:  
 $\llbracket la \notin \text{activeCells } \text{regS } h\ (\text{nat } (\text{the-Intg } (\text{the } (\text{sha } (\text{heapC}, \text{kf})))));$   
 $\text{ran } g = \text{activeCells } \text{regS } h\ (\text{nat } (\text{the-Intg } (\text{the } (\text{sha } (\text{heapC}, \text{kf})))));$   
 $g\ b = \text{Some } l'a;$   
 $\text{getFresh } hm \notin \text{dom } g;$   
 $\text{inj-on } g\ (\text{dom } hm) \rrbracket$   
 $\implies \text{ran } (g(b := \text{None})(\text{getFresh } hm \mapsto l'a)) =$   
 $\text{activeCells } \text{regS } (h(la \mapsto \text{Arr } \text{ty } m\ (S'a(\text{nat } n \mapsto vb'))))\ (\text{nat } (\text{the-Intg } (\text{the}$   
 $(\text{sha } (\text{heapC}, \text{kf}))))))$   
**apply** (*frule l-not-in-cellReg*)  
**apply** (*rule equalityI*)  
**apply** (*rule subsetI*)  
**apply** (*subgoal-tac x ∈ ran g,simp*)  
**apply** (*simp add: activeCells-def*)  
**apply** (*elim exE, elim conjE*)  
**apply** (*rule-tac x=j in exI,simp*)  
**apply** (*simp add: region-def*)  
**apply** (*elim conjE*)  
**apply** (*rule cellReg-step*)  
**apply** (*rule cellReg-basic*)  
**apply** (*erule-tac x=j in allE*)  
**apply** (*rule cellsReg-monotone-1,assumption+*)  
**apply** (*erule l-not-in-regs*)  
**apply** (*rule x-ran-update,assumption+*)  
**apply** (*rule subsetI*)  
**apply** (*subgoal-tac x ∈ activeCells regS h (nat (the-Intg (the (sha (heapC, kf))))))*)  
**apply** (*subgoal-tac x ∈ ran g*)  
**prefer** 2 **apply** *simp*  
**apply** (*rule x-ran-update-2,assumption+*)

```

apply (simp add: activeCells-def)
apply (elim exE, elim conjE)
apply (rule-tac x=j in exI, simp)
apply (simp add: region-def)
apply (elim conjE)
apply (rule cellReg-step)
  apply (rule cellReg-basic)
apply (erule-tac x=j in allE)
apply (rule cellsReg-monotone-2, assumption+)
by (erule l-not-in-regs)

```

```

declare equivH.simps [simp del]

```

```

axioms l-not-getFresh:
  l ≠ getFresh hm

```

```

lemma same-region-REUSE:
  la ∉ activeCells regS h (nat (the-Intg (the (sha (heapC, kf)))))
  ⇒ region regS h j = region regS (h(la ↦ Arr ty m (S'a(nat n ↦ vb^))) j
apply (frule l-not-in-cellReg)
apply (simp add: region-def)
apply (rule equalityI)
  apply (rule subsetI)
apply clarsimp
apply (rule cellReg-step)
  apply (rule cellReg-basic)
apply (erule-tac x=j in allE)
apply (rule cellsReg-monotone-1, assumption+)
apply (erule l-not-in-regs)
apply (rule subsetI)
apply clarsimp
apply (rule cellReg-step)
  apply (rule cellReg-basic)
apply (erule-tac x=j in allE)
apply (rule cellsReg-monotone-2, assumption+)
by (erule l-not-in-regs)

```

```

lemma equivV-REUSE:
  [ i < length vs;
    ∀ i < length vs. ∀ l''. vs!i = Val.Loc l'' ⇒ l'' ≠ b ∧ l'' ∈ dom hm;
    getFresh hm ∉ dom g;
    b ≠ getFresh hm;
    equivV (vs ! i) (argCell flds i) d g ]
  ⇒ equivV (vs ! i) (argCell flds i) d (g(b := None)(getFresh hm ↦ l'a))
apply (case-tac vs!i)

```

**apply** *clarsimp*  
**apply** *simp*  
**by** *simp*

**axioms** *vs-i-good*:

$\forall i < \text{length } vs. \forall l''. vs!i = \text{Val.Loc } l'' \longrightarrow l'' \neq b \wedge l'' \in \text{dom } hm$

**lemma** *domH-REUSE*:

$\llbracket la \notin \text{activeCells } regS \ h \ (nat \ (the\text{-Intg} \ (the \ (sha \ (heapC, \ kf))))); \text{dom } g = \text{dom } hm;$

$\forall l \in \text{dom } hm. \exists l'. l' = the \ (g \ l) \wedge equivC \ (the \ (hm \ l)) \ h \ l' \ (the \ (h \ l')) \ (nat \ (the\text{-Intg} \ (the \ (sha \ (heapC, \ kf)))) \ com \ regS \ d \ g;$

$g \ b = \text{Some } l'a \rrbracket$

$\implies \forall l \in \text{dom} \ (hm(b := None)(getFresh \ hm \mapsto the \ (hm \ b))).$

$\exists l'. l' = the \ ((g(b := None)(getFresh \ hm \mapsto l'a)) \ l) \wedge$

$equivC \ (the \ ((hm(b := None)(getFresh \ hm \mapsto the \ (hm \ b))) \ l)) \ (h(la$

$\mapsto Arr \ ty \ m \ (S'a(nat \ n \mapsto vb')))) \ l'$

$(the \ ((h(la \mapsto Arr \ ty \ m \ (S'a(nat \ n \mapsto vb')))) \ l')) \ (nat \ (the\text{-Intg} \ (the$

$(sha \ (heapC, \ kf)))) \ com \ regS \ d$

$(g(b := None)(getFresh \ hm \mapsto l'a))$

**apply** *(rule ballI)*

**apply** *(erule-tac x=l in ballE)*

**prefer** 2 **apply** *(subgoal-tac l  $\neq$  getFresh hm)*

**prefer** 2 **apply** *(rule l-not-getFresh)*

**apply** *simp*

**apply** *(elim exE)*

**apply** *(rule-tac x=l' in exI)*

**apply** *(rule conjI)*

**apply** *(subgoal-tac l  $\neq$  getFresh hm)*

**prefer** 2 **apply** *(rule l-not-getFresh)*

**apply** *simp* **apply** *(elim conjE)*

**apply** *(unfold equivC.simps)*

**apply** *(elim exE)* **apply** *(elim conjE)*

**apply** *(rule-tac x=j in exI)*

**apply** *(rule-tac x=C in exI)*

**apply** *(rule-tac x=vs in exI)*

**apply** *(rule-tac x=Obj in exI)*

**apply** *(rule-tac x=cname in exI)*

**apply** *(rule-tac x=flds in exI)*

**apply** *(rule-tac x=tag in exI)*

**apply** *(rule-tac x=na in exI)*

**apply** *(rule-tac x=reg-j in exI)*

**apply** *(rule-tac x=tag' in exI)*

**apply** *(rule conjI)*

**apply** *(subgoal-tac l  $\neq$  getFresh hm, simp)*

**apply** *(rule l-not-getFresh)*

**apply** *(rule conjI, assumption)*

```

apply (rule conjI)
  apply (subgoal-tac l' ≠ la, simp)
  apply (simp add: activeCells-def)
  apply (erule-tac x=j in allE, clarsimp)
apply (rule conjI, assumption)
apply (rule conjI, assumption)
apply (rule conjI, clarsimp)
  apply (rule same-region-REUSE, assumption+)
apply (rule conjI, assumption)
apply (rule conjI, assumption)
apply (rule conjI, assumption)
apply (rule conjI, assumption)
apply simp
apply (rule allI, rule impI)
apply (erule-tac x=i in allE, simp)
apply clarsimp
apply (subgoal-tac ∀ i < length vs. ∀ l''. vs!i = Val.Loc l'' → l'' ≠ b ∧ l'' ∈
  dom hm)
prefer 2 apply (rule vs-i-good)
apply (subgoal-tac getFresh hm ∉ dom g)
  prefer 2 apply (subgoal-tac getFresh hm ∉ dom hm, simp)
  apply (rule getFresh-notin-dom-h)
by (rule equivV-REUSE, assumption+, blast, assumption)

```

**lemma** equivH-REUSE:

```

[[ la ∉ activeCells regS h (nat (the-Intg (the (sha (heapC, kf)))));
  dom g = dom hm; g b = Some l'a; getFresh hm ∉ dom g; inj-on g (dom hm);
  equivH (hm, k) h (nat (the-Intg (the (sha (heapC, kf)))) com regS d g ]]
  ⇒ equivH (hm (b := None)(getFresh hm ↦ the (hm b)), k) (h(la ↦ Arr ty m
  (S'a(nat n ↦ vb^)))
  (nat (the-Intg (the (sha (heapC, kf)))) com regS d (g(b := None)(getFresh
  hm ↦ l'a)))
apply (unfold equivH.simps, elim conjE)
apply (rule conjI, clarsimp)
apply (rule conjI)
apply (rule activeCells-REUSE, assumption+)
apply (rule conjI, clarsimp)
by (rule domH-REUSE, assumption+)

```

```

declare extractBytecode-def [simp del]
declare initConsTable-def [simp del]
declare equivS.simps [simp del]

```

**axioms** vs-length:

$length\ vs = 10$

**axioms** *copyCell-method-jvm*:

```

[[ s = (None, shp, hp, ihp, [[(vb),loc,safeP, sigSafeMain,pc,z1,z2]]);
  shp (stackC,Sf) = Some (Addr l);
  shp (StackC,topf) = Some (Intg ntop);
  hp l = Some (Arr ty ma S');
  vb = Intg i;
  d (nat i) = Somde (Addr l');
  Invoke-static heapC copyCell [PrimT Integer] =
    fst(snd(snd(snd(snd(the(method' (P',safeP) sigSafeMain)))))) !
pc
]] ==>
P' |- s -jvm -> (None, shp, hp, ihp, [[(Intg ii),loc,safeP,sigSafeMain,pc+1,z1,z2]])

```

**axioms** *copyCell-method-equivState*:

```

[[ s = (None, shp, hp, ihp,
  ((vb),loc,safeP, sigSafeMain,pc,z1,z2)#[]);
  shp (stackC,Sf) = Some (Addr l);
  shp (StackC,topf) = Some (Intg ntop);
  hp l = Some (Arr ty ma S');
  vb = Intg i;
  d (nat i) = Somde (Addr l');
  Invoke-static heapC copyCell [PrimT Integer] =
    fst(snd(snd(snd(snd(the(method' (P',safeP) sigSafeMain)))))) !
pc
]] ==>
P' |- s -jvm -> (None, shp,
  hp,
  ihp, ((vb'),loc,safeP,sigSafeMain,pc+1,z1,z2)#[]) ^
vb' = Intg ii ^
d (nat i) = None ^
d (nat ii) = Some (Addr l')

```

**declare** *execSVMInst.simps* [*simp del*]

**lemma** *execSVMInstr-REUSE* :

```

[[ (P', cdm, ctm, com) = trSVM2JVM ((svms, ctmap), ini, ct, ah, ai, bc);
  cdm, ctm, com ⊢ ((hm, k), k0, (l, i), S) ≐ S1';
  (fst (the (map-of svms l)) ! i) = REUSE;
  execSVMInst REUSE (map-of ct) (hm, k) k0 (l, i) S = Either.Right S2;
  drop (the (cdm (l, i))) (extractBytecode P') =
  trInstr (the (cdm (l, i))) cdm' ctm' com pcc REUSE @ bytecode'
]] ==> ∃ v' sh' dh' ih' fms'.
P' ⊢ S1' -jvm -> (v',sh',dh',ih', fms') ^

```



$cdm, ctm, com \vdash S2 \triangleq (v', sh', dh', ih', fms')$

**apply** (*case-tac*  $S1'$ )  
**apply** (*rename-tac*  $v$   $tup$ )  
**apply** (*case-tac*  $tup$ )  
**apply** (*rename-tac*  $sh$   $tup$ )  
**apply** (*case-tac*  $tup$ )  
**apply** (*rename-tac*  $dh$   $tup$ )  
**apply** (*case-tac*  $tup$ )  
**apply** (*rename-tac*  $ih$   $fms$ )  
**apply** (*simp*)

**apply** (*subgoal-tac*  
 $\exists b S'. ((hm(b := None)(getFresh hm \mapsto the (hm b)), k), k0, incrPC (l, i), Val$   
 $(Val.Loc (getFresh hm)) \# S') = S2 \wedge$   
 $S = Val (Val.Loc b) \# S')$   
**prefer** 2 **apply** (*rule*  $S2-REUSE, assumption$ )  
**apply** (*elim*  $exE$ )

**apply** (*unfold*  $equivState-def$ )  
**apply** (*elim*  $exE, elim conjE$ )

**apply** (*subgoal-tac*  $equivS (Val (Val.Loc b) \# S') S'a n ctm d g$ )  
**prefer** 2 **apply**  $clarsimp$   
**apply** (*frule*  $equivS-Val-Loc$ )  
**apply** (*elim*  $exE, elim conjE$ )  
**apply** (*elim*  $exE, elim conjE$ )

**apply** (*subgoal-tac*  $length vs = 10$ )  
**prefer** 2 **apply** (*rule*  $vs-length$ )

**apply** (*rule-tac*  $x=None$  **in**  $exI$ )  
**apply** (*rule-tac*  $x=sha$  **in**  $exI$ )  
**apply** (*rule-tac*  $x=h(la \mapsto Arr ty m (S'a(nat n \mapsto Intg ia)))$  **in**  $exI$ )  
**apply** (*rule-tac*  $x=inih$  **in**  $exI$ )  
**apply** (*rule-tac*  $x=[([],$   
 $vs[Suc 0 := Intg ia],$   
 $safeP, sigSafeMain,$   
 $the (cdm (l, i)) + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1,$   
 $ref)]$  **in**  $exI$ )

**apply** (*subgoal-tac*  $\exists vs pc . fms=(([], vs, safeP, sigSafeMain, pc, ref))\#[]$ )  
**prefer** 2 **apply**  $simp$   
**apply** (*erule*  $exE$ )  
**apply** (*unfold*  $exec-all-def$ )

**apply** (*rule conjI*)

**apply** (*subgoal-tac PC = (l,i)*) **prefer** 2 **apply** *simp*  
**apply** *simp* **apply** (*elim conjE*) **apply** *clarsimp*  
**apply** (*subgoal-tac*  
   $P \vdash$   
    (*None,sha,h,inih*,[([], *vs, safeP, sigSafeMain, the (cdm (l,i)), ag,be*)])  
   $-jvm \rightarrow$   
    (*None,sha,h,inih*,[([*the (sha (stackC,Sf))*], *vs, safeP, sigSafeMain, the (cdm*  
    (*l,i*) + 1, *ag,be*)]))  
**apply** (*unfold exec-all-def*)  
**apply** (*erule rtrancl-trans*)

**prefer** 2

**apply** (*rule r-into-rtrancl*)  
**apply** (*clarify*)  
**apply** (*unfold JVMExec.exec.simps*)  
**apply** (*unfold Let-def*)  
**apply** (*unfold extractBytecode-def*)  
**apply** (*subgoal-tac*  
  (*fst (snd (snd (snd (snd (the (method' (P', safeP) sigSafeMain))))))*) !  
  *the (cdm (l, i)) = Getstatic Sf stackC,simp*)  
**apply** (*unfold trInstr.simps, erule nth-via-drop-append*)

**apply** (*drule drop-Suc-append*)

**apply** (*subgoal-tac*  
   $P \vdash$   
    (*None, sha, h, inih*,  
      [([*the (sha (stackC, Sf))*],  
      *vs, safeP, sigSafeMain, the (cdm (l, i)) + 1, ag, be*)]))  
   $-jvm \rightarrow$   
    (*None,sha,h,inih*,  
      [([*the (sha (stackC,topf))*],  
      *the (sha (stackC,Sf))*],  
      *vs, safeP, sigSafeMain, the (cdm (l, i)) + 1 + 1, ag, be*)]))

**apply** (*unfold exec-all-def*)

**apply** (*erule rtrancl-trans*)

**prefer** 2

**apply** (*rule r-into-rtrancl*)  
**apply** (*clarify*)  
**apply** (*unfold JVMExec.exec.simps*)  
**apply** (*unfold Let-def*)  
**apply** (*subgoal-tac*  
  (*fst (snd (snd (snd (snd (the (method' (P', safeP) sigSafeMain))))))*) ! *Suc (the*  
  (*cdm (l, i))) =*  
  *Getstatic topf stackC,simp*)

**apply** (*erule nth-via-drop-append*)

**apply** (*drule drop-Suc-append*)

**apply** (*subgoal-tac*

$P \vdash$

(*None, sha, h, inih,*  
[[*the (sha (stackC, topf)),*  
*the (sha (stackC, Sf))*]],  
*vs, safeP, sigSafeMain, the (cdm (l, i)) + 1 + 1, ag, be*]])

$-jvm \rightarrow$

(*None, sha, h, inih,*  
[[*the (S'a (nat n))*]],  
*vs, safeP, sigSafeMain, the (cdm (l, i)) + 1 + 1 + 1, ag, be*]])

**apply** (*unfold exec-all-def*)

**apply** (*erule rtrancl-trans*)

**prefer** 2

**apply** (*rule r-into-rtrancl*)

**apply** (*clarify*)

**apply** (*unfold JVMEexec.exec.simps*)

**apply** (*unfold Let-def*)

**apply** (*subgoal-tac*

(*fst (snd (snd (snd (snd (the (method' (P', safeP) sigSafeMain)))))) ! (Suc (Suc*  
*(the (cdm (l, i)))))) =*

*ArrLoad, simp*)

**apply** (*simp add: raise-system-xcpt-def*)

**apply** (*subgoal-tac n >= 0, simp*)

**apply** (*simp add: equivS.simps*)

**apply** (*erule nth-via-drop-append*)

**apply** (*drule drop-Suc-append*)

**apply** (*frule equivS-Val-Loc*)

**apply** (*elim exE, elim conjE*)

**apply** (*elim exE, elim conjE*)

**apply** (*subgoal-tac*

$P \vdash$

(*None, sha, h, inih,*  
[[*the (S'a (nat n))*]],  
*vs, safeP, sigSafeMain, the (cdm (l, i)) + 1 + 1 + 1, ag, be*]])

$-jvm \rightarrow$

(*None, sha, h, inih,*  
[[*Intg ia*]],  
*vs, safeP, sigSafeMain, the (cdm (l, i)) + 1 + 1 + 1 + 1, ag, be*]])

**prefer** 2

**apply** (*rule-tac vb=the (S'a (nat n)) in copyCell-method-jvm*)

**apply** (*simp, assumption+*)

**apply** (*rule sym*)

**apply** (*simp, erule nth-via-drop*)

**apply** (*unfold exec-all-def*)  
**apply** (*erule rtrancl-trans*)

**apply** (*drule drop-Suc-append*)  
**apply** (*subgoal-tac*)  
 $P \vdash$   
 (*None, sha, h, inih,*  
 [([*Intg ia*],  
*vs, safeP, sigSafeMain, the (cdm (l, i)) + 1 + 1 + 1 + 1, ag, be*)]])  
 $-jvm \rightarrow$   
 (*None, sha, h, inih,*  
 [([],  
*vs[Suc 0 := Intg ia],*  
*safeP, sigSafeMain, the (cdm (l, i)) + 1 + 1 + 1 + 1 + 1, ag, be*)]])  
**apply** (*unfold exec-all-def*)  
**apply** (*erule rtrancl-trans*)  
**prefer** 2  
**apply** (*rule r-into-rtrancl*)  
**apply** (*clarify*)  
**apply** (*unfold JVMExec.exec.simps*)  
**apply** (*unfold Let-def*)  
**apply** (*subgoal-tac*)  
 (*fst (snd (snd (snd (snd (the (method' (P', safeP) sigSafeMain)))))) ! (Suc (Suc*  
 (*Suc (Suc (the (cdm (l, i)))))) =*  
*Store 1, simp*)  
**apply** (*erule nth-via-drop-append*)

**apply** (*drule drop-Suc-append*)  
**apply** (*subgoal-tac*)  
 $P \vdash$   
 (*None, sha, h, inih,*  
 [([],  
*vs[Suc 0 := Intg ia],*  
*safeP, sigSafeMain, the (cdm (l, i)) + 1 + 1 + 1 + 1 + 1, ag, be*)]])  
 $-jvm \rightarrow$   
 (*None, sha, h, inih,*  
 [([*the (sha (stackC, Sf))*],  
*vs[Suc 0 := Intg ia],*  
*safeP, sigSafeMain, the (cdm (l, i)) + 1 + 1 + 1 + 1 + 1 + 1, ag, be*)]])  
**apply** (*unfold exec-all-def*)  
**apply** (*erule rtrancl-trans*)  
**prefer** 2  
**apply** (*rule r-into-rtrancl*)  
**apply** (*clarify*)  
**apply** (*unfold JVMExec.exec.simps*)

```

apply (unfold Let-def)
apply (subgoal-tac
  (fst (snd (snd (snd (snd (the (method' (P', safeP) sigSafeMain))))))) !
  (Suc (Suc (Suc (Suc (Suc (the (cdm (l, i)))))))) =
  Getstatic Sf stackC,simp)
apply (erule nth-via-drop-append)

```

```

apply (drule drop-Suc-append)
apply (subgoal-tac
  P⊢
  (None,sha,h,inih,
  [[the (sha (stackC, Sf))],
  vs[Suc 0 := Intg ia],
  safeP, sigSafeMain, the (cdm (l, i)) + 1 + 1 + 1 + 1 + 1 + 1, ag, be]])
  -jvm→
  (None,sha,h,inih,
  [[the (sha (stackC,topf)),
  the (sha (stackC, Sf))],
  vs[Suc 0 := Intg ia],
  safeP, sigSafeMain, the (cdm (l, i)) + 1 + 1 + 1 + 1 + 1 + 1 + 1, ag,
  be]))
apply (unfold exec-all-def)
apply (erule rtrancl-trans)
prefer 2
apply (rule r-into-rtrancl)
apply (clarify)
apply (unfold JVMExec.exec.simps)
apply (unfold Let-def)
apply (subgoal-tac
  (fst (snd (snd (snd (snd (the (method' (P', safeP) sigSafeMain))))))) !
  (Suc (Suc (Suc (Suc (Suc (Suc (the (cdm (l, i)))))))))) =
  Getstatic topf stackC,simp)
apply (erule nth-via-drop-append)

```

```

apply (drule drop-Suc-append)
apply (subgoal-tac
  P⊢
  (None,sha,h,inih,
  [[the (sha (stackC,topf)),
  the (sha (stackC, Sf))],
  vs[Suc 0 := Intg ia],
  safeP, sigSafeMain, the (cdm (l, i)) + 1 + 1 + 1 + 1 + 1 + 1 + 1, ag, be]])
  -jvm→
  (None,sha,h,inih,
  [[Intg ia,
  the (sha (stackC,topf)),
  the (sha (stackC, Sf))],

```

```

    vs[Suc 0 := Intg ia],
    safeP, sigSafeMain, the (cdm (l, i)) + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1,
    ag, be]))
apply (unfold exec-all-def)
apply (erule rtrancl-trans)
prefer 2
apply (rule r-into-rtrancl)
apply (clarify)
apply (unfold JVMExec.exec.simps)
apply (unfold Let-def)
apply (subgoal-tac
  (fst (snd (snd (snd (snd (the (method' (P', safeP) sigSafeMain)))))) !
    (Suc (Suc (Suc (Suc (Suc (Suc (Suc (the (cdm (l, i)))))))))) =
      Load 1 ,simp)
apply (erule nth-via-drop-append)

apply (drule drop-Suc-append)
apply (subgoal-tac
  P ⊢
  (None, sha, h, inih,
  [[Intg ia,
    the (sha (stackC, topf)),
    the (sha (stackC, Sf))]],
  vs[Suc 0 := Intg ia],
  safeP, sigSafeMain, the (cdm (l, i)) + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1,
  ag, be]))
  -jvm→
  (None, sha,
  h(la ↦ Arr ty m (S'a(nat n ↦ Intg ia))),
  inih,
  [[[]],
  vs[Suc 0 := Intg ia],
  safeP, sigSafeMain, the (cdm (l, i)) + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 +
  1, ag, be]))
apply (unfold exec-all-def)
apply (erule rtrancl-trans)
prefer 2
apply (rule r-into-rtrancl)
apply (clarify)
apply (unfold JVMExec.exec.simps)
apply (unfold Let-def)
apply (subgoal-tac
  (fst (snd (snd (snd (snd (the (method' (P', safeP) sigSafeMain)))))) !
    (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (the (cdm (l, i)))))))))) =
      ArrStore, simp)
apply (simp add: raise-system-xcpt-def)

```

```

apply (subgoal-tac n ≥ 0, simp)
apply (simp add: equivS.simps)
apply (erule nth-via-drop-append)
apply simp

```

```

apply (frule nonJumping-Suc-pc)
apply (erule-tac sym [where t=REUSE])
apply (simp add: nonJumping.simps)
apply (simp add: trInstr.simps)

```

```

apply (rule-tac x=hm(b := None)(getFresh hm ↦ the (hm b)) in exI)
apply (rule-tac x=k in exI)
apply (rule-tac x=k0 in exI)
apply (rule-tac x=(l, Suc i) in exI)
apply (rule-tac x=Val (Val.Loc (getFresh hm)) # S' in exI)

```

```

apply (rule-tac x=sha in exI)
apply (rule-tac x=h(la ↦ Arr ty m (S'a(nat n ↦ Intg ia))) in exI)
apply (rule-tac x=inih in exI)
apply (rule-tac x=vs[Suc 0 := Intg ia] in exI)
apply (rule-tac x=the (cdm (l, i)) + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 in
exI)
apply (rule-tac x=ref in exI)

```

```

apply (rule-tac x=nat (the-Intg (the (sha (heapC, kf)))) in exI)
apply (rule-tac x=nat (the-Intg (the (sha (heapC, k0f)))) in exI)

```

```

apply (rule-tac x=la in exI)
apply (rule-tac x=ty in exI)
apply (rule-tac x=m in exI)
apply (rule-tac x=(S'a(nat n ↦ Intg ia)) in exI)

```

```

apply (rule-tac x=n in exI)
apply (rule-tac x=l' in exI)
apply (rule-tac x=regS in exI)
apply (rule-tac x=l'' in exI)
apply (rule-tac x=m' in exI)
apply (rule-tac x=m'' in exI)
apply (rule-tac x=d in exI)
apply (rule-tac x=g(b := None)(getFresh hm ↦ l'a) in exI)

```

```

apply (rule conjI, clarsimp)
apply (rule conjI, rule refl)
apply (rule conjI, rule refl)
apply (rule conjI, rule refl)
apply (rule conjI) apply clarsimp

```

```

apply (rule conjI, assumption)
apply (rule conjI) apply clarsimp
      apply (simp add: stackC-def add: heapC-def)
      apply (simp add: kf-def add: k0f-def)
apply (rule conjI) apply (erule activeCells)
apply (rule conjI) apply (erule activeCells-2,assumption)
      apply (erule activeCells-2,assumption)
apply (rule conjI, simp)
apply (rule conjI) apply (simp add: stackC-def add: heapC-def)
apply (rule conjI) apply clarsimp
apply (rule conjI) apply clarsimp
apply (rule conjI) apply clarsimp
apply (rule conjI) apply clarsimp
apply (rule conjI) apply clarsimp
apply (rule conjI) apply (subgoal-tac getFresh hm  $\notin$  dom hm)
      prefer 2 apply (rule getFresh-notin-dom-h)
      apply (subgoal-tac getFresh hm  $\notin$  dom g)
      prefer 2 apply simp
      apply (rule inj-on-REUSE,simp,assumption,simp,simp,assumption+)
apply (rule conjI) apply clarsimp
      apply (subgoal-tac getFresh hm  $\notin$  dom hm)
      prefer 2 apply (rule getFresh-notin-dom-h)
      apply (subgoal-tac getFresh hm  $\notin$  dom g)
      prefer 2 apply simp
      apply (rule equivH-REUSE,assumption+)
apply (rule conjI) apply clarsimp
apply (rule conjI) apply clarsimp
      apply (frule-tac hm=hm in equivS2-REUSE,assumption+)
apply (rule conjI) apply clarsimp
by simp

end

```

```

theory dem-MATCH
imports ../JVMSAFE/JVMEexec SVM2JVM SVMSemantics CertifSVM2JVM
begin

```

```

no-translations Norm s == (None,s)
no-translations ex-table-of m == snd (snd (snd m))

```

```

declare trInstr.simps [simp del]
declare equivS.simps [simp del]
declare extractBytecode-def [simp del]
declare initConsTable-def [simp del]

```

```

consts

```



*upt-rev* :: nat => nat => nat list ((1[->/..-]))

**primrec**

*upt-0-rev*: [0>..i] = []

*upt-Suc-rev*: [(Suc j)>..i] = (if i <= j then j # [j>..i] else [])

**lemma** *length-upt-rev*: length [i>..j] = i - j

**by** (induct i) (auto simp add: Suc-diff-le)

**lemma** *nth-upt-2*: k < j ∧ k >= i ==> [i..<j] ! (k - i) = k

**apply** (induct j)

**apply** (auto simp add: less-Suc-eq nth-append split: nat-diff-split)

**done**

**constdefs** *diff* :: jvm-state => nat

*diff* s ≡ (case s of (none, sh, h, inih, [(], vs, C, M, pc, ref)))  
 => case vs!5 of Intg i =>  
 case vs!4 of Intg nargs =>  
 nat (nargs - i)

**constdefs** *oneStep* :: jvm-prog => jvm-state => jvm-state

*oneStep* P s ≡ (THE s'. P ⊢ s -jvm→ s' ∧  
 (case s of (none, sh, h, inih, [(], vs, C, M, pc, ref)))  
 => (case s' of (none', sh', h', inih', [(], vs', C', M', pc', ref'))  
 => case vs!5 of Intg i  
 => case sh (stackC, Sf) of Some (Addr la)  
 => case h la of Some (Arr ty m S')  
 => case vs!6 of Addr l  
 => case vs!4 of Intg nargs  
 => case h l of Some (Obj cell fs) =>  
 vs' = vs[5 := Intg (i + 1)] ∧  
 none' = none ∧  
 sh' = sh ∧  
 h' = h(the-Addr (the (sh (stackC, Sf))) ↦ Arr ty m  
 (S'(nat (the-Intg (the (sh (stackC, topf))) - i)  
 ↦ argCell fs (nat i + 1)))) ∧  
 inih' = inih ∧  
 C' = C ∧  
 M' = M ∧  
 pc' = pc ∧  
 ref' = ref)))

**consts**

*steps* :: nat => jvm-prog => jvm-state => jvm-state

**primrec**

*steps* 0 P s = s

$steps (Suc\ n)\ P\ s = steps\ n\ P\ (oneStep\ P\ s)$

**constdefs**  $lastState :: jvm-prog \Rightarrow jvm-state \Rightarrow jvm-state$   
 $lastState\ P\ s \equiv steps\ (diff\ s)\ P\ s$

**axioms** *undefined-false*:  
 $undefined \Longrightarrow False$

**axioms** *i-0*:  
 $i = 0$

**axioms** *top-good*:  
 $(\neg\ int\ m \leq the-Intg\ (the\ (sh\ (stackC,\ topf))) - i \wedge \neg\ the-Intg\ (the\ (sh\ (stackC,\ topf))) < i)$

**lemma** *drop-append-3*:  
 $drop\ n\ xs = [] @ (ys @ zs) \Longrightarrow drop\ n\ xs = ys @ zs$   
**by** *simp*

**lemma** *drop-length-append*:  
 $drop\ n\ xs = (ys @ zs) \Longrightarrow drop\ (n + length\ ys)\ xs = zs$   
**by** (*induct\ n\ arbitrary: xs*) (*auto, case-tac xs, auto*)

**lemma** *drop-endlabel1*:  
 $[ [ drop\ (Suc\ (Suc\ (Suc\ (Suc\ (Suc\ (Suc\ (Suc\ (Suc\ (Suc\ pe)))))))))) )$   
 $(fst\ (snd\ (snd\ (snd\ (the\ (method'\ (P',\ safeP)\ sigSafeMain)))))) =$   
 $[Goto\ endlab1] @ InstLabel2 @ InstLabel3 @ InstLabel4 @ InstLabel5 @$   
 $InstLabel6 @$   
 $InstLabel7 @ InstLabel8 @ Match2 @ bytecode ]$   
 $\Longrightarrow drop\ (nat\ (12 + int\ pc + endlab1))\ (fst\ (snd\ (snd\ (snd\ (the\ (method'\$   
 $(P',\ safeP)\ sigSafeMain)))))) =$   
 $Match2 @ bytecode$

**apply** (*drule drop-Suc-append*)  
**apply** (*drule drop-append-3*)  
**apply** (*drule drop-length-append*)  
**apply** (*drule drop-length-append*)  
**apply** (*drule drop-length-append*)  
**apply** (*drule drop-length-append*)  
**apply** (*drule drop-length-append*)  
**apply** (*drule drop-length-append*)  
**apply** (*drule drop-length-append*)  
**apply** (*drule drop-length-append*)  
**apply** (*drule drop-length-append*)  
**apply** (*simp add: endlab1-def*)  
**apply** (*simp add: nlab2-def*)  
**apply** (*simp add: nlab3-def*)  
**apply** (*simp add: nlab4-def*)  
**apply** (*simp add: nlab5-def*)

```

apply (simp add: nlab6-def)
apply (simp add: nlab7-def)
apply (simp add: nlab8-def)
apply (simp add: InstLabel2-def)
apply (simp add: InstLabel3-def)
apply (simp add: InstLabel4-def)
apply (simp add: InstLabel5-def)
apply (simp add: InstLabel6-def)
apply (simp add: InstLabel7-def)
apply (simp add: InstLabel8-def)
apply (subgoal-tac
  (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc
    (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc
    (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc
    (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc
    (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc
    pc)))))))))))))))))))))))))))))))))))))))))))))))))))))) = nat (69 + (int
pc)))
by auto

```

**lemma** *oneStep-MATCH*:

```

[[ (P', cdm, ctm, com) = trSVM2JVM ((svms, ctmap), ini, ct, ah, ai, bc);
  length vs = 10;
  vs ! 6 = Addr l;
  l ≠ la;
  sh (stackC, Sf) = Some (Addr la);
  h la = Some (Arr ty m S');
  h l = Some (Obj cellC fs);
  drop pc (fst (snd (snd (snd (snd (the (method' (P', safeP) sigSafeMain)))))))]
=
  Match12 @ InstLabel1 @ InstLabel2 @ InstLabel3 @ InstLabel4 @ InstLabel5
@
  InstLabel6 @ InstLabel7 @ InstLabel8 @ Match2 @ bytecode;
  s = (None, sh, h, inih, [[[], vs[4:= Intg nargs, 5:= Intg i], safeP, sigSafeMain,
pc, ref]]);
  i < nargs ]]
  ⇒ P' ⊢ (None, sh, h, inih, [[[], vs[4:= Intg nargs, 5:= Intg i], safeP, sigSafe-
Main, pc, ref]])
  -jvm→ (None, sh, h(the-Addr (the (sh (stackC, Sf))) ↦ Arr ty m
    (S'(nat (the-Intg (the (sh (stackC, topf))) - i)
    ↦ argCell fs (nat i + 1))))),
    inih, [[[], vs[4:= Intg nargs, 5:= Intg (i + 1)], safeP,
sigSafeMain, pc, ref]])
apply (unfold Match12-def)

```

```

apply (subgoal-tac

```

$P \vdash (None, sh, h, inih, [([], vs[4 := Intg nargs, 5 := Intg i], safeP, sigSafeMain, pc, ref)])$   
 $-jvm \rightarrow$   
 $(None, sh, h, inih, [([Intg nargs], vs[4 := Intg nargs, 5 := Intg i], safeP, sigSafeMain, pc + 1, ref)])$   
**apply** (unfold exec-all-def)  
**apply** (erule rtrancl-trans)  
**prefer** 2  
**apply** (rule r-into-rtrancl)  
**apply** clarify  
**apply** (unfold JVMEexec.exec.simps)  
**apply** (unfold Let-def)  
**apply** (subgoal-tac (fst (snd (snd (snd (snd (the (method' (P', safeP) sigSafeMain))))))) !  
 $pc = Load\ 4$ ) **apply** simp  
**apply** (rule nth-via-drop-append) **apply** force

**apply** (drule drop-Suc-append)  
**apply** (subgoal-tac  
 $P \vdash (None, sh, h, inih, [([Intg nargs], vs[4 := Intg nargs, 5 := Intg i], safeP, sigSafeMain, pc + 1, ref)])$   
 $-jvm \rightarrow$   
 $(None, sh, h, inih, [([Intg i, Intg nargs], vs[4 := Intg nargs, 5 := Intg i], safeP, sigSafeMain, pc + 1 + 1, ref)])$   
**apply** (unfold exec-all-def)  
**apply** (erule rtrancl-trans)  
**prefer** 2  
**apply** (rule r-into-rtrancl)  
**apply** clarify  
**apply** (unfold JVMEexec.exec.simps)  
**apply** (unfold Let-def)  
**apply** (subgoal-tac (fst (snd (snd (snd (snd (the (method' (P', safeP) sigSafeMain))))))) !  
 $Suc\ pc = Load\ 5$ ) **apply** simp  
**apply** (rule nth-via-drop-append) **apply** force

**apply** (drule drop-Suc-append)  
**apply** (subgoal-tac  
 $P \vdash (None, sh, h, inih, [([Intg i, Intg nargs], vs[4 := Intg nargs, 5 := Intg i], safeP, sigSafeMain, pc + 1 + 1, ref)])$   
 $-jvm \rightarrow$   
 $(None, sh, h, inih, [([], vs[4 := Intg nargs, 5 := Intg i], safeP, sigSafeMain, pc + 1 + 1 + 1, ref)])$   
**apply** (unfold exec-all-def)  
**apply** (erule rtrancl-trans)  
**prefer** 2  
**apply** (rule r-into-rtrancl)

**apply** *clarify*  
**apply** (*unfold JVMExec.exec.simps*)  
**apply** (*unfold Let-def*)  
**apply** (*subgoal-tac (fst (snd (snd (snd (snd (the (method' (P', safeP) sigSafeMain)))))) !*  
*Suc (Suc pc)) = Ifcmp GreaterEqual labelEndLoop)* **apply** *simp*  
**apply** (*rule nth-via-drop-append*) **apply** *force*

**apply** (*drule drop-Suc-append*)  
**apply** (*subgoal-tac*  
*P ⊢ (None, sh, h, inih, [([], vs[4 := Intg nargs, 5 := Intg i], safeP, sigSafeMain, pc*  
*+ 1 + 1 + 1, ref)])*  
*-jvm→*  
*(None, sh, h, inih, [( [Intg i], vs[4 := Intg nargs, 5 := Intg i], safeP, sigSafeMain,*  
*pc + 1 + 1 + 1 + 1, ref)])*)  
**apply** (*unfold exec-all-def*)  
**apply** (*erule rtrancl-trans*)  
**prefer** 2  
**apply** (*rule r-into-rtrancl*)  
**apply** *clarify*  
**apply** (*unfold JVMExec.exec.simps*)  
**apply** (*unfold Let-def*)  
**apply** (*subgoal-tac (fst (snd (snd (snd (snd (the (method' (P', safeP) sigSafeMain)))))) !*  
*Suc (Suc (Suc pc))) = Load 5)* **apply** *simp*  
**apply** (*rule nth-via-drop-append*) **apply** *force*

**apply** (*drule drop-Suc-append*)  
**apply** (*subgoal-tac*  
*P ⊢ (None, sh, h, inih, [( [Intg i], vs[4 := Intg nargs, 5 := Intg i], safeP, sigSafeMain,*  
*pc + 1 + 1 + 1 + 1, ref)])*  
*-jvm→*  
*(None, sh, h, inih, [( [], vs[4 := Intg nargs, 5 := Intg i], safeP, sigSafeMain,*  
*pc + 1 + 1 + 1 + 1 + 1, ref)])*)  
**apply** (*unfold exec-all-def*)  
**apply** (*erule rtrancl-trans*)  
**prefer** 2  
**apply** (*rule r-into-rtrancl*)  
**apply** *clarify*  
**apply** (*unfold JVMExec.exec.simps*)  
**apply** (*unfold Let-def*)  
**apply** (*subgoal-tac (fst (snd (snd (snd (snd (the (method' (P', safeP) sigSafeMain)))))) !*  
*Suc (Suc (Suc (Suc pc))) =*  
*Tableswitch 0 7 [label1, label2, label3, label4, label5, label6, label7,*  
*label8])* **apply** *simp*

**apply** (*subgoal-tac*  $i = 0$ )  
**prefer** 2 **apply** (*rule*  $i-0$ ) **apply** *simp* **apply** (*simp* *add: label1-def*)  
**apply** (*rule* *nth-via-drop-append*) **apply** *force*

**apply** (*drule* *drop-Suc-append*)  
**apply** (*unfold* *InstLabel1-def*)  
**apply** (*drule* *drop-append-3*)  
**apply** (*subgoal-tac*  
 $P \vdash (None, sh, h, inih, [([]), vs[4 := Intg nargs, 5 := Intg i], safeP, sigSafeMain, pc$   
 $+ 1 + 1 + 1 + 1 + 1, ref])$ )  
 $-jvm \rightarrow$   
 $(None, sh, h, inih, [(the (sh (stackC, Sf))],$   
 $vs[4 := Intg nargs, 5 := Intg i], safeP, sigSafeMain, Suc (Suc (Suc (Suc (Suc$   
 $(Suc pc))))), ref]))$ )  
**apply** (*unfold* *exec-all-def*)  
**apply** (*erule* *rtrancl-trans*)  
**prefer** 2  
**apply** (*rule* *r-into-rtrancl*)  
**apply** *clarify*  
**apply** (*unfold* *JVMExec.exec.simps*)  
**apply** (*unfold* *Let-def*)  
**apply** (*subgoal-tac* (*fst* (*snd* (*snd* (*snd* (*snd* (*the* (*method'* ( $P'$ , *safeP*) *sigSafe-*  
*Main*)))))) !  
 $Suc (Suc (Suc (Suc (Suc pc)))) =$   
 $Getstatic Sf stackC$ ) **apply** *simp*  
**apply** (*rule* *nth-via-drop-append*) **apply** *force*

**apply** (*drule* *drop-Suc-append*)  
**apply** (*subgoal-tac*  
 $P \vdash (None, sh, h, inih,$   
 $[(the (sh (stackC, Sf))],$   
 $vs[4 := Intg nargs, 5 := Intg i], safeP, sigSafeMain,$   
 $Suc (Suc (Suc (Suc (Suc pc))))), ref])$ )  
 $-jvm \rightarrow$   
 $(None, sh, h, inih,$   
 $[(the (sh (stackC, topf)), the (sh (stackC, Sf))],$   
 $vs[4 := Intg nargs, 5 := Intg i], safeP, sigSafeMain,$   
 $Suc (Suc (Suc (Suc (Suc (Suc pc))))), ref]))$ )  
**apply** (*unfold* *exec-all-def*)  
**apply** (*erule* *rtrancl-trans*)  
**prefer** 2  
**apply** (*rule* *r-into-rtrancl*)  
**apply** *clarify*  
**apply** (*unfold* *JVMExec.exec.simps*)  
**apply** (*unfold* *Let-def*)  
**apply** (*subgoal-tac* (*fst* (*snd* (*snd* (*snd* (*snd* (*the* (*method'* ( $P'$ , *safeP*) *sigSafe-*  
*Main*)))))) !

$Suc (Suc (Suc (Suc (Suc (Suc pc)))))) =$   
 $Getstatic\ topf\ stackC) \mathbf{apply}\ simp$   
**apply** (rule nth-via-drop-append) **apply** force

**apply** (drule drop-Suc-append)  
**apply** (subgoal-tac  
 $P \vdash (None, sh, h, inih,$   
 $[[the (sh (stackC, topf)), the (sh (stackC, Sf))],$   
 $vs[4 := Intg\ nargs, 5 := Intg\ i], safeP, sigSafeMain,$   
 $Suc (Suc (Suc (Suc (Suc (Suc (Suc pc)))))), ref]])$   
 $-jvm \rightarrow$   
 $(None, sh, h, inih,$   
 $[[Intg\ i, the (sh (stackC, topf)), the (sh (stackC, Sf))],$   
 $vs[4 := Intg\ nargs, 5 := Intg\ i], safeP, sigSafeMain,$   
 $Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc pc)))))), ref]])$ )  
**apply** (unfold exec-all-def)  
**apply** (erule rtrancl-trans)  
**prefer** 2  
**apply** (rule r-into-rtrancl)  
**apply** clarify  
**apply** (unfold JVMEexec.exec.simps)  
**apply** (unfold Let-def)  
**apply** (subgoal-tac (fst (snd (snd (snd (snd (the (method' (P', safeP) sigSafe-  
Main))))))) !  
 $Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc pc)))))) =$   
 $Load\ 5) \mathbf{apply}\ simp$   
**apply** (rule nth-via-drop-append) **apply** force

**apply** (drule drop-Suc-append)  
**apply** (subgoal-tac  
 $P \vdash (None, sh, h, inih,$   
 $[[Intg\ i, the (sh (stackC, topf)), the (sh (stackC, Sf))],$   
 $vs[4 := Intg\ nargs, 5 := Intg\ i], safeP, sigSafeMain,$   
 $Suc (Suc (Suc (Suc (Suc (Suc (Suc pc)))))), ref]])$   
 $-jvm \rightarrow$   
 $(None, sh, h, inih,$   
 $[[Intg (the-Intg (the (sh (stackC, topf))) - i), the (sh (stackC, Sf))],$   
 $vs[4 := Intg\ nargs, 5 := Intg\ i], safeP, sigSafeMain,$   
 $Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc pc)))))), ref]])$ )  
**apply** (unfold exec-all-def)  
**apply** (erule rtrancl-trans)  
**prefer** 2  
**apply** (rule r-into-rtrancl)  
**apply** clarify  
**apply** (unfold JVMEexec.exec.simps)  
**apply** (unfold Let-def)

**apply** (*subgoal-tac* (*fst* (*snd* (*snd* (*snd* (*snd* (*the* (*method'* (*P'*, *safeP*) *sigSafeMain*))))))) !

*Suc* (*Suc* (*Suc* (*Suc* (*Suc* (*Suc* (*Suc* (*Suc* (*Suc* (*Suc* (*Suc* (*pc*)))))))))) =  
*BinOp Subtract*) **apply** *simp*

**apply** (*rule nth-via-drop-append*) **apply** *force*

**apply** (*drule drop-Suc-append*)

**apply** (*subgoal-tac*

$P \vdash$  (*None*, *sh*, *h*, *inih*,

[[*Intg* (*the-Intg* (*the* (*sh* (*stackC*, *topf*))) - *i*), *the* (*sh* (*stackC*, *Sf*))],  
*vs*[*4*:= *Intg nargs*, *5*:= *Intg i*], *safeP*, *sigSafeMain*,  
*Suc* (*Suc* (*Suc* (*Suc* (*Suc* (*Suc* (*Suc* (*Suc* (*Suc* (*Suc* (*Suc* (*pc*))))))))))], *ref*]])

-*jvm*→

(*None*, *sh*, *h*, *inih*,  
[[*Addr l*, *Intg* (*the-Intg* (*the* (*sh* (*stackC*, *topf*))) - *i*), *the* (*sh* (*stackC*,  
*Sf*))],  
*vs*[*4*:= *Intg nargs*, *5*:= *Intg i*], *safeP*, *sigSafeMain*,  
*Suc* (*Suc* (*Suc* (*Suc* (*Suc* (*Suc* (*Suc* (*Suc* (*Suc* (*Suc* (*Suc* (*pc*))))))))))], *ref*]])

**apply** (*unfold exec-all-def*)

**apply** (*erule rtrancl-trans*)

**prefer** 2

**apply** (*rule r-into-rtrancl*)

**apply** *clarify*

**apply** (*unfold JVMExec.exec.simps*)

**apply** (*unfold Let-def*)

**apply** (*subgoal-tac* (*fst* (*snd* (*snd* (*snd* (*snd* (*the* (*method'* (*P'*, *safeP*) *sigSafeMain*))))))) !

*Suc* (*Suc* (*Suc* (*Suc* (*Suc* (*Suc* (*Suc* (*Suc* (*Suc* (*Suc* (*Suc* (*pc*)))))))))) =  
*Load 6*) **apply** *simp*

**apply** (*rule nth-via-drop-append*) **apply** *force*

**apply** (*drule drop-Suc-append*)

**apply** (*subgoal-tac*

$P \vdash$  (*None*, *sh*, *h*, *inih*,

[[*Addr l*, *Intg* (*the-Intg* (*the* (*sh* (*stackC*, *topf*))) - *i*), *the* (*sh* (*stackC*,  
*Sf*))],

*vs*[*4*:= *Intg nargs*, *5*:= *Intg i*], *safeP*, *sigSafeMain*,

*Suc* (*Suc* (*Suc* (*Suc* (*Suc* (*Suc* (*Suc* (*Suc* (*Suc* (*Suc* (*Suc* (*pc*))))))))))], *ref*]])

-*jvm*→

(*None*, *sh*, *h*, *inih*,  
[[*the* (*fs* (*arg1*, *cellC*)), *Intg* (*the-Intg* (*the* (*sh* (*stackC*, *topf*))) - *i*), *the* (*sh* (*stackC*,  
*Sf*))],

*vs*[*4*:= *Intg nargs*, *5*:= *Intg i*], *safeP*, *sigSafeMain*,

*Suc* (*Suc* (*Suc* (*Suc* (*Suc* (*Suc* (*Suc* (*Suc* (*Suc* (*Suc* (*Suc* (*pc*))))))))))], *ref*]])

**apply** (*unfold exec-all-def*)



```

apply (erule rtrancl-trans)
prefer 2
apply (rule r-into-rtrancl)
apply clarify
apply (unfold JVMExec.exec.simps)
apply (unfold Let-def)
apply (subgoal-tac (fst (snd (snd (snd (snd (the (method' (P', safeP) sigSafe-
Main))))))) !
      Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc pc))))))))))
=
      Getfield arg1 cellC) apply simp
apply (simp add: raise-system-xcpt-def)
apply (rule nth-via-drop-append) apply force

apply (drule drop-Suc-append)
apply (subgoal-tac
  P⊢ (None, sh, h, inih,
    [(the (fs (arg1, cellC)), Intg (the-Intg (the (sh (stackC, topf))) - i), the
      (sh (stackC, Sf))),
      vs[4:= Intg nargs, 5:= Intg i], safeP, sigSafeMain,
      Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc pe)))))))))), ref]])
    -jvm→
      (None, sh,
        h(the-Addr (the (sh (stackC, Sf))) ↦ Arr ty m (S'(nat (the-Intg (the (sh
          (stackC, topf))) - i)
          ↦ argCell fs (nat i + 1)))),
        inih,
        [([], vs[4:= Intg nargs, 5:= Intg i], safeP, sigSafeMain,
          Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc pe)))))))))),
          ref]]))
apply (unfold exec-all-def)
apply (erule rtrancl-trans)
prefer 2
apply (rule r-into-rtrancl)
apply clarify
apply (unfold JVMExec.exec.simps)
apply (unfold Let-def)
apply (subgoal-tac (fst (snd (snd (snd (snd (the (method' (P', safeP) sigSafe-
Main))))))) !
      Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc
pc)))))))))) =
      ArrStore) apply simp
apply (subgoal-tac (¬ int m ≤ the-Intg (the (sh (stackC, topf))) - i ∧ ¬ the-Intg
(the (sh (stackC, topf))) < i))
apply (simp add: raise-system-xcpt-def)
apply (subgoal-tac i=0)
prefer 2 apply (rule i-0)
apply (subgoal-tac the (fs (arg1, cellC)) = argCell fs (Suc (nat i)),simp)

```

**apply** (*simp add: argCell-def*)  
**apply** (*rule top-good*)  
**apply** (*rule nth-via-drop-append*) **apply force**

**apply** (*drule drop-Suc-append*)  
**apply** (*subgoal-tac*  
 $P \vdash (None, sh,$   
 $h(the-Addr (the (sh (stackC, Sf))) \mapsto Arr ty m (S'(nat (the-Intg (the (sh$   
 $(stackC, topf))) - i)$   
 $\mapsto argCell fs (nat i + 1))))),$   
 $inih,$   
 $[[([], vs[4 := Intg nargs, 5 := Intg i], safeP, sigSafeMain,$   
 $Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc pc)))))))))))]$   
 $ref)])$   
 $-jvm \rightarrow$   
 $(None, sh,$   
 $h(the-Addr (the (sh (stackC, Sf))) \mapsto Arr ty m (S'(nat (the-Intg (the (sh$   
 $(stackC, topf))) - i)$   
 $\mapsto argCell fs (nat i + 1))))),$   
 $inih,$   
 $[[([], vs[4 := Intg nargs, 5 := Intg i], safeP, sigSafeMain, nat (12 + int pc +$   
 $endlabel1), ref)])$ )  
**apply** (*unfold exec-all-def*)  
**apply** (*erule rtrancl-trans*)  
**prefer** 2  
**apply** (*rule r-into-rtrancl*)  
**apply** *clarify*  
**apply** (*unfold JVMExec.exec.simps*)  
**apply** (*unfold Let-def*)  
**apply** (*subgoal-tac (fst (snd (snd (snd (snd (the (method' (P', safeP) sigSafe-*  
 $Main)))))) !$   
 $Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc$   
 $pc)))))))))) =$   
 $Goto endlabel1)$  **apply simp**  
**apply** (*rule nth-via-drop-append*) **apply force**

**apply** (*frule drop-endlabel1*)  
**apply** (*unfold Match2-def*)  
**apply** (*subgoal-tac*  
 $P \vdash (None, sh,$   
 $h(the-Addr (the (sh (stackC, Sf))) \mapsto Arr ty m (S'(nat (the-Intg (the (sh$   
 $(stackC, topf))) - i)$   
 $\mapsto argCell fs (nat i + 1))))),$   
 $inih,$   
 $[[([], vs[4 := Intg nargs, 5 := Intg i], safeP, sigSafeMain, nat (12 + int pc +$   
 $endlabel1), ref)])$   
 $-jvm \rightarrow$

```

      (None, sh,
        h(the-Addr (the (sh (stackC, Sf)))  $\mapsto$  Arr ty m (S'(nat (the-Intg (the (sh
(stackC, topf))) - i)
           $\mapsto$  argCell fs (nat i + 1))))),
      inih,
      [([Intg i],
        vs[4:= Intg nargs, 5:= Intg i], safeP, sigSafeMain,
        Suc (nat (12 + int pc + endlabel1)), ref)])
apply (unfold exec-all-def)
apply (erule rtrancl-trans)
prefer 2
apply (rule r-into-rtrancl)
apply clarify
apply (unfold JVMExec.exec.simps)
apply (unfold Let-def)
apply (subgoal-tac (fst (snd (snd (snd (snd (the (method' (P', safeP) sigSafe-
Main))))))) !
      (nat (12 + int pc + endlabel1))) =
      Load 5) apply simp
apply (rule nth-via-drop-append) apply force

apply (drule drop-Suc-append)
apply (drule drop-Suc-append)
apply (subgoal-tac
  P $\vdash$  (None, sh,
    h(the-Addr (the (sh (stackC, Sf)))  $\mapsto$  Arr ty m (S'(nat (the-Intg (the (sh
(stackC, topf))) - i)
       $\mapsto$  argCell fs (nat i + 1))))),
    inih,
    [([Intg i],
      vs[4:= Intg nargs, 5:= Intg i], safeP, sigSafeMain,
      Suc (nat (12 + int pc + endlabel1)), ref)])
    -jvm $\rightarrow$ 
    (None, sh,
      h(the-Addr (the (sh (stackC, Sf)))  $\mapsto$  Arr ty m (S'(nat (the-Intg (the (sh
(stackC, topf))) - i)
         $\mapsto$  argCell fs (nat i + 1))))),
      inih,
      [([Intg 1, Intg i],
        vs[4:= Intg nargs, 5:= Intg i], safeP, sigSafeMain,
        Suc (Suc (nat (12 + int pc + endlabel1))), ref)])
apply (unfold exec-all-def)
apply (erule rtrancl-trans)
prefer 2
apply (rule r-into-rtrancl)
apply clarify
apply (unfold JVMExec.exec.simps)
apply (unfold Let-def)

```

**apply** (*subgoal-tac* (*fst* (*snd* (*snd* (*snd* (*snd* (*the* (*method'* (*P'*, *safeP*) *sigSafeMain*))))))) !

*Suc* (*nat* (*12* + *int pc* + *endlabel1*))) =  
*LitPush* (*Intg 1*) **apply simp**

**apply** (*rule nth-via-drop-append*) **apply force**

**apply** (*drule drop-Suc-append*)

**apply** (*subgoal-tac*

*P* ⊢ (*None*, *sh*,

*h*(*the-Addr* (*the* (*sh* (*stackC*, *Sf*)))) ⊢ *Arr ty m* (*S'*(*nat* (*the-Intg* (*the* (*sh* (*stackC*, *topf*))) – *i*)

⊢ *argCell fs* (*nat i* + 1))),

*inih*,

[[*Intg 1*, *Intg i*],

*vs*[*4*:= *Intg nargs*, *5*:= *Intg i*], *safeP*, *sigSafeMain*,

*Suc* (*Suc* (*nat* (*12* + *int pc* + *endlabel1*))), *ref*)]

–*jvm*→

(*None*, *sh*,

*h*(*the-Addr* (*the* (*sh* (*stackC*, *Sf*)))) ⊢ *Arr ty m* (*S'*(*nat* (*the-Intg* (*the* (*sh* (*stackC*, *topf*))) – *i*)

⊢ *argCell fs* (*nat i* + 1))),

*inih*,

[[*Intg* (*i* + 1)],

*vs*[*4*:= *Intg nargs*, *5*:= *Intg i*], *safeP*, *sigSafeMain*,

*Suc* (*Suc* (*Suc* (*nat* (*12* + *int pc* + *endlabel1*))), *ref*)]])

**apply** (*unfold exec-all-def*)

**apply** (*erule rtrancl-trans*)

**prefer** 2

**apply** (*rule r-into-rtrancl*)

**apply** *clarify*

**apply** (*unfold JVMEexec.exec.simps*)

**apply** (*unfold Let-def*)

**apply** (*subgoal-tac* (*fst* (*snd* (*snd* (*snd* (*snd* (*the* (*method'* (*P'*, *safeP*) *sigSafeMain*))))))) !

*Suc* (*Suc* (*nat* (*12* + *int pc* + *endlabel1*)))) =

*BinOp Add*) **apply simp**

**apply** (*rule nth-via-drop-append*) **apply force**

**apply** (*drule drop-Suc-append*)

**apply** (*subgoal-tac*

*P* ⊢ (*None*, *sh*,

*h*(*the-Addr* (*the* (*sh* (*stackC*, *Sf*)))) ⊢ *Arr ty m* (*S'*(*nat* (*the-Intg* (*the* (*sh* (*stackC*, *topf*))) – *i*)

⊢ *argCell fs* (*nat i* + 1))),

*inih*,

[[*Intg* (*i* + 1)],

$vs[4 := \text{Intg nargs}, 5 := \text{Intg } i], \text{ safeP}, \text{ sigSafeMain},$   
 $\text{Suc} (\text{Suc} (\text{Suc} (\text{nat} (12 + \text{int pc} + \text{endlabel1}))))), \text{ ref}]]$   
 $-jvm \rightarrow$   
 $(\text{None}, \text{ sh}, h(\text{the-Addr} (\text{the} (\text{sh} (\text{stackC}, \text{ Sf})))) \mapsto \text{Arr ty m} (S'(\text{nat} (\text{the-Intg}$   
 $(\text{the} (\text{sh} (\text{stackC}, \text{ topf}))) - i)$   
 $\mapsto \text{argCell fs} (\text{nat } i + 1))))),$   
 $\text{ inih},$   
 $[(\square,$   
 $vs[4 := \text{Intg nargs}, 5 := \text{Intg } (i + 1)], \text{ safeP}, \text{ sigSafeMain},$   
 $\text{Suc} (\text{Suc} (\text{Suc} (\text{Suc} (\text{nat} (12 + \text{int pc} + \text{endlabel1}))))), \text{ ref}]]$   
**apply** (*unfold exec-all-def*)  
**apply** (*erule rtrancl-trans*)  
**prefer** 2  
**apply** (*rule r-into-rtrancl*)  
**apply** (*clarify*)  
**apply** (*unfold JVMExec.exec.simps*)  
**apply** (*unfold Let-def*)  
**apply** (*subgoal-tac* (*fst* (*snd* (*snd* (*snd* (*snd* (*the* (*method'* (*P'*, *safeP*) *sigSafe-* *Main*)))))) !  
 $\text{Suc} (\text{Suc} (\text{Suc} (\text{nat} (12 + \text{int pc} + \text{endlabel1})))) =$   
 $\text{Store } 5)$  **apply** (*simp*)  
**apply** (*rule nth-via-drop-append*) **apply** (*force*)

**apply** (*drule drop-Suc-append*)  
**apply** (*subgoal-tac*  
 $P \vdash (\text{None}, \text{ sh}, h(\text{the-Addr} (\text{the} (\text{sh} (\text{stackC}, \text{ Sf})))) \mapsto \text{Arr ty m} (S'(\text{nat} (\text{the-Intg}$   
 $(\text{the} (\text{sh} (\text{stackC}, \text{ topf}))) - i)$   
 $\mapsto \text{argCell fs} (\text{nat } i + 1))))),$   
 $\text{ inih},$   
 $[(\square,$   
 $vs[4 := \text{Intg nargs}, 5 := \text{Intg } (i + 1)], \text{ safeP}, \text{ sigSafeMain},$   
 $\text{Suc} (\text{Suc} (\text{Suc} (\text{Suc} (\text{nat} (12 + \text{int pc} + \text{endlabel1}))))), \text{ ref}]]$   
 $-jvm \rightarrow$   
 $(\text{None}, \text{ sh},$   
 $h(\text{the-Addr} (\text{the} (\text{sh} (\text{stackC}, \text{ Sf})))) \mapsto \text{Arr ty m} (S'(\text{nat} (\text{the-Intg} (\text{the} (\text{sh}$   
 $(\text{stackC}, \text{ topf}))) - i)$   
 $\mapsto \text{argCell fs} (\text{nat } i + 1))))),$   
 $\text{ inih},$   
 $[(\square,$   
 $vs[4 := \text{Intg nargs}, 5 := \text{Intg } (i + 1)],$   
 $\text{ safeP}, \text{ sigSafeMain},$   
 $\text{nat} (16 + (\text{int pc} + \text{endlabel1}) + \text{labelLoop}), \text{ ref}]]$   
**apply** (*unfold exec-all-def*)  
**apply** (*erule rtrancl-trans*)  
**prefer** 2  
**apply** (*rule r-into-rtrancl*)

```

apply clarify
apply (unfold JVMExec.exec.simps)
apply (unfold Let-def)
apply (subgoal-tac (fst (snd (snd (snd (snd (the (method' (P', safeP) sigSafe-
Main))))))) !
      Suc (Suc (Suc (Suc (nat (12 + int pc + endlabel1)))))) =
      Goto labelLoop) apply simp
apply (simp add: endlabel1-def)
apply (rule nth-via-drop-append) apply force
apply (simp add: endlabel1-def)
apply (simp add: labelLoop-def)
apply (simp add: nlab1-def)
apply (simp add: InstLabel1-def)
apply (simp add: nMatch2-def)
apply (simp add: nMatch12-def)
by (simp add: Match12-def)

```

**lemma** oneStep-state:

```

[[ (P', cdm, ctm, com) = trSVM2JVM ((svms, ctmap), ini, ct, ah, ai, bc);
  length vs = 10;
  vs ! 6 = Addr l;
  sh (stackC, Sf) = Some (Addr la);
  l ≠ la;
  h la = Some (Arr ty m S');
  h l = Some (Obj cellC fs);
  drop pc (fst (snd (snd (snd (snd (the (method' (P', safeP) sigSafeMain)))))))
=
  Match12 @ InstLabel1 @ InstLabel2 @ InstLabel3 @ InstLabel4 @ InstLabel5
@
  InstLabel6 @ InstLabel7 @ InstLabel8 @ Match2 @ bytecode;
  i < nargs ]]
⇒ oneStep P' (None, sh, h, inih, [[[], vs[4:= Intg nargs, 5:= Intg i], safeP,
sigSafeMain, pc, ref]])
  = (None, sh, h(the-Addr (the (sh (stackC, Sf))) ↦
      Arr ty m (S'(nat (the-Intg (the (sh (stackC,
topf))) - i) ↦ argCell fs (nat i + 1))))),
      inih, [[[], vs[4:= Intg nargs, 5:= Intg (i + 1)], safeP,
sigSafeMain, pc, ref]])
apply (simp (no-asm) only: oneStep-def)
apply (rule the-equality)
apply (rule conjI)
apply (rule oneStep-MATCH,assumption+,simp)
apply auto
apply (rename-tac v1' sh' dh' ih' frms')
apply (case-tac frms',simp)
apply (frule undefined-false,simp)
apply simp
apply (case-tac a,simp)

```

```

apply (case-tac aa,simp)
apply (case-tac b,simp)
apply (case-tac ba,simp)
apply (case-tac bb,simp)
apply (case-tac bca,simp)
apply (case-tac list,simp)
apply (simp,frule undefined-false,simp)
apply (simp,frule undefined-false,simp)
apply (case-tac y,simp)
apply (frule undefined-false,simp)
apply simp
apply (case-tac a,simp)
apply (case-tac aa,simp)
apply (case-tac b,simp)
apply (case-tac ba,simp)
apply (case-tac bb,simp)
apply (case-tac bca,simp)
apply (case-tac list,simp)
apply (simp,frule undefined-false,simp)
apply (simp,frule undefined-false,simp)
apply (case-tac y,simp)
apply (frule undefined-false,simp)
apply simp
apply (case-tac a,simp)
apply (case-tac aa,simp)
apply (case-tac b,simp)
apply (case-tac ba,simp)
apply (case-tac bb,simp)
apply (case-tac bca,simp)
apply (case-tac list,simp)
apply (simp,frule undefined-false,simp)
apply (simp,frule undefined-false,simp)
apply (rename-tac v1' sh' dh' ih' frms')
apply (case-tac frms')
apply simp
apply (frule undefined-false,simp)
apply simp
apply (case-tac a,simp)
apply (case-tac aa,simp)
apply (case-tac b,simp)
apply (case-tac ba,simp)
apply (case-tac bb,simp)
apply (case-tac bca,simp)
apply (case-tac list,simp)
apply (simp,frule undefined-false,simp)
apply (simp,frule undefined-false,simp)
apply (case-tac y)
apply simp
apply (frule undefined-false,simp)

```

```

apply simp
apply (case-tac a, simp)
apply (case-tac aa, simp)
apply (case-tac b, simp)
apply (case-tac ba, simp)
apply (case-tac bb, simp)
apply (case-tac bca, simp)
apply (case-tac list, simp)
apply (simp, frule undefined-false, simp)
by (simp, frule undefined-false, simp)

```

**lemma** *diff-monotone*:

```

[[ (P', cdm, ctm, com) = trSVM2JVM ((svms, ctmap), ini, ct, ah, ai, bc);
  length vs = 10;
  vs ! 6 = Addr l;
  sh (stackC, Sf) = Some (Addr la);
  h la = Some (Arr ty m S');
  h l = Some (Obj cellC fs);
  l ≠ la;
  i < nargs;
  drop pc (fst (snd (snd (snd (the (method' (P', safeP) sigSafeMain)))))))]
=
  Match12 @ InstLabel1 @ InstLabel2 @ InstLabel3 @ InstLabel4 @ InstLabel5
  @
    InstLabel6 @ InstLabel7 @ InstLabel8 @ Match2 @ bytecode;
  Suc n = diff (None, sh, h, inih, [[[], vs[ 4:= Intg nargs, 5 := Intg i], safeP,
  sigSafeMain, pc, ref]]) ]
  ⇒ n = diff (oneStep P' (None, sh, h, inih, [[[], vs[ 4:= Intg nargs, 5 := Intg
  i], safeP, sigSafeMain, pc, ref]]))
apply (simp add: diff-def)
apply (subgoal-tac
  oneStep P' (None, sh, h, inih, [[[], vs[ 4:= Intg nargs, 5 := Intg i], safeP,
  sigSafeMain, pc, ref]]) =
  (None, sh, h(the-Addr (the (sh (stackC, Sf))) ↦
  Arr ty m (S'(nat (the-Intg (the (sh (stackC,
  topf))) - i) ↦ argCell fs (nat i + 1))))),
  inih, [[[], vs[4:= Intg nargs, 5:= Intg (i + 1)], safeP,
  sigSafeMain, pc, ref]]), simp)
by (rule oneStep-state, assumption+)

```

**lemma** *steps-monotone*:

```

[[ (P', cdm, ctm, com) = trSVM2JVM ((svms, ctmap), ini, ct, ah, ai, bc);
  length vs = 10;
  vs ! 6 = Addr l;
  sh (stackC, Sf) = Some (Addr la);
  h la = Some (Arr ty m S');
  h l = Some (Obj cellC fs);
  l ≠ la;

```



$i < nargs;$   
 $drop\ pc\ (fst\ (snd\ (snd\ (snd\ (snd\ (the\ (method'\ (P',\ safeP)\ sigSafeMain)))))))))$   
 $=$   
 $Match12\ @\ InstLabel1\ @\ InstLabel2\ @\ InstLabel3\ @\ InstLabel4\ @\ InstLabel5$   
 $@$   
 $InstLabel6\ @\ InstLabel7\ @\ InstLabel8\ @\ Match2\ @\ bytecode;$   
 $Suc\ n = diff\ (None,\ sh,\ h,\ inih,\ [([],\ vs[4 := Intg\ nargs,\ 5 := Intg\ i],\ safeP,$   
 $sigSafeMain,\ pc,\ ref)])$   
 $\implies\ steps\ (diff\ (None,\ sh,\ h,\ inih,\ [([],\ vs[4 := Intg\ nargs,\ 5 := Intg\ i],\ safeP,$   
 $sigSafeMain,\ pc,\ ref))))\ P'$   
 $(None,\ sh,\ h,\ inih,\ [([],\ vs[4 := Intg\ nargs,\ 5 := Intg\ i],\ safeP,$   
 $sigSafeMain,\ pc,\ ref))) =$   
 $steps\ (diff\ (oneStep\ P'\ (None,\ sh,\ h,\ inih,\ [([],\ vs[4 := Intg\ nargs,\ 5 := Intg$   
 $i],\ safeP,\ sigSafeMain,\ pc,\ ref))))\ P'$   
 $(oneStep\ P'\ (None,\ sh,\ h,\ inih,\ [([],\ vs[4 := Intg\ nargs,\ 5 := Intg$   
 $i],\ safeP,\ sigSafeMain,\ pc,\ ref))))$   
 $apply\ (frule\ diff-monotone,assumption+)$   
 $by\ (case-tac\ (diff\ (None,\ sh,\ h,\ inih,\ [([],\ vs[4 := Intg\ nargs,\ 5 := Intg\ i],\ safeP,$   
 $sigSafeMain,\ pc,\ ref)])),simp,simp)$

**lemma** *MATCH-loop* [rule-format]:

$\forall\ sh\ h\ inih\ vs\ i\ pc\ ref\ S'.$   
 $n = diff\ (None,\ sh,\ h,\ inih,\ [([],\ vs[4 := Intg\ nargs,\ 5 := Intg\ i],\ safeP,$   
 $sigSafeMain,\ pc,\ ref)])$   
 $\longrightarrow\ (P',\ cdm,\ ctm,\ com) = trSVM2JVM\ ((svms,\ ctmap),\ ini,\ ct,\ ah,\ ai,\ bc)$   
 $\longrightarrow\ length\ vs = 10$   
 $\longrightarrow\ vs\ !\ 6 = Addr\ l$   
 $\longrightarrow\ sh\ (stackC,\ Sf) = Some\ (Addr\ la)$   
 $\longrightarrow\ h\ la = Some\ (Arr\ ty\ m\ S')$   
 $\longrightarrow\ h\ l = Some\ (Obj\ cellC\ fs)$   
 $\longrightarrow\ l \neq la$   
 $\longrightarrow\ drop\ pc\ (fst\ (snd\ (snd\ (snd\ (snd\ (the\ (method'\ (P',\ safeP)\ sigSafe-$   
 $Main))))))))) =$   
 $Match12\ @\ InstLabel1\ @\ InstLabel2\ @\ InstLabel3\ @\ InstLabel4\ @\ InstLabel5$   
 $@$   
 $InstLabel6\ @\ InstLabel7\ @\ InstLabel8\ @\ Match2\ @\ bytecode$   
 $\longrightarrow\ P' \vdash (None,\ sh,\ h,\ inih,\ [([],\ vs[4 := Intg\ nargs,\ 5 := Intg\ i],\ safeP,$   
 $sigSafeMain,\ pc,\ ref)])$   
 $\quad -jvm \longrightarrow steps\ n\ P'\ (None,\ sh,\ h,\ inih,\ [([],\ vs[4 := Intg\ nargs,\ 5 := Intg\ i],$   
 $safeP,\ sigSafeMain,\ pc,\ ref)])$   
 $apply\ (induct\ n)$   
 $apply\ (simp\ add:\ steps.simps)$   
 $apply\ (simp\ add:\ exec-all-def)$   
 $apply\ (rule\ allI)+$   
 $apply\ (rule\ impI)$

$apply\ (case-tac\ i < nargs)$

**apply** (*erule-tac*  $x=sh$  **in** *allE*)  
**apply** (*erule-tac*  $x=h(\text{the-Addr } (\text{the } (sh \text{ (stackC, Sf))}) \mapsto$   
 $\text{Arr ty m } (S'(\text{nat } (\text{the-Intg } (\text{the } (sh \text{ (stackC, topf))}) - i) \mapsto$   
 $\text{argCell fs } (\text{nat } i + 1))))$  **in** *allE*)  
**apply** (*erule-tac*  $x=inih$  **in** *allE*)  
**apply** (*erule-tac*  $x=vs$  **in** *allE*)  
**apply** (*erule-tac*  $x=i + 1$  **in** *allE*)  
**apply** (*erule-tac*  $x=pc$  **in** *allE*)  
**apply** (*erule-tac*  $x=ref$  **in** *allE*)  
**apply** (*erule-tac*  $x=S'(\text{nat } (\text{the-Intg } (\text{the } (sh \text{ (stackC, topf))}) - i) \mapsto \text{argCell fs}$   
 $(\text{nat } i + 1))$  **in** *allE*)  
**apply** (*rule impI*)  
**apply** (*drule mp*)  
**apply** (*frule diff-monotone,assumption+*)  
**apply** (*frule-tac*  $inih=inih$  **and**  $nargs=nargs$  **and**  $i=i$  **and**  $ref=ref$  **in** *oneStep-state,assumption+*)

**apply** (*simp add: diff-def*)  
**apply** (*drule mp,force+*)  
**apply** (*frule diff-monotone,assumption+*)  
**apply** (*frule-tac*  
 $inih=inih$  **and**  $nargs=nargs$  **and**  $i=i$  **and**  $ref=ref$   
**in** *oneStep-MATCH,assumption+,simp,assumption*)  
**apply** (*frule-tac*  
 $inih=inih$  **and**  $nargs=nargs$  **and**  $i=i$  **and**  $ref=ref$   
**in** *oneStep-state,assumption+*)  
**apply** (*drule-tac*  
 $s=oneStep P' (None, sh, h, inih, [([], vs[4 := Intg nargs, 5 := Intg i],$   
 $safeP, sigSafeMain, pc, ref)])$  **in** *sym*)  
**apply** *clarify*  
**apply** *simp*  
**apply** (*frule steps-monotone, assumption+*)  
**apply** (*drule-tac*  
 $s=steps (diff (None, sh, h, inih, [([], vs[4 := Intg nargs, 5 := Intg i], safeP,$   
 $sigSafeMain, pc, a, b)]) P'$   
 $(None, sh, h, inih, [([], vs[4 := Intg nargs, 5 := Intg i], safeP,$   
 $sigSafeMain, pc, a, b)])$  **in** *sym,simp*)  
**apply** (*simp add: exec-all-def*)

**apply** (*rule impI*)  
**by** (*simp add: diff-def*)

**lemma** *execSVMInstr-MATCH-loop*:

$\llbracket (P', \text{cdm}, \text{ctm}, \text{com}) = \text{trSVM2JVM } ((\text{svms}, \text{ctmap}), \text{ini}, \text{ct}, \text{ah}, \text{ai}, \text{bc});$   
 $\text{length } vs = 10;$   
 $vs ! 6 = \text{Addr } l;$   
 $sh \text{ (stackC, Sf)} = \text{Some } (\text{Addr } la);$   
 $h \text{ la} = \text{Some } (\text{Arr ty m } S');$

```

    h l = Some (Obj cellC fs);
    l ≠ la;
    drop pc (fst (snd (snd (snd (snd (the (method' (P', safeP) sigSafeMain)))))))
  =
    Match12 @ InstLabel1 @ InstLabel2 @ InstLabel3 @ InstLabel4 @ InstLabel5
  @
    InstLabel6 @ InstLabel7 @ InstLabel8 @ Match2 @ bytecode;
    s = (None, sh, h, inih, [[[], vs[ 4:= Intg nargs, 5:= Intg i], safeP, sigSafeMain,
pc, ref]]) ]
    ⇒ P' ⊢ s -jvm→ lastState P' s
apply (simp add: lastState-def)
apply (intro MATCH-loop)
apply (simp add: diff-def)
by assumption+

```

**axioms** map-of-update-add:

```

x ∉ set xs
⇒ S(x ↦ y) ++ map-of (zip xs ys) = S ++ map-of (zip (xs@[x]) (ys@[y]))

```

**axioms** map-upt-rev:

```

m < n ⇒ map f [n>..Suc m] @ [f m] = map f [n>..m]

```

**lemma** map-argCell-upt-rev:

```

[[ i < nargs; i ≥ 0 ]]
⇒ (map (λj. argCell fs (Suc j)) [nat nargs>..nat (i + 1)]) @ [argCell fs (Suc
(nat i))]
  = (map (λj. argCell fs (Suc j)) [nat nargs>..nat i])
apply (subgoal-tac nat (i + 1) = Suc (nat i), simp)
apply (subgoal-tac nat i < nat nargs)
apply (rule-tac f=(λj. argCell fs (Suc j)) in map-upt-rev, assumption)
by force+

```

**lemma** empty-map-add-2: m ++ empty = m

**by** (rule ext) (simp add: map-add-def split: option.split)

**lemma** upt-Suc-append-topf:

```

[[ i < nargs; i ≥ 0; (the-Intg (the (sh (stackC, topf))) - nargs) ≥ 0 ]]
⇒ [Suc (nat (the-Intg (the (sh (stackC, topf))) - nargs))..<nat (the-Intg (the
(sh (stackC, topf))) - i)] @
  [nat (the-Intg (the (sh (stackC, topf))) - i)]
  = [Suc (nat (the-Intg (the (sh (stackC, topf))) - nargs))..<nat (the-Intg (the
(sh (stackC, topf))) - i + 1)]
apply (rule sym)
apply (subgoal-tac

```

$\text{nat (the-Intg (the (sh (stackC, topf))) - i + 1) =}$   
 $\text{Suc (nat (the-Intg (the (sh (stackC, topf))) - i)),simp)}$   
**by** *arith+*

**lemma** *execSVMInstr-MATCH-loop-post* [rule-format]:

$\forall$  *sh h inih vs i pc ref S'*.  
 $n = \text{diff (None, sh, h, inih, [([], vs[ 4:= Intg nargs, 5:= Intg i], safeP, sigSafeMain, pc, ref])}$   
 $\longrightarrow (P', \text{cdm, ctm, com}) = \text{trSVM2JVM ((svms, ctmap), ini, ct, ah, ai, bc)}$   
 $\longrightarrow \text{length vs} = 10$   
 $\longrightarrow i \geq 0$   
 $\longrightarrow (\text{the-Intg (the (sh (stackC, topf))) - nargs}) \geq 0$   
 $\longrightarrow \text{nargs} \geq i$   
 $\longrightarrow \text{vs ! 6} = \text{Addr } l$   
 $\longrightarrow \text{sh (stackC, Sf)} = \text{Some (Addr } la)$   
 $\longrightarrow h \text{ la} = \text{Some (Arr ty m } S')$   
 $\longrightarrow h \text{ l} = \text{Some (Obj cellC fs)}$   
 $\longrightarrow l \neq la$   
 $\longrightarrow \text{drop pc (fst (snd (snd (snd (snd (the (method' (P', safeP) sigSafeMain)))))))} =$   
 $\text{Match12 @ InstLabel1 @ InstLabel2 @ InstLabel3 @ InstLabel4 @ InstLabel5}$   
 $\text{@}$   
 $\text{InstLabel6 @ InstLabel7 @ InstLabel8 @ Match2 @ bytecode}$   
 $\longrightarrow \text{steps } n \text{ P' (None, sh, h, inih, [([], vs[ 4:= Intg nargs, 5:= Intg i], safeP, sigSafeMain, pc, ref])} =$   
 $(\text{None, sh,}$   
 $\text{(let vitems = map (\%j. argCell fs (j + 1)) [nat nargs>..nat i];}$   
 $\text{idxs = map (\%j. j}$   
 $\text{[nat (the-Intg (the (sh (stackC, topf))) - nargs) + 1..<nat}$   
 $\text{(the-Intg (the (sh (stackC, topf))) - i + 1) ]};$   
 $\text{S'' = S' ++ map-of (zip idxs vitems)}$   
 $\text{in h(the-Addr (the (sh (stackC, Sf))) \mapsto Arr ty m S''),}$   
 $\text{in ih, [([], vs[ 4:= Intg nargs, 5:= Intg nargs],}$   
 $\text{safeP, sigSafeMain, pc, ref])}$   
**apply** (*induct n*)  
**apply** (*simp add: steps.simps, clarsimp*)  
**apply** (*simp add: diff-def*)  
**apply** (*subgoal-tac nargs = i, clarsimp*)  
**apply** (*subst upt-conv-Nil, simp*)  
**apply** (*subst zip-Nil*)  
**apply** (*subst map-of.simps, subst empty-map-add-2*)  
**apply** (*rule ext, force*)  
**apply** *simp*  
  
**apply** (*rule allI*)  
**apply** (*rule impI*)  
  
**apply** (*case-tac i < nargs*)

```

apply (erule-tac x=sh in allE)
apply (erule-tac x=h(the-Addr (the (sh (stackC, Sf)))  $\mapsto$  Arr ty m
      (S'(nat (the-Intg (the (sh (stackC, topf))) - i)  $\mapsto$  argCell fs
        (nat i + 1)))) in allE)
apply (erule-tac x=inih in allE)
apply (erule-tac x=vs in allE)
apply (erule-tac x=i + 1 in allE)
apply (erule-tac x=pc in allE)
apply (erule-tac x=ref in allE)
apply (erule-tac x=(S'(nat (the-Intg (the (sh (stackC, topf))) - i)  $\mapsto$  argCell fs
      (nat i + 1))) in allE)
apply (drule mp)
apply (frule diff-monotone,assumption+)
apply (frule-tac inih=inih and nargs=nargs and i=i and ref=ref in oneStep-state,assumption+)

apply (simp add: diff-def)
apply (drule mp,force)+

apply (frule diff-monotone,assumption+)
apply (frule-tac
      inih=inih and nargs=nargs and i=i and ref=ref
      in oneStep-MATCH,assumption+,simp,assumption)
apply (frule-tac
      inih=inih and nargs=nargs and i=i and ref=ref
      in oneStep-state,assumption+)
apply (drule-tac
      s=oneStep P' (None, sh, h, inih, [([]), vs[4 := Intg nargs, 5 := Intg i],
      safeP, sigSafeMain, pc, ref])) in sym)
apply clarify
apply simp
apply (frule steps-monotone, assumption+)
apply (drule-tac
      s=steps (diff (None, sh, h, inih, [([]), vs[4 := Intg nargs, 5 := Intg i], safeP,
      sigSafeMain, pc, a, b])) P'
      (None, sh, h, inih, [([]), vs[4 := Intg nargs, 5 := Intg i], safeP,
      sigSafeMain, pc, a, b])) in sym,simp)
apply (drule-tac
      t=oneStep P' (None, sh, h, inih, [([]), vs[4 := Intg nargs, 5 := Intg i],
      safeP, sigSafeMain, pc, a, b])) in sym)
apply (simp add: diff-def)
apply (subst map-of-update-add,simp)
apply (subst upt-Suc-append-topf,assumption+)
apply (subst map-argCell-upt-rev,assumption+)
apply (rule refl)

by (simp add: diff-def)

```

**lemma** *execSVMInstr-MATCH-loop-state*:

```

[[ (P', cdm, ctm, com) = trSVM2JVM ((svms, ctmap), ini, ct, ah, ai, bc);
  length vs = 10;
  vs ! 6 = Addr l;
  sh (stackC, Sf) = Some (Addr la);
  h la = Some (Arr ty m S');
  h l = Some (Obj cellC fs);
  l ≠ la;
  0 ≤ i;
  0 ≤ the-Intg (the (sh (stackC, topf))) - nargs;
  i ≤ nargs;
  drop pc (fst (snd (snd (snd (snd (the (method' (P', safeP) sigSafeMain))))))))
=
  Match12 @ InstLabel1 @ InstLabel2 @ InstLabel3 @ InstLabel4 @ InstLabel5
  @
    InstLabel6 @ InstLabel7 @ InstLabel8 @ Match2 @ bytecode;
  s = (None, sh, h, inih, [[([], vs[ 4:= Intg nargs, 5:= Intg i], safeP, sigSafeMain,
pc, ref)]) ]
⇒ P' ⊢ s -jvm→ (None, sh,
  (let vitems = map (%j. argCell fs (j + 1)) [nat nargs>..nat i];
    idxs = map (%j. j)
      [nat (the-Intg (the (sh (stackC, topf))) - nargs) + 1..<nat
(the-Intg (the (sh (stackC, topf))) - i + 1) ];
    S'' = S' ++ map-of (zip idxs vitems)
      in h(the-Addr (the (sh (stackC, Sf))) ↦ Arr ty m S'')),
  inih, [[([], vs[ 4:= Intg nargs, 5:= Intg nargs],
safeP, sigSafeMain, pc, ref)])
apply (frule execSVMInstr-MATCH-loop,assumption+)
apply (simp add: lastState-def)
apply (subgoal-tac
  (nat (nargs - i) = diff (None, sh, h, inih, [[([], vs[4 := Intg nargs, 5 := Intg
i], safeP, sigSafeMain, pc, ref)])))
apply (frule execSVMInstr-MATCH-loop-post,assumption+)
apply (simp only: diff-def,force)
by (simp add: diff-def)

```

**axioms** *equivS-MATCH-bound*:

```

[[ (fst (the (map-of svms l)) ! i) = MATCH off ps;
  sh (stackC,topf) = Some (Intg n) ] ⇒ int off ≤ n

```

**axioms** *maxPush-MATCH*:

```

fst (the (map-of svms l)) ! i = MATCH off ps ⇒ n + nargs' < int m

```

**axioms** *equivS-S'-n-minus-off*:

$\llbracket \text{equivS } S \ S' \ n \ \text{ctm } d \ g; \ \text{int } \text{off} \leq n; \ h \ l' = \text{Some } (\text{Arr } \text{ty } m \ S') \rrbracket$   
 $\implies \exists \ l \ \text{fs-S'-n-minus-off}.$   
 $\text{the } (d \ (\text{nat } (\text{the-Intg } (\text{the } (S' \ (\text{nat } (n - \text{int } \text{off})))))) = \text{Addr } l \wedge$   
 $h \ l = \text{Some } (\text{Obj } \text{cellC } \text{fs-S'-n-minus-off}) \wedge$   
 $l' \neq l$

**axioms** *consTableC-def*:

$\text{sha } (\text{consTableC}, \text{tablef}) = \text{Some } (\text{Addr } l''')$   
 $h \ l''' = \text{Some } (\text{Arr } \text{ty } m''' \ \text{consTableS})$

**axioms** *consTableC-n-minus-off*:

$\llbracket \text{sha } (\text{consTableC}, \text{tablef}) = \text{Some } (\text{Addr } l''');$   
 $h \ l''' = \text{Some } (\text{Arr } \text{ty } m''' \ \text{consTableS}) \rrbracket$   
 $\implies \exists \ l' \ \text{fs-consTableS-S'-n-minus-off-tagGf}.$   
 $\text{the } (\text{consTableS } (\text{nat}$   
 $(\text{the-Intg } (\text{the } (\text{fs-S'-n-minus-off } (\text{tagGf}, \text{cellC})))))) = \text{Addr } l' \wedge$   
 $h \ l' = \text{Some } (\text{Obj } \text{cellC } \text{fs-consTableS-S'-n-minus-off-tagGf}) \wedge$   
 $\text{the } (\text{fs-consTableS-S'-n-minus-off-tagGf } (\text{nargsf}, \text{consDataC}))$   
 $= \text{Intg } \text{nargs}' \wedge$   
 $\text{nargs}' \geq 0$

**axioms** *consTable-bounds*:

$\text{int } m''' > \text{the-Intg } (\text{the } (\text{fs-S'-n-minus-off } (\text{tagGf}, \text{cellC}))) \wedge$   
 $\text{the-Intg } (\text{the } (\text{fs-S'-n-minus-off } (\text{tagGf}, \text{cellC}))) \geq 0$

**axioms** *MATCH-bounds*:

$0 \leq i \wedge$   
 $0 \leq \text{the-Intg } (\text{the } (\text{sh } (\text{stackC}, \text{topf}))) - \text{nargs} \wedge$   
 $i \leq \text{nargs}$

**axioms** *tag-bounds*:

$\neg \text{the-Intg } (\text{the } (\text{fs-consTableS-S'-n-minus-off-tagGf}$   
 $(\text{tagLf}, \text{consDataC}))) < 0 \wedge$   
 $\neg \text{int } (\text{length } \text{ps} - \text{Suc } 0) <$   
 $\text{the-Intg } (\text{the } (\text{fs-consTableS-S'-n-minus-off-tagGf}$   
 $(\text{tagLf}, \text{consDataC})))$

**lemma** *drop-append-3*:

$\text{drop } n \ xs = [] \ @ \ (\text{ys} \ @ \ \text{zs}) \implies \text{drop } n \ xs = \text{ys} \ @ \ \text{zs}$   
**by** *simp*

**lemma** *drop-Suc-2*:

$\text{drop } n \ xs = (y \ # \ \text{ys}) \implies \text{drop } (\text{Suc } n) \ xs = \text{ys}$   
**apply** (*induct xs arbitrary: n, simp*)  
**apply** (*simp add:drop-Cons nth-Cons split:nat.splits*)  
**done**

**lemma** *drop-length-append*:

*drop n xs = (ys @ zs)  $\implies$  drop (n + length ys) xs = zs*  
**by** (*induct n arbitrary: xs*) (*auto, case-tac xs, auto*)

**lemma** *go-End*:

```

drop (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc
  (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc
  (Suc (Suc (Suc (Suc (Suc (the (cdm (l, i)))))))))))))))))))))) =
[Ifcmp GreaterEqual labelEndLoop, Load 5, Tableswitch 0 7
[label1, label2, label3, label4, label5, label6, label7, label8]] @
InstLabel1 @
InstLabel2 @
InstLabel3 @
InstLabel4 @
InstLabel5 @
InstLabel6 @
InstLabel7 @
InstLabel8 @
Match2 @
Load 3 #
Tableswitch 0 (int (length ps - Suc 0))
  (map ( $\lambda n$ . trAddr n (the (cdm (l, i)) + incMatch))
  (map ( $\lambda p$ . the (cdm' (p, 0))) ps)) #
bytecode'  $\implies$ 
drop (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc
  (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc
  (Suc (Suc (Suc (Suc (Suc (the (cdm (l, i)) + nat labelEndLoop))))))))))))))))))
  (fst (snd (snd (snd (snd (the (method' (P', safeP) sigSafeMain)))))))
= Load 3 #
Tableswitch 0 (int (length ps - Suc 0))
  (map ( $\lambda n$ . trAddr n (the (cdm (l, i)) + incMatch))
  (map ( $\lambda p$ . the (cdm' (p, 0))) ps)) #
bytecode'
apply (drule drop-Suc-append)+
apply (drule drop-append-3)
apply (drule drop-length-append)
apply (drule drop-length-append)
apply (drule drop-length-append)
apply (drule drop-length-append)
apply (drule drop-length-append)
apply (drule drop-length-append)
apply (drule drop-length-append)
apply (drule drop-length-append)
apply (drule drop-length-append)
apply (simp add: labelEndLoop-def)
apply (simp add: nlab1-def)

```



```

apply (simp add: nlab2-def)
apply (simp add: nlab3-def)
apply (simp add: nlab4-def)
apply (simp add: nlab5-def)
apply (simp add: nlab6-def)
apply (simp add: nlab7-def)
apply (simp add: nlab8-def)
apply (simp add: nMatch2-def)
apply (simp add: InstLabel1-def)
apply (simp add: InstLabel2-def)
apply (simp add: InstLabel3-def)
apply (simp add: InstLabel4-def)
apply (simp add: InstLabel5-def)
apply (simp add: InstLabel6-def)
apply (simp add: InstLabel7-def)
apply (simp add: InstLabel8-def)
apply (drule drop-length-append)
apply (simp add: Match2-def)
apply (subgoal-tac
  (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc
    (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc
    (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc
    (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc
    (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc
    (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc
    (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (the (cdm (l, i)
    ))))))))))))))))))))))))))))))))))))))))))))
  )))
  = (101 + the (cdm (l, i))))
by simp+

axioms equivS-MATCH:
  equivS S S' n ctm d g
  ⇒ equivS (map Val vsa @ S)
    (S' ++ map-of (zip [Suc (nat n)..<nat (n + nargs' + 1)]
      (map (λj. argCell fs-S'-n-minus-off (Suc j)) [nat nargs'>..0])))
    (n + nargs') ctm d g

axioms pc-MATCH:
  Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc
    (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc
    (Suc (Suc (Suc (the (cdm (l, i)) + nat labelEndLoop +
    nat (map (λn. trAddr n (the (cdm (l, i)) + incMatch)) (map (λp. the (cdm'
    (p, 0))) ps) !
    nat (the-Intg (the (fs-consTableS-S'-n-minus-off-tagGf (tagLf, consDataC))))))))))))))
  ))))))) =
  the (cdm (ps ! r, 0))

```

**axioms** *Obj-cells-equals*:

*Obj cellC fs-S'-n-minus-off* = *Obja cellC flds*  
 $\implies$  *flds* = *fs-S'-n-minus-off*

**lemma** *equivV-MATCH*:

$\llbracket$  *hm b* = *Some (j, C, vs)*;  
*equivH (hm, k) h (nat (the-Intg (the (sha (heapC, kf)))))) com regS d g*;  
*g b* = *Some l*;  
*h l* = *Some (Obj cellC fs-S'-n-minus-off)*  $\rrbracket$   
 $\implies \forall i < \text{length } vs. \text{equivV } (vs ! i) (\text{argCell } fs-S'-n-minus-off \ i) \ d \ g$

**apply** (*simp add: equivH.simps*)

**apply** *clarsimp*

**apply** (*erule-tac x=b in ballE*)

**apply** *clarsimp*

**apply** (*frule Obj-cells-equals, simp*)

**by** (*simp add: dom-def*)

**axioms** *equivS-extend*:

*equivS (map Val vs @ S)*  
*(S' ++ map-of (zip [Suc (nat n)..<nat (2 + (n + int (length vs))])*  
*(argCell fs-S'-n-minus-off (Suc (length vs)) #*  
*map ( $\lambda j. \text{argCell } fs-S'-n-minus-off (Suc j) [length vs > ..0]$ ))*  
*(n + int (length vs)) ctm d g*

**axioms** *argCell-Suc-j*:

*argCell fs-S'-n-minus-off 0* =  
*(argCell fs-S'-n-minus-off (Suc (length vs)) # map ( $\lambda j. \text{argCell } fs-S'-n-minus-off$*   
*(Suc j) [length vs > ..0]) !*  
*(nat (1 + (n + int (length vs))) - Suc (nat n))*

**lemma** *S'-map-of-0*:

*n*  $\geq$  *0*  
 $\implies (S' ++ \text{map-of } (\text{zip } [\text{Suc } (\text{nat } n) .. < \text{nat } (2 + (\text{n} + \text{int } (\text{length } \text{vs}))])$   
 $(\text{argCell } fs-S'-n-minus-off (\text{Suc } (\text{length } \text{vs})) \#$   
 $\text{map } (\lambda j. \text{argCell } fs-S'-n-minus-off (\text{Suc } j) [length$   
 $\text{vs} > ..0])))$   
 $(\text{nat } (\text{n} + (1 + \text{int } (\text{length } \text{vs})))) = \text{Some } (\text{argCell } fs-S'-n-minus-off$   
 $0)$

**apply** (*subst map-add-Some-iff*)

**apply** (*rule disjI1*)

**apply** (*rule map-of-is-SomeI*)

**apply** (*subst map-fst-zip*)

**apply** (*subst length-upt, simp*)

```

apply (subst length-upt-rev)
apply arith
apply simp
apply (subst set-zip)
apply simp
apply (rule-tac x=nat (1 + (n + int (length vs))) - Suc (nat n) in exI)
apply (rule conjI)
apply (subst nth-upt-2)
apply simp apply simp
apply (rule conjI)
apply (rule argCell-Suc-j)
apply (rule conjI)
apply simp
by (subst length-upt-rev,simp)

```

**axioms** equivV-HI:

```


$$\llbracket \forall i < \text{Suc } (\text{length } vs). \text{equivV } ((a \# vs) ! i) (\text{argCell } fs\text{-}S'\text{-}n\text{-minus-off } i) d g \rrbracket$$


$$\implies (\forall i < \text{length } vs. \text{equivV } (vs ! i) (\text{argCell } fs\text{-}S'\text{-}n\text{-minus-off } i) d g)$$


```

**lemma** equivS-MATCH [rule-format]:

```

equivS S S' n ctm d g
   $\longrightarrow n \geq 0$ 
   $\longrightarrow (\forall i < \text{length } vs. \text{equivV } (vs ! i) (\text{argCell } fs\text{-}S'\text{-}n\text{-minus-off } i) d g)$ 
   $\longrightarrow \text{equivS } (\text{map Val } vs @ S)$ 
    (S' ++ map-of (zip [Suc (nat n)..<nat (n + int (length vs) + 1)]
      (map ( $\lambda j. \text{argCell } fs\text{-}S'\text{-}n\text{-minus-off } (\text{Suc } j) [\text{length } vs > ..0]$ ))
      (n + int (length vs)) ctm d g)
apply (induct vs)
apply (simp add: equivS.simps)
apply clarsimp
apply (simp add: equivS.simps)
apply (rule conjI)
apply (subgoal-tac
  equivV a (argCell fs-S'-n-minus-off 0) d g)
apply (subgoal-tac
  the ((S' ++ map-of (zip [Suc (nat n)..<nat (2 + (n + int (length vs))])
    (argCell fs-S'-n-minus-off (Suc (length vs)) #
    map ( $\lambda j. \text{argCell } fs\text{-}S'\text{-}n\text{-minus-off } (\text{Suc } j) [\text{length } vs > ..0]$ ))
    (nat (n + (1 + int (length vs)))))) = argCell fs-S'-n-minus-off 0,simp)
apply (subst S'-map-of-0,assumption,simp)
apply (erule-tac x=0 in allE,simp)
apply (erule equivV-HI,simp)
by (rule equivS-extend)

```

**axioms** equivS-MATCH-axiom:

```

equivS (map Val vsa @ S)
  (S' ++ map-of (zip [Suc (nat n)..<nat (n + nargs' + 1)]
    (map ( $\lambda j. \text{argCell } fs\text{-}S'\text{-}n\text{-minus-off } (\text{Suc } j) [\text{nat } nargs' > ..0]$ ))

```

$(n + \text{nargs}') \text{ ctm } d \ g$

**declare** *equivH.simps* [*simp del*]

**lemma** *activeCells-MATCH*:

```

[[ la ∉ activeCells regS h (nat (the-Intg (the (sha (heapC, kf))))))]
  ⇒ activeCells regS h (nat (the-Intg (the (sha (heapC, kf)))) =
    activeCells regS
      (h(la ↦
        Arr ty m
          (S' ++ map-of (zip [Suc (nat n)..<nat (n + nargs' + 1)] (map (λj.
            argCell fs-S'-n-minus-off (Suc j)) [nat nargs'>..0])))))
        (nat (the-Intg (the (sha (heapC, kf))))))
apply (frule l-not-in-cellReg)
apply (unfold activeCells-def,auto)
apply (unfold region-def, simp add: Let-def, elim conjE)
apply (rule-tac x=j in exI,simp)
apply (rule cellReg-step)
apply (rule cellReg-basic)
apply (erule-tac x=j in allE)
apply (rule cellsReg-monotone-1,assumption+)
apply (erule l-not-in-regs)
apply (simp add: Let-def, elim conjE)
apply (rule-tac x=j in exI, simp)
apply (rule cellReg-step)
apply (rule cellReg-basic)
apply (erule-tac x=j in allE)
apply (erule cellsReg-monotone-2,assumption+)
by (erule l-not-in-regs)

```

**lemma** *domH-MATCH*:

```

[[ la ∉ activeCells regS h (nat (the-Intg (the (sha (heapC, kf)))));
  ∀ l ∈ dom H. ∃ l'. l' = the (g l) ∧ equivC (the (H l)) h l' (the (h l')) (nat
  (the-Intg (the (sha (heapC, kf)))))) com regS d g ]]
  ⇒ ∀ l ∈ dom H.
    ∃ l'. l' = the (g l) ∧
      equivC (the (H l))
        (h(la ↦
          Arr ty m
            (S' ++ map-of (zip [Suc (nat n)..<nat (n + nargs' + 1)]
              (map (λj. argCell fs-S'-n-minus-off (Suc j)) [nat
nargs'>..0])))))
        l' (the ((h(la ↦
          Arr ty m
            (S' ++ map-of (zip [Suc (nat n)..<nat (n + nargs' + 1)]

```

```

                                                                    (map (λj. argCell fs-S'-n-minus-off (Suc j))
[nat nargs'>..0])))
      l')
      (nat (the-Intg (the (sha (heapC, kf)))))) com regS d g
apply (frule l-not-in-cellReg)
apply (rule ballI)
apply (erule-tac x=l in ballE)
  prefer 2 apply simp
apply (erule exE)
apply (rule-tac x=l' in exI) apply (elim conjE)
  apply (rule conjI, assumption)
apply clarsimp
apply (rule conjI) apply (rule impI)
apply clarsimp apply (simp add: region-def)
apply clarsimp
apply (rule-tac x=Obj in exI)
apply (rule-tac x=flds in exI) apply simp
apply (simp add: region-def) apply (elim conjE)
apply (rule cellReg-step)
  apply (rule cellReg-basic)
  apply (erule-tac x=j in allE)
  apply (rule cellsReg-monotone-1, assumption+)
by (erule l-not-in-regs)

```

**lemma** *equivH-MATCH*:

```

[ k' = nat (the-Intg (the (sha (heapC, kf))));
  equivH (hm, k) h (nat (the-Intg (the (sha (heapC, kf)))))) com regS d g;
  sha (stackC, Sf) = Some (Addr la);
  activeCells regS h k' ∩ {la, l', l''} = {}]
⇒ equivH (hm, k)
  (h(la ↦
    Arr ty m
      (S' ++ map-of (zip [Suc (nat n)..<nat (n + nargs' + 1)] (map (λj.
argCell fs-S'-n-minus-off (Suc j)) [nat nargs'>..0])))
      (nat (the-Intg (the (sha (heapC, kf)))))) com regS d g
apply (simp add: heapC-def add: stackC-def)
apply (simp add: kf-def add: kf-def)
apply (fold heapC-def, fold kf-def)
apply (unfold equivH.simps, elim conjE)
apply (rule conjI, assumption)
apply (fold stackC-def)
apply (rule conjI)
apply (subgoal-tac
  la ∉ activeCells regS h (nat (the-Intg (the (sha (heapC, kf))))))
  apply (frule activeCells-MATCH, simp, simp)
apply (rule conjI, assumption)
by (rule domH-MATCH, assumption, simp)

```

**lemma** *execSVMInstr-MATCH* :

$$\begin{aligned} & \llbracket (P', cdm, ctm, com) = trSVM2JVM ((svms, ctmap), ini, ct, ah, ai, bc); \\ & \quad cdm, ctm, com \vdash ((hm, k), k0, (l, i), S) \triangleq S1'; \\ & \quad (fst (the (map-of svms l)) ! i) = MATCH \text{ off } ps; \\ & \quad execSVMInst (MATCH \text{ off } ps) (map-of ct) (hm, k) k0 (l, i) S = \\ & \quad \quad \text{Either.Right } S2; \\ & \quad drop (the (cdm (l, i))) (extractBytecode P') = \\ & \quad trInstr (the (cdm (l, i))) cdm' ctm' com pcc (MATCH \text{ off } ps) @ bytecode'; \\ & \quad cdm' \subseteq_m cdm \\ & \rrbracket \implies \exists v' sh' dh' ih' fms'. \\ & P' \vdash S1' \text{ --jvm} \rightarrow (v', sh', dh', ih', fms') \wedge \\ & cdm, ctm, com \vdash S2 \triangleq (v', sh', dh', ih', fms') \end{aligned}$$

**apply** (*case-tac*  $S1'$ )  
**apply** (*rename-tac*  $v1$  *tup*)  
**apply** (*case-tac* *tup*)  
**apply** (*rename-tac*  $sh$  *tup*)  
**apply** (*case-tac* *tup*)  
**apply** (*rename-tac*  $dh$  *tup*)  
**apply** (*case-tac* *tup*)  
**apply** (*rename-tac*  $ih$  *fms*)  
**apply** (*simp del: execSVMInst.simps*)  
**apply** (*unfold equivState-def*)  
**apply** (*elim exE, elim conjE*)  
**apply** (*subgoal-tac int off  $\leq n$* )  
**prefer** 2 **apply** (*erule equivS-MATCH-bound, simp*)  
**apply** *clarify*

**apply** (*simp only: execSVMInst.simps*)  
**apply** (*insert RightNotUndefined*)  
**apply** (*erule-tac x=S2 in allE*)  
**apply** (*case-tac S!+ off*)  
**defer** **apply** (*force, force*)  
**apply** (*case-tac Vala*)  
**defer** **apply** (*force, force*)  
**apply** (*case-tac hm nata*)  
**apply** *force*  
**apply** (*case-tac aha, rename-tac j rest, case-tac rest, rename-tac C usa*)  
**apply** (*case-tac map-of ct C*)  
**apply** *force*  
**apply** (*rename-tac tup, case-tac tup,*  
*rename-tac nargs rest', case-tac rest', rename-tac r xa*)  
**apply** *clarsimp*

**apply** (*frule equivS-S'-n-minus-off, assumption+*)

```

apply (elim exE, elim conjE)
apply (subgoal-tac
  sha (consTableC, tablef) = Some (Addr l''') ∧
  h l''' = Some (Arr ty m''' consTableS))
prefer 2 apply (rule consTableC-def)
apply (erule conjE)
apply (subgoal-tac
  ∃ l' fs-consTableS-S'-n-minus-off-tagGf.
  the (consTableS (nat (the-Intg (the (fs-S'-n-minus-off (tagGf, cellC)))))) =
  Addr l' ∧
  h l' = Some (Obj cellC fs-consTableS-S'-n-minus-off-tagGf) ∧
  the (fs-consTableS-S'-n-minus-off-tagGf (nargsf, consDataC)) = Intg nargs'
  ∧
  nargs' ≥ 0)
apply (elim exE, elim conjE)
prefer 2
apply (rule-tac l'''=l''' in consTableC-n-minus-off,assumption+)

apply (rule-tac x=sha((stackC, topf) ↦ Intg (n + nargs')) in exI)
apply (rule-tac x=h(la ↦ Arr ty m (S' ++ map-of (zip [Suc (nat n)..<nat (n +
nargs' + 1)]
  (map (λj. argCell fs-S'-n-minus-off (Suc j))
  [nat nargs'>..0]))) in exI)

apply (rule-tac x=inih in exI)
apply (rule-tac x=[([],vs[Suc 0 := the (S' (nat (n - int off))),
  6 := Addr laa,
  2 := the (fs-S'-n-minus-off (tagGf, cellC)),
  3 := the (fs-consTableS-S'-n-minus-off-tagGf (tagLf,
consDataC)),
  4 := Intg nargs',
  5 := Intg nargs'],
safeP, sigSafeMain,
  Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc
(Suc
  (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc
(Suc
  (Suc (Suc (Suc (Suc (the (cdm (l, i)) + nat labelEndLoop +
  nat (map (λn. trAddr n (the (cdm (l, i)) + incMatch))
  (map (λp. the (cdm' (p, 0))) ps) !
  nat (the-Intg (the (fs-consTableS-S'-n-minus-off-tagGf (tagLf,
consDataC))
  )))))))))))))))))))))))))] in exI)

apply (subgoal-tac length vs = 10)
prefer 2 apply (rule length-vs)
apply (rule conjI)

```

```

apply (subgoal-tac
  P⊢ (None, sha, h, inih, [([]),
    vs, safeP, sigSafeMain, the (cdm (l, i)), ag, bd)])
  -jvm→
    (None, sha, h, inih, [(Addr la],
      vs, safeP, sigSafeMain, the (cdm (l, i)) + 1, ag, bd)))
apply (unfold exec-all-def)
apply (erule rtrancl-trans)

prefer 2
apply (rule r-into-rtrancl)
apply clarify
apply (unfold JVMEexec.exec.simps)
apply (unfold Let-def)
apply (unfold extractBytecode-def)
apply (subgoal-tac (fst (snd (snd (snd (snd (the (method' (P', safeP) sigSafe-
  Main))))))) !
  the (cdm (l, i))) = Getstatic Sf stackC) apply simp
apply (unfold trInstr.simps, unfold Let-def)
apply (unfold Match11-def)
apply (rule nth-via-drop-append) apply force

apply (drule-tac ?ms=bytecode' and ?y= Getstatic Sf stackC
  in drop-Suc-append-2)
apply (subgoal-tac
  P⊢ (None, sha, h, inih,
    [(Addr la],
      vs, safeP, sigSafeMain, the (cdm (l, i)) + 1, ag, bd)])
  -jvm→
    (None, sha, h, inih,
      [(Intg n, Addr la],
        vs, safeP, sigSafeMain, the (cdm (l, i)) + 1 + 1, ag, bd)))
apply (unfold exec-all-def)
apply (simp del:append.append-Cons)
apply (erule rtrancl-trans)
prefer 2
apply (rule r-into-rtrancl)
apply (clarify)
apply (unfold JVMEexec.exec.simps)
apply (unfold Let-def)
apply (subgoal-tac (fst (snd (snd (snd (snd (the (method' (P', safeP) sigSafe-
  Main)))))))
  ! Suc (the (cdm (l, i))) = Getstatic topf stackC, simp)
apply (erule nth-via-drop-append)

apply (drule drop-Suc-append)

```



```

apply (subgoal-tac
  P⊢ (None,sha,h,inih,
    [[Intg n,Addr la],
     vs, safeP, sigSafeMain, the (cdm (l, i)) + 1 + 1, ag, bd]])
  -jvm→
    (None,sha,h,inih,
     [[Intg (int off),
      Intg n,Addr la],
      vs, safeP, sigSafeMain, the (cdm (l, i)) + 1 + 1 + 1, ag, bd]])
apply (unfold exec-all-def)
apply (simp del:append.append-Cons)
apply (erule rtrancl-trans)
prefer 2
apply (rule r-into-rtrancl)
apply (clarify)
apply (unfold JVMExec.exec.simps)
apply (unfold Let-def)
apply (subgoal-tac (fst (snd (snd (snd (snd (the (method' (P', safeP) sigSafe-
Main)))))))
  ! (Suc (Suc (the (cdm (l, i)))))) = LitPush (Intg (int off)),simp)
apply (erule nth-via-drop-append)

```

```

apply (drule drop-Suc-append)
apply (subgoal-tac P⊢ (None,sha,h,inih,
  [[Intg (int off),
   Intg n,Addr la],
   vs, safeP, sigSafeMain, the (cdm (l, i)) + 1 + 1 + 1, ag, bd]])
  -jvm→
    (None,sha,h,inih,
     [[Intg ( n - int off),
      Addr la],
      vs, safeP, sigSafeMain, the (cdm (l, i)) + 1 + 1 + 1 + 1,
      ag, bd]])
apply (unfold exec-all-def)
apply (simp del:append.append-Cons)
apply (erule rtrancl-trans)
prefer 2
apply (rule r-into-rtrancl)
apply (clarify)
apply (unfold JVMExec.exec.simps)
apply (unfold Let-def)
apply (subgoal-tac (fst (snd (snd (snd (snd (the (method' (P', safeP) sigSafe-
Main)))))))
  ! (Suc (Suc (Suc (the (cdm (l, i))))))) = BinOp Subtract,simp)
apply (erule nth-via-drop-append)

```

```

apply (drule drop-Suc-append)

```

```

apply (subgoal-tac P⊢ (None,sha,h,inih,
  [[Intg (n - int off),
   Addr la],
   vs, safeP, sigSafeMain, the (cdm (l, i)) + 1 + 1 + 1 + 1,
  ag, bd]))
  -jvm→
  (None,sha,h,inih,
  [[the (S' (nat (n - int off))),
   vs, safeP, sigSafeMain, the (cdm (l, i)) + 1 + 1 + 1 + 1 +
  1, ag, bd]]))
apply (unfold exec-all-def)
apply (simp del:append.append-Cons)
apply (erule rtrancl-trans)
prefer 2
apply (rule r-into-rtrancl)
apply (clarify)
apply (unfold JVMEexec.exec.simps)
apply (unfold Let-def)
apply (subgoal-tac (fst (snd (snd (snd (snd (the (method' (P', safeP) sigSafe-
Main))))))) !
  (Suc (Suc (Suc (Suc (the (cdm (l, i))))))) = ArrLoad)
apply (simp del: Let-def execSVMInst.simps)
apply (simp (no-asm) add: raise-system-xcpt-def)
apply (rule conjI) apply (force,force)
apply (erule nth-via-drop-append)

apply (drule drop-Suc-append)
apply (subgoal-tac
  P⊢ (None,sha,h,inih,[[the (S' (nat (n - int off))),
   vs, safeP, sigSafeMain, the (cdm (l, i)) + 1 + 1 + 1 + 1 + 1, ag, bd]])
  -jvm→
  (None,sha,h,inih,[[[],vs[Suc 0 := the (S' (nat (n - int off))),
   safeP, sigSafeMain, the (cdm (l, i)) + 1 + 1 + 1 + 1 + 1 + 1, ag, bd]]))
apply (unfold exec-all-def)
apply (simp del:append.append-Cons)
apply (erule rtrancl-trans)
prefer 2
apply (rule r-into-rtrancl)
apply (clarify)
apply (unfold JVMEexec.exec.simps)
apply (unfold Let-def)
apply (subgoal-tac (fst (snd (snd (snd (snd (the (method' (P', safeP) sigSafe-
Main))))))) !
  (Suc (Suc (Suc (Suc (Suc (the (cdm (l, i))))))) =
  Store (Suc 0)) apply (simp)
apply (erule nth-via-drop-append)

```

```

apply (drule drop-Suc-append)
apply (subgoal-tac P⊢ (None, sha, h, inih,
  [[[], vs[Suc 0 := the (S' (nat (n - int off))]], safeP, sigSafeMain,
  the (cdm (l, i)) + 1 + 1 + 1 + 1 + 1 + 1, ag, bd]])
  -jvm→
  (None, sha, h, inih,
  [[[Addr l'], vs[Suc 0 := the (S' (nat (n - int off))]],
  safeP, sigSafeMain, the (cdm (l, i)) + 1 + 1 + 1 + 1 + 1 + 1 + 1, ag,
  bd]]])
apply (unfold exec-all-def)
apply (simp del:append.append-Cons)
apply (erule rtrancl-trans)
prefer 2
apply (rule r-into-rtrancl)
apply (clarify)
apply (unfold JVMExec.exec.simps)
apply (unfold Let-def)
apply (subgoal-tac (fst (snd (snd (snd (snd (the (method' (P', safeP) sigSafe-
Main)))))) !
  (Suc (Suc (Suc (Suc (Suc (Suc (the (cdm (l, i)))))))) =
  Getstatic safeDirf dirCellC, simp)
apply (erule nth-via-drop-append)

```

```

apply (drule drop-Suc-append)
apply (subgoal-tac P⊢ (None, sha, h, inih,
  [[[Addr l'], vs[Suc 0 := the (S' (nat (n - int off))]],
  safeP, sigSafeMain, the (cdm (l, i)) + 1 + 1 + 1 + 1 + 1 + 1 + 1, ag,
  bd]])
  -jvm→
  (None, sha, h, inih,
  [[[the (S' (nat (n - int off))), Addr l'],
  vs[Suc 0 := the (S' (nat (n - int off))]],
  safeP, sigSafeMain, the (cdm (l, i)) + 1 + 1 + 1 + 1 + 1 + 1 + 1 +
  1, ag, bd]]])
apply (unfold exec-all-def)
apply (simp del: append.append-Cons)
apply (erule rtrancl-trans)
prefer 2
apply (rule r-into-rtrancl)
apply (clarify)
apply (unfold JVMExec.exec.simps)
apply (unfold Let-def)
apply (subgoal-tac fst (snd (snd (snd (snd (the (method' (P', safeP) sigSafe-
Main)))))) !
  (Suc (Suc (Suc (Suc (Suc (Suc (Suc (the (cdm (l, i)))))))) =
  Load (Suc 0))) apply (simp (no-asm-simp))
apply (erule nth-via-drop-append)

```

**apply** (*drule drop-Suc-append*)

**apply** (*subgoal-tac*  $P \vdash$  (*None, sha, h, inih*,  
 [[*the* ( $S'$  (*nat* ( $n - \text{int off}$ ))), *Addr l'*],  
*vs*[*Suc 0 := the* ( $S'$  (*nat* ( $n - \text{int off}$ ))),  
*safeP*, *sigSafeMain*, *the* (*cdm* ( $l, i$ )) + 1 + 1 + 1 + 1 + 1 + 1 + 1 +  
 1, *ag*, *bd*]])  
 $\rightarrow$   
 (*None, sha, h, inih*,  
 [[*Addr laa*],  
*vs*[*Suc 0 := the* ( $S'$  (*nat* ( $n - \text{int off}$ ))),  
*safeP*, *sigSafeMain*, *the* (*cdm* ( $l, i$ )) + 1 + 1 + 1 + 1 + 1 + 1 + 1 +  
 + 1, *ag*, *bd*]]))  
**apply** (*unfold exec-all-def*)  
**apply** (*simp del: append.append-Cons*)  
**apply** (*erule rtrancl-trans*)  
**prefer 2**  
**apply** (*rule r-into-rtrancl*)  
**apply** (*clarify*)  
**apply** (*unfold JVMExec.exec.simps*)  
**apply** (*unfold Let-def*)  
**apply** (*subgoal-tac fst* (*snd* (*snd* (*snd* (*snd* (*the* (*method'* ( $P'$ , *safeP*) *sigSafe-*  
*Main*)))))) !  
*Suc* (*Suc* (*Suc* (*Suc* (*Suc* (*Suc* (*Suc* (*Suc* (*Suc* (*the* (*cdm* ( $l, i$ )))))))))))  
 =  
*ArrLoad*) **apply** (*simp* (*no-asm-simp*) *add: raise-system-xcpt-def*)  
**apply** (*subgoal-tac the-Intg* (*the* ( $S'$  (*nat* ( $n - \text{int off}$ ))) < *int m''*  
 $\wedge 0 \leq \text{the-Intg}$  (*the* ( $S'$  (*nat* ( $n - \text{int off}$ ))))))  
**prefer 2** **apply** (*rule safeDir-bounds*)  
**apply** (*rule conjI*) **apply** (*force, force*)  
**apply** (*erule nth-via-drop-append*)

**apply** (*drule drop-Suc-append*)

**apply** (*subgoal-tac*  $P \vdash$  (*None, sha, h, inih*,  
 [[*Addr laa*],  
*vs*[*Suc 0 := the* ( $S'$  (*nat* ( $n - \text{int off}$ ))),  
*safeP*, *sigSafeMain*, *the* (*cdm* ( $l, i$ )) + 1 + 1 + 1 + 1 + 1 + 1 + 1 +  
 + 1, *ag*, *bd*]])  
 $\rightarrow$   
 (*None, sha, h, inih*,  
 [[[],  
*vs*[*Suc 0 := the* ( $S'$  (*nat* ( $n - \text{int off}$ ))),  
   *6 := Addr laa*,  
*safeP*, *sigSafeMain*, *the* (*cdm* ( $l, i$ )) + 1 + 1 + 1 + 1 + 1 + 1 + 1 +  
 + 1 + 1, *ag*, *bd*]]))  
**apply** (*unfold exec-all-def*)  
**apply** (*simp del: append.append-Cons*)

```

apply (erule rtrancl-trans)
prefer 2
apply (rule r-into-rtrancl)
apply (clarify)
apply (unfold JVMEexec.exec.simps)
apply (unfold Let-def)
apply (subgoal-tac fst (snd (snd (snd (snd (the (method' (P', safeP) sigSafe-
Main)))))) !
      (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (the (cdm (l,
i)))))))))) =
      Store 6) apply (simp (no-asm-simp))
apply (erule nth-via-drop-append)

```

```

apply (drule drop-Suc-append)
apply (subgoal-tac P+ (None, sha, h, inih,
  [([],
    vs[Suc 0 := the (S' (nat (n - int off))),
      6 := Addr laa],
    safeP, sigSafeMain, the (cdm (l, i)) + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1
+ 1 + 1, ag, bd)])
  -jvm→
  (None, sha, h, inih,
  [([Addr laa],
    vs[Suc 0 := the (S' (nat (n - int off))),
      6 := Addr laa],
    safeP, sigSafeMain, the (cdm (l, i)) + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1
+ 1 + 1 + 1, ag, bd)]))
apply (unfold exec-all-def)
apply (simp del: append.append-Cons)
apply (erule rtrancl-trans)
prefer 2
apply (rule r-into-rtrancl)
apply (clarify)
apply (unfold JVMEexec.exec.simps)
apply (unfold Let-def)
apply (subgoal-tac fst (snd (snd (snd (snd (the (method' (P', safeP) sigSafe-
Main)))))) !
      (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (the (cdm
(l, i)))))))))) =
      Load 6) apply (simp (no-asm-simp))
apply (erule nth-via-drop-append)

```

```

apply (drule drop-Suc-append)
apply (subgoal-tac P+ (None, sha, h, inih,
  [([Addr laa],
    vs[Suc 0 := the (S' (nat (n - int off))),
      6 := Addr laa],

```

```

      safeP, sigSafeMain, the (cdm (l, i)) + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1
+ 1 + 1 + 1 + 1, ag, bd)])
    -jvm→
      (None, sha, h, inih,
      [[the (fs-S'-n-minus-off (tagGf, cellC))],
      vs[Suc 0 := the (S' (nat (n - int off))),
      6 := Addr laa],
      safeP, sigSafeMain, the (cdm (l, i)) + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1
+ 1 + 1 + 1 + 1, ag, bd)])
apply (unfold exec-all-def)
apply (simp del: append.append-Cons)
apply (erule rtrancl-trans)
prefer 2
apply (rule r-into-rtrancl)
apply (clarify)
apply (unfold JVMEexec.exec.simps)
apply (unfold Let-def)
apply (subgoal-tac fst (snd (snd (snd (snd (the (method' (P', safeP) sigSafe-
Main)))))) !
      (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (the
(cdm (l, i)))))))))))))) =
      Getfield tagGf cellC) apply (simp (no-asm-simp) add:
raise-system-xcpt-def)
apply (simp (no-asm-simp))
apply (erule nth-via-drop-append)

apply (drule drop-Suc-append)
apply (subgoal-tac P ⊢ (None, sha, h, inih,
      [[the (fs-S'-n-minus-off (tagGf, cellC))],
      vs[Suc 0 := the (S' (nat (n - int off))),
      6 := Addr laa],
      safeP, sigSafeMain, the (cdm (l, i)) + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1
+ 1 + 1 + 1 + 1, ag, bd)])
    -jvm→
      (None, sha, h, inih,
      [[[],
      vs[Suc 0 := the (S' (nat (n - int off))),
      6 := Addr laa,
      2 := the (fs-S'-n-minus-off (tagGf, cellC))],
      safeP, sigSafeMain, the (cdm (l, i)) + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1
+ 1 + 1 + 1 + 1 + 1, ag, bd)])
apply (unfold exec-all-def)
apply (simp del: append.append-Cons)
apply (erule rtrancl-trans)
prefer 2
apply (rule r-into-rtrancl)
apply (clarify)
apply (unfold JVMEexec.exec.simps)

```

**apply** (*unfold Let-def*)  
**apply** (*subgoal-tac fst (snd (snd (snd (snd (the (method' (P', safeP) sigSafeMain)))))) !*  
*(Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (the (cdm (l, i)))))))))))))) =*  
*Store 2)* **apply** (*simp (no-asm-simp)*)  
**apply** (*erule nth-via-drop-append*)

**apply** (*drule drop-Suc-append*)  
**apply** (*subgoal-tac*  
*sha (consTableC, tablef) = Some (Addr l''') ∧*  
*h l''' = Some (Arr ty m''' consTableS)*)  
**apply** (*elim conjE*)  
**prefer** 2 **apply** (*rule consTableC-def*)  
**apply** (*subgoal-tac P<sup>⊢</sup> (None, sha, h, inih,*  
*[[[],*  
*vs[Suc 0 := the (S' (nat (n - int off))),*  
*6 := Addr laa,*  
*2 := the (fs-S'-n-minus-off (tagGf, cellC))],*  
*safeP, sigSafeMain, the (cdm (l, i)) + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1*  
*+ 1 + 1 + 1 + 1 + 1, ag, bd]]*)  
*-jvm→*  
*(None, sha, h, inih,*  
*[[[Addr l'''],*  
*vs[Suc 0 := the (S' (nat (n - int off))),*  
*6 := Addr laa,*  
*2 := the (fs-S'-n-minus-off (tagGf, cellC))],*  
*safeP, sigSafeMain, the (cdm (l, i)) + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1*  
*+ 1 + 1 + 1 + 1 + 1 + 1 + 1, ag, bd]]*)  
**apply** (*unfold exec-all-def*)  
**apply** (*simp del: append.append-Cons*)  
**apply** (*erule rtrancl-trans*)  
**prefer** 2  
**apply** (*rule r-into-rtrancl*)  
**apply** (*clarify*)  
**apply** (*unfold JVMEexec.exec.simps*)  
**apply** (*unfold Let-def*)  
**apply** (*subgoal-tac fst (snd (snd (snd (snd (the (method' (P', safeP) sigSafeMain)))))) !*  
*(Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (the (cdm (l, i)))))))))))))))) =*  
*Getstatic tablef consTableC)* **apply** (*simp (no-asm-simp)*)  
**apply** (*erule nth-via-drop-append*)

**apply** (*drule drop-Suc-append*)  
**apply** (*subgoal-tac P<sup>⊢</sup> (None, sha, h, inih,*  
*[[[Addr l'''],*

```

vs[Suc 0 := the (S' (nat (n - int off))),
  6 := Addr laa,
  2 := the (fs-S'-n-minus-off (tagGf, cellC))],
safeP, sigSafeMain, the (cdm (l, i) + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1
+ 1 + 1 + 1 + 1 + 1 + 1, ag, bd)])
-jvm→
(None, sha, h, inih,
  [(the (fs-S'-n-minus-off (tagGf, cellC)),
    Addr l''),
  vs[Suc 0 := the (S' (nat (n - int off))),
    6 := Addr laa,
    2 := the (fs-S'-n-minus-off (tagGf, cellC))],
  safeP, sigSafeMain, the (cdm (l, i) + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1
+ 1 + 1 + 1 + 1 + 1 + 1, ag, bd))])
apply (unfold exec-all-def)
apply (simp del: append.append-Cons)
apply (erule rtrancl-trans)
prefer 2
apply (rule r-into-rtrancl)
apply (clarify)
apply (unfold JVMExec.exec.simps)
apply (unfold Let-def)
apply (subgoal-tac fst (snd (snd (snd (snd (the (method' (P', safeP) sigSafe-
Main)))))) !
  (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc
(Suc (Suc (the (cdm (l, i)))))))))))))))))) =
  Load 2) apply (simp (no-asm-simp))
apply (erule nth-via-drop-append)

apply (drule drop-Suc-append)
apply (subgoal-tac P⊢ (None, sha, h, inih,
  [(the (fs-S'-n-minus-off (tagGf, cellC)),
    Addr l''),
  vs[Suc 0 := the (S' (nat (n - int off))),
    6 := Addr laa,
    2 := the (fs-S'-n-minus-off (tagGf, cellC))],
  safeP, sigSafeMain, the (cdm (l, i) + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1
+ 1 + 1 + 1 + 1 + 1 + 1, ag, bd))])
-jvm→
(None, sha, h, inih,
  [(Addr l'a],
  vs[Suc 0 := the (S' (nat (n - int off))),
    6 := Addr laa,
    2 := the (fs-S'-n-minus-off (tagGf, cellC))],
  safeP, sigSafeMain, the (cdm (l, i) + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1
+ 1 + 1 + 1 + 1 + 1 + 1, ag, bd))])
apply (unfold exec-all-def)
apply (simp del: append.append-Cons)

```



```

apply (erule rtrancl-trans)
prefer 2
apply (rule r-into-rtrancl)
apply (clarify)
apply (unfold JVMExec.exec.simps)
apply (unfold Let-def)
apply (subgoal-tac fst (snd (snd (snd (snd (the (method' (P', safeP) sigSafe-
Main)))))) !
      (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc
(Suc (Suc (Suc (the (cdm (l, i)))))))))))))))))) =
      ArrLoad) apply (simp (no-asm-simp) add: raise-system-xcpt-def)
apply (subgoal-tac
  int m''' > the-Intg (the (fs-S'-n-minus-off (tagGf, cellC))) ∧
  the-Intg (the (fs-S'-n-minus-off (tagGf, cellC))) ≥ 0)
apply (simp (no-asm-simp))
apply simp
apply (rule consTable-bounds)
apply (erule nth-via-drop-append)

apply (drule drop-Suc-append)
apply (subgoal-tac P ⊢ (None, sha, h, inih,
  [[Addr l'a],
  vs[Suc 0 := the (S' (nat (n - int off))),
  6 := Addr laa,
  2 := the (fs-S'-n-minus-off (tagGf, cellC))],
  safeP, sigSafeMain, the (cdm (l, i)) + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1
+ 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1, ag, bd]))
  -jvm →
  (None, sha, h, inih,
  [[Addr l'a,
  Addr l'a],
  vs[Suc 0 := the (S' (nat (n - int off))),
  6 := Addr laa,
  2 := the (fs-S'-n-minus-off (tagGf, cellC))],
  safeP, sigSafeMain, the (cdm (l, i)) + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1
+ 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1, ag, bd]))
apply (unfold exec-all-def)
apply (simp del: append.append-Cons)
apply (erule rtrancl-trans)
prefer 2
apply (rule r-into-rtrancl)
apply (clarify)
apply (unfold JVMExec.exec.simps)
apply (unfold Let-def)
apply (subgoal-tac fst (snd (snd (snd (snd (the (method' (P', safeP) sigSafe-
Main)))))) !
      (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc
(Suc (Suc (Suc (Suc (the (cdm (l, i)))))))))))))))))) =

```



```

-jvm→
  (None,sha,h,inih,
  [[Addr l'a],
  vs[Suc 0 := the (S' (nat (n - int off))),
    6 := Addr laa,
    2 := the (fs-S'-n-minus-off (tagGf, cellC)),
    3 := the (fs-consTableS-S'-n-minus-off-tagGf (tagLf, consDataC))],
  safeP, sigSafeMain,
  the (cdm (l, i)) + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1
+ 1 + 1 + 1 + 1 + 1 + 1 + 1,
  ag, bd]])
apply (unfold exec-all-def)
apply (simp del: append.append-Cons)
apply (erule rtrancl-trans)
prefer 2
apply (rule r-into-rtrancl)
apply (clarify)
apply (unfold JVMExec.exec.simps)
apply (unfold Let-def)
apply (subgoal-tac fst (snd (snd (snd (snd (the (method' (P', safeP) sigSafe-
Main)))))) !
  (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc
(Suc (Suc (Suc (Suc
  (Suc (Suc (the (cdm (l, i)))))))))))))))))) =
  Store 3) apply (simp (no-asm-simp))
apply (erule nth-via-drop-append)

apply (drule drop-Suc-append)
apply (subgoal-tac P⊢ (None,sha,h,inih,
  [[Addr l'a],
  vs[Suc 0 := the (S' (nat (n - int off))),
    6 := Addr laa,
    2 := the (fs-S'-n-minus-off (tagGf, cellC)),
    3 := the (fs-consTableS-S'-n-minus-off-tagGf (tagLf, consDataC))],
  safeP, sigSafeMain,
  the (cdm (l, i)) + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1
+ 1 + 1 + 1 + 1 + 1 + 1 + 1,
  ag, bd]])
-jvm→
  (None,sha,h,inih,
  [[Intg nargs'],
  vs[Suc 0 := the (S' (nat (n - int off))),
    6 := Addr laa,
    2 := the (fs-S'-n-minus-off (tagGf, cellC)),
    3 := the (fs-consTableS-S'-n-minus-off-tagGf (tagLf, consDataC))],
  safeP, sigSafeMain,
  the (cdm (l, i)) + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1
+ 1 + 1 + 1 + 1 + 1 + 1 + 1,

```

```

    ag, bd)))
apply (unfold exec-all-def)
apply (simp del: append.append-Cons)
apply (erule rtrancl-trans)
prefer 2
apply (rule r-into-rtrancl)
apply (clarify)
apply (unfold JVMEexec.exec.simps)
apply (unfold Let-def)
apply (subgoal-tac fst (snd (snd (snd (snd (the (method' (P', safeP) sigSafeMain)))))) !
      (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc
(Suc (Suc (Suc (Suc
      (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc
      (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc
      (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc
      Getfield nargsf consDataC) apply (simp (no-asm-simp) add:
raise-system-xcpt-def)
apply (erule nth-via-drop-append)

apply (erule drop-Suc-append)
apply (subgoal-tac P ⊢ (None, sha, h, inih,
  [([Intg nargs'],
    vs[Suc 0 := the (S' (nat (n - int off))),
      6 := Addr laa,
      2 := the (fs-S'-n-minus-off (tagGf, cellC)),
      3 := the (fs-consTableS-S'-n-minus-off-tagGf (tagLf, consDataC))],
    safeP, sigSafeMain,
    the (cdm (l, i)) + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1,
    ag, bd]))
  -jvm→
  (None, sha, h, inih,
  ([],
    vs[Suc 0 := the (S' (nat (n - int off))),
      6 := Addr laa,
      2 := the (fs-S'-n-minus-off (tagGf, cellC)),
      3 := the (fs-consTableS-S'-n-minus-off-tagGf (tagLf, consDataC)),
      4 := the (fs-consTableS-S'-n-minus-off-tagGf (nargsf, consDataC))],
    safeP, sigSafeMain,
    the (cdm (l, i)) + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1,
    ag, bd]))))
apply (unfold exec-all-def)
apply (simp del: append.append-Cons)
apply (erule rtrancl-trans)
prefer 2
apply (rule r-into-rtrancl)
apply (clarify)
apply (unfold JVMEexec.exec.simps)

```

```

apply (unfold Let-def)
apply (subgoal-tac fst (snd (snd (snd (snd (the (method' (P', safeP) sigSafe-
Main)))))) !
      (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc
(Suc (Suc (Suc (Suc
      (Suc (Suc (Suc (Suc (the (cdm (l, i))))))))))))))))) =
      Store 4) apply (simp (no-asm-simp))
apply (erule nth-via-drop-append)

apply (drule drop-Suc-append)
apply (subgoal-tac P+ (None, sha, h, inih,
  [([],
    vs[Suc 0 := the (S' (nat (n - int off))),
      6 := Addr laa,
      2 := the (fs-S'-n-minus-off (tagGf, cellC)),
      3 := the (fs-consTableS-S'-n-minus-off-tagGf (tagLf, consDataC)),
      4 := Intg nargs^],
    safeP, sigSafeMain,
    the (cdm (l, i)) + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1
+ 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1,
    ag, bd]))
  -jvm→
  (None, sha, h, inih,
  [([Intg 0],
    vs[Suc 0 := the (S' (nat (n - int off))),
      6 := Addr laa,
      2 := the (fs-S'-n-minus-off (tagGf, cellC)),
      3 := the (fs-consTableS-S'-n-minus-off-tagGf (tagLf, consDataC)),
      4 := Intg nargs^],
    safeP, sigSafeMain,
    the (cdm (l, i)) + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1
+ 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1,
    ag, bd]]))
apply (unfold exec-all-def)
apply (simp del: append.append-Cons)
apply (erule rtrancl-trans)
prefer 2
apply (rule r-into-rtrancl)
apply (clarify)
apply (unfold JVMExec.exec.simps)
apply (unfold Let-def)
apply (subgoal-tac fst (snd (snd (snd (snd (the (method' (P', safeP) sigSafe-
Main)))))) !
      (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc
(Suc (Suc (Suc (Suc
      (Suc (Suc (Suc (Suc (the (cdm (l, i))))))))))))))))) =
      LitPush (Intg 0)) apply (simp (no-asm-simp))
apply (erule nth-via-drop-append)

```



```

        6 := Addr laa,
        2 := the (fs-S'-n-minus-off (tagGf, cellC)),
        3 := the (fs-consTableS-S'-n-minus-off-tagGf (tagLf, consDataC)),
        4 := Intg nargs',
        5 := Intg 0],
    safeP, sigSafeMain,
    the (cdm (l, i)) + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1
+ 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1,
    ag, bd)])
-jvm→
(None, sha, h, inih,
 [([Intg n],
 vs[Suc 0 := the (S' (nat (n - int off))),
    6 := Addr laa,
    2 := the (fs-S'-n-minus-off (tagGf, cellC)),
    3 := the (fs-consTableS-S'-n-minus-off-tagGf (tagLf, consDataC)),
    4 := Intg nargs',
    5 := Intg 0],
    safeP, sigSafeMain,
    the (cdm (l, i)) + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1
+ 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1,
    ag, bd]]))
apply (unfold exec-all-def)
apply (simp del: append.append-Cons)
apply (erule rtrancl-trans)
prefer 2
apply (rule r-into-rtrancl)
apply (clarify)
apply (unfold JVMEexec.exec.simps)
apply (unfold Let-def)
apply (subgoal-tac fst (snd (snd (snd (snd (the (method' (P', safeP) sigSafe-
Main)))))) !
    (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc
(Suc (Suc (Suc (Suc
    (Suc (Suc (Suc (Suc (Suc (Suc (Suc (the (cdm (l, i))))))))))))))))))))))
=
    Getstatic topf stackC) apply (simp (no-asm-simp))
apply (erule nth-via-drop-append)

apply (drule drop-Suc-append)
apply (subgoal-tac P+ (None, sha, h, inih,
 [([Intg n],
 vs[Suc 0 := the (S' (nat (n - int off))),
    6 := Addr laa,
    2 := the (fs-S'-n-minus-off (tagGf, cellC)),
    3 := the (fs-consTableS-S'-n-minus-off-tagGf (tagLf, consDataC)),
    4 := Intg nargs',
    5 := Intg 0],

```

```

    safeP, sigSafeMain,
    the (cdm (l, i)) + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1
+ 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1,
    ag, bd)])
-jvm→
(None, sha, h, inih,
 [([Intg nargs',
   Intg n],
  vs[Suc 0 := the (S' (nat (n - int off))),
     6 := Addr laa,
     2 := the (fs-S'-n-minus-off (tagGf, cellC)),
     3 := the (fs-consTableS-S'-n-minus-off-tagGf (tagLf, consDataC)),
     4 := Intg nargs',
     5 := Intg 0],
   safeP, sigSafeMain,
   the (cdm (l, i)) + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1
+ 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1,
   ag, bd]]))
apply (unfold exec-all-def)
apply (simp del: append.append-Cons)
apply (erule rtrancl-trans)
prefer 2
apply (rule r-into-rtrancl)
apply (clarify)
apply (unfold JVMExec.exec.simps)
apply (unfold Let-def)
apply (subgoal-tac fst (snd (snd (snd (snd (the (method' (P', safeP) sigSafe-
Main)))))) !
(Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc
(Suc (Suc (Suc (Suc
(Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (the (cdm (l,
i)))))))))))))))))) =
Load 4) apply (simp (no-asm-simp))
apply (erule nth-via-drop-append)

apply (drule drop-Suc-append)
apply (subgoal-tac P'⊢ (None, sha, h, inih,
 [([Intg nargs',
   Intg n],
  vs[Suc 0 := the (S' (nat (n - int off))),
     6 := Addr laa,
     2 := the (fs-S'-n-minus-off (tagGf, cellC)),
     3 := the (fs-consTableS-S'-n-minus-off-tagGf (tagLf, consDataC)),
     4 := Intg nargs',
     5 := Intg 0],
   safeP, sigSafeMain,
   the (cdm (l, i)) + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1
+ 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1,

```





```

[[[],
  vs[Suc 0 := the (S' (nat (n - int off))),
    6 := Addr laa,
    2 := the (fs-S'-n-minus-off (tagGf, cellC)),
    3 := the (fs-consTableS-S'-n-minus-off-tagGf (tagLf, consDataC)),
    4 := Intg nargs',
    5 := Intg 0],
  safeP, sigSafeMain,
  the (cdm (l, i)) + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1
+ 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1,
  ag, bd]]))
apply (unfold exec-all-def)
apply (simp del: append.append-Cons)
apply (erule rtrancl-trans)
prefer 2
apply (rule r-into-rtrancl)
apply (clarify)
apply (unfold JVMExec.exec.simps)
apply (unfold Let-def)
apply (subgoal-tac fst (snd (snd (snd (snd (the (method' (P', safeP) sigSafe-
Main)))))) !
  (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc
(Suc (Suc (Suc (Suc
  (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (the (cdm
(l, i)))))))))))))))))))))) =
  Putstatic topf stackC) apply (simp (no-asm-simp))
apply (erule nth-via-drop-append)

apply (drule drop-Suc-append)
apply (drule drop-append-3)
apply (subgoal-tac
  0 ≤ the-Intg (vs!5) ∧
  0 ≤ the-Intg (the (sh (stackC, topf))) - the-Intg (vs!4) ∧
  the-Intg (vs!5) ≤ the-Intg (vs!4))
apply (elim conjE)
prefer 2 apply (rule MATCH-bounds)
apply (frule-tac inih=inih and
  ref=(ag,bd) and
  nargs=nargs' and
  i = 0 and
  l = laa and
  h = h and
  ty = ty and
  m = m and
  S' = S' and
  vs=vs[Suc 0 := the (S' (nat (n - int off))),
    6 := Addr laa,

```



```

      safeP, sigSafeMain,
      the (cdm (l, i)) + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1
+ 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1
      + 1,
      ag, bd]))
apply (unfold exec-all-def)
apply (simp del: append.append-Cons)
apply (erule rtrancl-trans)
prefer 2
apply (rule r-into-rtrancl)
apply (clarify)
apply (unfold JVMEexec.exec.simps)
apply (unfold Let-def)
apply (subgoal-tac fst (snd (snd (snd (snd (the (method' (P', safeP) sigSafe-
Main)))))) !
      (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc
(Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc
      (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (the
(cdm (l, i)))))))))))))))))))))))))))))) =
      Load 4) apply (simp (no-asm-simp))
apply (simp, erule nth-via-drop)

apply (drule drop-Suc-append)
apply (subgoal-tac
  P† (None,
    sha((stackC, topf) ↦ Intg (n + nargs')),
    h(la ↦
      Arr ty m
      (S' ++ map-of (zip [Suc (nat n)..<nat (n + nargs' + 1)]
        (map (λj. argCell fs-S'-n-minus-off (Suc j)) [nat nargs'>..0])))),
    inih,
    [([Intg nargs'],
      vs[Suc 0 := the (S' (nat (n - int off))),
        6 := Addr laa,
        2 := the (fs-S'-n-minus-off (tagGf, cellC)),
        3 := the (fs-consTableS-S'-n-minus-off-tagGf (tagLf, consDataC)),
        4 := Intg nargs',
        5 := Intg nargs'],
      safeP, sigSafeMain,
      the (cdm (l, i)) + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1
+ 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1
      + 1,
      ag, bd]))
  -jvm→
  (None,
    sha((stackC, topf) ↦ Intg (n + nargs')),
    h(la ↦
      Arr ty m

```

```

      (S' ++ map-of (zip [Suc (nat n)..<nat (n + nargs' + 1)]
        (map (λj. argCell fs-S'-n-minus-off (Suc j)) [nat nargs'>..0]))),
    inih,
    [([Intg nargs',
      Intg nargs']),
      vs[Suc 0 := the (S' (nat (n - int off))),
        6 := Addr laa,
        2 := the (fs-S'-n-minus-off (tagGf, cellC)),
        3 := the (fs-consTableS-S'-n-minus-off-tagGf (tagLf, consDataC)),
        4 := Intg nargs',
        5 := Intg nargs']),
      safeP, sigSafeMain,
      the (cdm (l, i)) + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 +
+ 1 + 1 + 1 + 1 + 1 +
      1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1,
      ag, bd]]))
apply (unfold exec-all-def)
apply (simp del: append.append-Cons)
apply (erule rtrancl-trans)
prefer 2
apply (rule r-into-rtrancl)
apply (clarify)
apply (unfold JVMExec.exec.simps)
apply (unfold Let-def)
apply (subgoal-tac fst (snd (snd (snd (snd (the (method' (P', safeP) sigSafe-
Main)))))) !
      (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc
(Suc (Suc (Suc (Suc
      (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc
(the (cdm (l, i)))))))))))))))))))))))))) =
      Load 5) apply (simp (no-asm-simp)))
apply (simp, erule nth-via-drop)

apply (drule drop-Suc-append)
apply (subgoal-tac
  P ⊢ (None,
    sha((stackC, topf) ↦ Intg (n + nargs')),
    h(la ↦
      Arr ty m
      (S' ++ map-of (zip [Suc (nat n)..<nat (n + nargs' + 1)]
        (map (λj. argCell fs-S'-n-minus-off (Suc j)) [nat nargs'>..0]))),
    inih,
    [([Intg nargs',
      Intg nargs']),
      vs[Suc 0 := the (S' (nat (n - int off))),
        6 := Addr laa,
        2 := the (fs-S'-n-minus-off (tagGf, cellC)),
        3 := the (fs-consTableS-S'-n-minus-off-tagGf (tagLf, consDataC)),

```

```

      4 := Intg nargs',
      5 := Intg nargs'],
    safeP, sigSafeMain,
    the (cdm (l, i)) + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1
+ 1 + 1 + 1 + 1 + 1 + 1 + 1 +
    1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1,
    ag, bd)])
-jvm→
(None,
sha((stackC, topf) ↦ Intg (n + nargs')),
h(la ↦
  Arr ty m
  (S' ++ map-of (zip [Suc (nat n)..<nat (n + nargs' + 1)]
    (map (λj. argCell fs-S'-n-minus-off (Suc j)) [nat nargs'>..0])))),
inih,
[([],
vs[Suc 0 := the (S' (nat (n - int off))),
6 := Addr laa,
2 := the (fs-S'-n-minus-off (tagGf, cellC)),
3 := the (fs-consTableS-S'-n-minus-off-tagGf (tagLf, consDataC)),
4 := Intg nargs',
5 := Intg nargs'],
safeP, sigSafeMain,
the (cdm (l, i)) + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1
+ 1 + 1 + 1 + 1 +
    1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + nat
labelEndLoop,
ag, bd]]))
apply (unfold exec-all-def)
apply (simp del: append.append-Cons)
apply (erule rtrancl-trans)
prefer 2
apply (rule r-into-rtrancl)
apply (clarify)
apply (unfold JVMEexec.exec.simps)
apply (unfold Let-def)
apply (subgoal-tac fst (snd (snd (snd (snd (the (method' (P', safeP) sigSafe-
Main)))))) !
  (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc
(Suc (Suc (Suc (Suc (Suc
  (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc
(the (cdm (l, i)))))))))))))))))))))))))) =
    Ifcmp GreaterEqual labelEndLoop) apply (simp (no-asm-simp))
apply (simp add: labelEndLoop-def)
apply (simp, erule nth-via-drop)

apply (frule go-End)
apply (subgoal-tac

```

```

P⊢ (None,
  sha((stackC, topf) ↦ Intg (n + nargs')),
  h(la ↦
    Arr ty m
      (S' ++ map-of (zip [Suc (nat n)..<nat (n + nargs' + 1)]
        (map (λj. argCell fs-S'-n-minus-off (Suc j)) [nat nargs'>..0])))),
  inih,
  [([],
    vs[Suc 0 := the (S' (nat (n - int off))),
      6 := Addr laa,
      2 := the (fs-S'-n-minus-off (tagGf, cellC)),
      3 := the (fs-consTableS-S'-n-minus-off-tagGf (tagLf, consDataC)),
      4 := Intg nargs',
      5 := Intg nargs'],
    safeP, sigSafeMain,
    the (cdm (l, i)) + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1
+ 1 + 1 + 1 +
      1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + nat
labelEndLoop,
    ag, bd)])
  -jvm→
  (None,
  sha((stackC, topf) ↦ Intg (n + nargs')),
  h(la ↦
    Arr ty m
      (S' ++ map-of (zip [Suc (nat n)..<nat (n + nargs' + 1)]
        (map (λj. argCell fs-S'-n-minus-off (Suc j)) [nat nargs'>..0])))),
  inih,
  [([the (fs-consTableS-S'-n-minus-off-tagGf (tagLf, consDataC))],
    vs[Suc 0 := the (S' (nat (n - int off))),
      6 := Addr laa,
      2 := the (fs-S'-n-minus-off (tagGf, cellC)),
      3 := the (fs-consTableS-S'-n-minus-off-tagGf (tagLf, consDataC)),
      4 := Intg nargs',
      5 := Intg nargs'],
    safeP, sigSafeMain,
    the (cdm (l, i)) + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1
+ 1 + 1 + 1 +
      1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + 1 + nat
labelEndLoop + 1,
    ag, bd)))]
apply (unfold exec-all-def)
apply (simp del: append.append-Cons)
apply (erule rtrancl-trans)
prefer 2
apply (rule r-into-rtrancl)
apply (clarify)
apply (unfold JVMEExec.exec.simps)
apply (unfold Let-def)

```





```

      nat (the-Intg (the (fs-consTableS-S'-n-minus-off-tagGf (tagLf,
consDataC))))))))))))))))))))))))))))))))))))))))))))))))))))))))))))))
      ag, bd)))
apply (unfold exec-all-def)
apply (simp del: append.append-Cons)
apply simp

apply (rule r-into-rtrancl)
apply (clarify)
apply (unfold JVMEexec.exec.simps)
apply (unfold Let-def)
apply (subgoal-tac (fst (snd (snd (snd (snd (the (method' (P', safeP) sigSafe-
Main))))))) !
      Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc
      (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc
(Suc
      (Suc (Suc (Suc (Suc (the (cdm (l, i)) + nat labelEndLoop))))))))))))))))))))))))))))))))))))))))))
=
      Tableswitch 0 (int (length ps - Suc 0))
      (map (λn. trAddr n (the (cdm (l, i)) + incMatch))
      (map (λp. the (cdm' (p, 0))) ps))) apply (simp (no-asm-simp))
apply (subgoal-tac
  ¬ the-Intg (the (fs-consTableS-S'-n-minus-off-tagGf (tagLf, consDataC))) < 0
  ∧
  ¬ int (length ps - Suc 0) < the-Intg (the (fs-consTableS-S'-n-minus-off-tagGf
(tagLf, consDataC))),simp)
apply (rule tag-bounds)
apply (simp, erule nth-via-drop)

apply (rule-tac x=vs[Suc 0 := the (S' (nat (n - int off))),
      6 := Addr laa,
      2 := the (fs-S'-n-minus-off (tagGf, cellC)),
      3 := the (fs-consTableS-S'-n-minus-off-tagGf (tagLf,
consDataC)),
      4 := Intg nargs',
      5 := Intg nargs'] in exI)
apply (rule-tac x=Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc
(Suc
      (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc (Suc
      (Suc (Suc (Suc (Suc (the (cdm (l, i)) + nat labelEndLoop +
      nat (map (λn. trAddr n (the (cdm (l, i)) + incMatch))
      (map (λp. the (cdm' (p, 0))) ps) !
      nat (the-Intg (the (fs-consTableS-S'-n-minus-off-tagGf (tagLf,
consDataC))
      )))))))))))))))))))))))))))))))))))))))) in exI)
apply (rule conjI)

```

```

apply (rule-tac x=ag in exI)
apply (rule-tac x=bd in exI)
apply (rule refl)
apply (rule-tac x=la in exI)
apply (rule conjI, simp add: Sf-def add: topf-def)
apply (rule-tac x=ty in exI)
apply (rule-tac x=m in exI)
apply (rule-tac x=S' ++ map-of (zip [Suc (nat n)..<nat (n + nargs' + 1)]
    (map (λj. argCell fs-S'-n-minus-off (Suc j))
      [nat nargs'>..0]))) in exI)
apply (rule-tac x=(n + nargs') in exI)
apply (rule-tac x=l' in exI)
apply (rule conjI, assumption)
apply (rule-tac x=regS in exI)
apply (rule-tac x=l'' in exI)
apply (rule conjI, assumption)
apply (rule conjI, assumption)
apply (rule conjI) apply (rule activeCells)
    apply (simp add: heapC-def add: stackC-def)
apply (rule conjI) apply (rule activeCells-2)
    apply (simp add: heapC-def add: stackC-def, simp)
apply (rule conjI) apply (rule activeCells-2)
    apply (simp add: heapC-def add: stackC-def, simp)
apply (rule conjI, simp)
apply (rule conjI, simp)
apply (rule conjI, simp add: heapC-def add: stackC-def)
apply (rule conjI)
apply (rule-tac x=m' in exI, simp)
apply (rule conjI, simp add: dirCellC-def add: stackC-def)
apply (rule-tac x=m'' in exI)
apply (rule-tac x=d in exI)
apply (rule conjI, simp)
apply (rule-tac x=g in exI, simp add: heapC-def add: stackC-def)
apply (fold heapC-def, fold stackC-def)
apply (rule conjI) apply (rule equivH-MATCH, simp, assumption+, force)
apply (rule conjI) apply (rule maxPush-MATCH, assumption)
apply (rule conjI) apply (rule equivS-MATCH-axiom)
apply (rule pc-MATCH)
done

```

end

## 26 Main correctness theorem of the translation SVM to JVM

```

theory mainTheoremSVM2JVM

```

```

imports CertifSVM2JVM dem-translation dem-PUSHCONT dem-POPCONT dem-PRIMOP

```

*dem-SLIDE dem-MATCHN dem-BUILDENV dem-CALL*  
*dem-DECREGION*  
*dem-BUILDCLS dem-REUSE dem-MATCH*

**begin**

This theorem certifies that, whenever the SVM and the JVM are started in equivalent states, the SVM executes its next instruction, and the JVM executes its translation to bytecode, then both machines arrive at equivalent states.

**theorem** *correctSVM2JVM* :

$$\begin{aligned} & \llbracket (P', cdm, ctm, com) = trSVM2JVM P; \\ & \quad cdm, ctm, com \vdash S1 \triangleq S1'; \\ & \quad execSVM P S1 = Either.Right S2 \\ & \rrbracket \implies \\ & \exists S2'. P' \vdash S1' -jvm \rightarrow S2' \wedge cdm, ctm, com \vdash S2 \triangleq S2' \end{aligned}$$

**apply** (*case-tac P*)  
**apply** (*rename-tac cs triple*)  
**apply** (*case-tac cs*)  
**apply** (*rename-tac svms ctnmap*)  
**apply** (*case-tac triple*)  
**apply** (*rename-tac ini pair*)  
**apply** (*case-tac pair*)  
**apply** (*rename-tac ct st*)  
**apply** (*unfold execSVM-def*)  
**apply** (*case-tac S1*)  
**apply** (*rename-tac h rest1*)  
**apply** (*case-tac rest1*)  
**apply** (*rename-tac k0 rest2*)  
**apply** (*case-tac rest2*)  
**apply** (*rename-tac pc S*)  
**apply** (*case-tac pc*)  
**apply** (*rename-tac l i*)  
**apply** (*case-tac h*)  
**apply** (*rename-tac hm k*)  
**apply** (*clarsimp*)

**apply** (*subgoal-tac l : dom (map-of svms)*  
 $\wedge i < length (fst (the (map-of svms l)))$ )  
**prefer** 2 **apply** (*erule PC-good,simp*)  
**apply** (*elim conjE*)  
**apply** (*subgoal-tac  $\exists j$  .*  
 $j < length svms \wedge (l, the (map-of svms l)) = svms ! j$ )  
**prefer** 2 **apply** (*rule map-of-distinct2*) **apply** *force*  
**apply** (*erule exE,elim conjE*)  
**apply** (*case-tac the (map-of svms l),rename-tac seq fn,clarify*)

**apply** (*subgoal-tac* ( $\exists$  *cdm' ctm' pcc inss bytecode'* .  
*inss = trInstr (the (cdm (l,i))) cdm' ctm' com pcc*  
*(fst (the (map-of svms l)) ! i)*  
 $\wedge$  *drop (the (cdm (l,i))) (extractBytecode P') = inss @ bytecode'*  
 $\wedge$  *cdm'  $\subseteq_m$  cdm*))  
**prefer** 2 **apply** (*rule fun-SVM2JVM*)  
**apply** *simp*  
**apply** *simp*  
**apply** (*rule sym, simp*)  
**apply** *simp*  
**apply** *simp*  
**apply** *simp*  
**apply** *simp*

**apply** (*elim exE, elim conjE*)  
**apply** (*simp, elim conjE, clarify*)

**apply** (*case-tac fst (the (map-of svms l)) ! i*)

**apply** (*erule execSVMInstr-DECREGION*)  
**apply** (*simp, simp, simp, force*)

**apply** (*erule execSVMInstr-POPCONT*)  
**apply** (*simp, simp, simp, force*)

**apply** (*erule execSVMInstr-PUSHCONT*)  
**apply** (*simp, simp, simp, force*)

**apply** (*erule execSVMInstr-COPY*)  
**apply** (*simp, simp, simp, force*)

**apply** (*erule execSVMInstr-REUSE*)  
**apply** (*simp, simp, simp, force*)

**apply** (*erule execSVMInstr-CALL*)  
**apply** (*simp, simp, simp, force, simp*)

**apply** (*erule execSVMInstr-PRIMOP*)  
**apply** (*simp, simp, simp, force*)

```
apply (erule execSVMInstr-MATCH)  
apply (simp,force,simp,force,simp)
```

```
apply (erule execSVMInstr-MATCHD)  
apply (simp,force,simp,force)
```

```
apply (erule execSVMInstr-MATCHN)  
apply (simp,force,simp,force,simp)
```

```
apply (erule execSVMInstr-BUILDENV)  
apply (simp,simp,simp,force)
```

```
apply (erule execSVMInstr-BUILDCLS)  
apply (simp,force,simp,force)
```

```
apply (frule-tac m=nat1 and  
          n=nat2 in execSVMInstr-SLIDE)  
apply (simp+)  
done  
  
end
```